

Multilevel Monte Carlo Path Simulation

Mike Giles

`mike.giles@maths.ox.ac.uk`

Oxford University Mathematical Institute
Oxford-Man Institute of Quantitative Finance

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SDEs in Finance

In computational finance, stochastic differential equations are used to model the behaviour of

- stocks
- interest rates
- exchange rates
- weather
- electricity/gas demand
- crude oil prices
- ...

The stochastic term accounts for the uncertainty of unpredictable day-to-day events.

SDEs in Finance

These models are then used to calculate “fair” prices for a huge range of financial options:

- an option to sell a stock portfolio at a specific price in 2 years time
- an option to buy aviation fuel at a specific price in 6 months time
- an option to sell US dollars at a specific exchange rate in 3 years time

In most cases, the buyer of the financial option is trying to reduce their risk.

SDEs in Finance

Examples:

- Geometric Brownian motion (Black-Scholes model for stock prices)

$$dS = r S dt + \sigma S dW$$

- Cox-Ingersoll-Ross model (interest rates)

$$dr = \alpha(b - r) dt + \sigma \sqrt{r} dW$$

- Heston stochastic volatility model (stock prices)

$$dS = r S dt + \sqrt{V} S dW_1$$

$$dV = \lambda(\sigma^2 - V) dt + \xi \sqrt{V} dW_2$$

with correlation ρ between dW_1 and dW_2

Generic Problem

Stochastic differential equation with general drift and volatility terms: SDE with general drift and volatility terms:

$$dS(t) = a(S, t) dt + b(S, t) dW(t)$$

$W(t)$ is a Wiener variable with the properties that for any $q < r < s < t$, $W(t) - W(s)$ is Normally distributed with mean 0 and variance $t - s$, independent of $W(r) - W(q)$.

In many finance applications, we want to compute the expected value of an option dependent on the terminal state

$$P \equiv f(S(T))$$

with a uniform Lipschitz bound,

$$|f(U) - f(V)| \leq c \|U - V\|, \quad \forall U, V.$$

Standard MC Approach

Euler discretisation with timestep h :

$$\widehat{S}_{n+1} = \widehat{S}_n + a(\widehat{S}_n, t_n) h + b(\widehat{S}_n, t_n) \Delta W_n$$

where ΔW_n are Normal with mean 0, variance h .

Simplest estimator for expected payoff is an average of N independent path simulations:

$$\widehat{Y} = N^{-1} \sum_{i=1}^N \widehat{P}^{(i)}$$

where $\widehat{P} \equiv f(\widehat{S}_{T/h})$ is an approximation to $P \equiv f(S(T))$ for a given Brownian path $W(t)$.

Standard MC Approach

The mean square error is defined as

$$\begin{aligned}\mathbb{E} \left[\left(\hat{Y} - \mathbb{E}[P] \right)^2 \right] &= \mathbb{E} \left[\left(\hat{Y} - \mathbb{E}[\hat{P}] + \mathbb{E}[\hat{P}] - \mathbb{E}[P] \right)^2 \right] \\ &= \mathbb{E} \left[\left(\hat{Y} - \mathbb{E}[\hat{P}] \right)^2 \right] + \left(\mathbb{E}[\hat{P}] - \mathbb{E}[P] \right)^2 \\ &= N^{-1} \mathbb{V}[\hat{P}] + \left(\mathbb{E}[\hat{P}] - \mathbb{E}[P] \right)^2\end{aligned}$$

- first term is due to variance of estimator
- second term is due to bias due to finite timestep
– weak convergence

Standard MC Approach

Weak convergence:

- error in the expected value, $\mathbb{E}[\hat{P}] - \mathbb{E}[P]$
- most important error in most applications
- $O(h)$ for the Euler discretisation

Strong convergence:

- error in path approximation

$$\sqrt{\mathbb{E} \left[\left\| \hat{S}_{T/h} - S(T) \right\|^2 \right]} \quad \text{or} \quad \sqrt{\mathbb{E} \left[\max_{0 < t < T} \left\| \hat{S}(t) - S(t) \right\|^2 \right]}$$

- usually not relevant, but important for multilevel method
- $O(h^{1/2})$ for the Euler discretisation

Standard MC Approach

Combined mean-square-error is $O(N^{-1} + h^2)$.

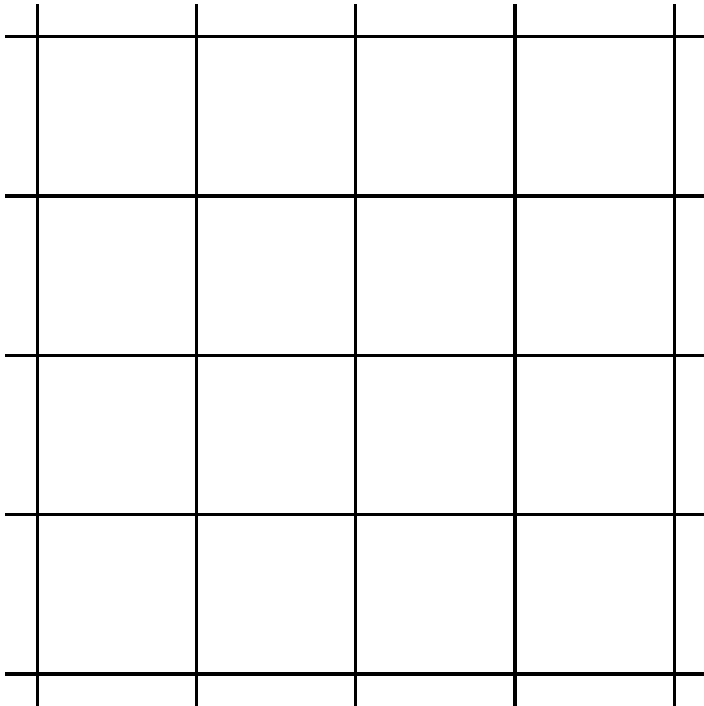
To make this equal to ε^2 requires

$$N = O(\varepsilon^{-2}), \quad h = O(\varepsilon) \quad \implies \quad \text{cost} = O(N h^{-1}) = O(\varepsilon^{-3})$$

Aim is to improve this cost to $O(\varepsilon^{-2}(\log \varepsilon)^2)$, by combining simulations with different numbers of timesteps – same accuracy as finest calculations, but at a much lower computational cost.

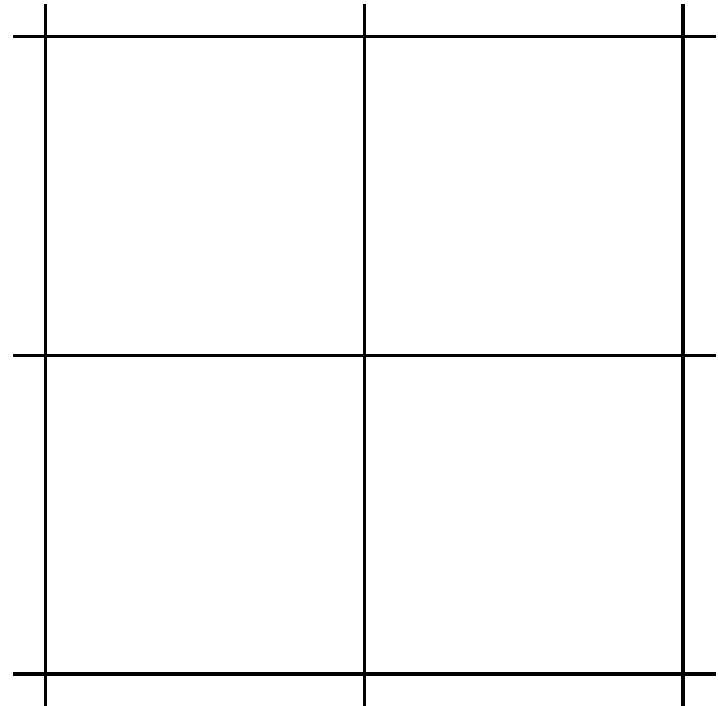
Multigrid

A powerful technique for solving PDE discretisations:



Fine grid

more accurate
more expensive



Coarse grid

less accurate
less expensive

Multigrid

Multigrid combines calculations on a nested sequence of grids to get the accuracy of the finest grid at a much lower computational cost.

We use a similar idea in Monte Carlo path calculations, combining simulations with different numbers of timesteps – same accuracy as finest calculations, but at a much lower computational cost.

Multilevel MC Approach

Consider multiple sets of simulations with different timesteps $h_l = 2^{-l} T$, $l = 0, 1, \dots, L$, and payoff \hat{P}_l

$$\mathbb{E}[\hat{P}_L] = \mathbb{E}[\hat{P}_0] + \sum_{l=1}^L \mathbb{E}[\hat{P}_l - \hat{P}_{l-1}]$$

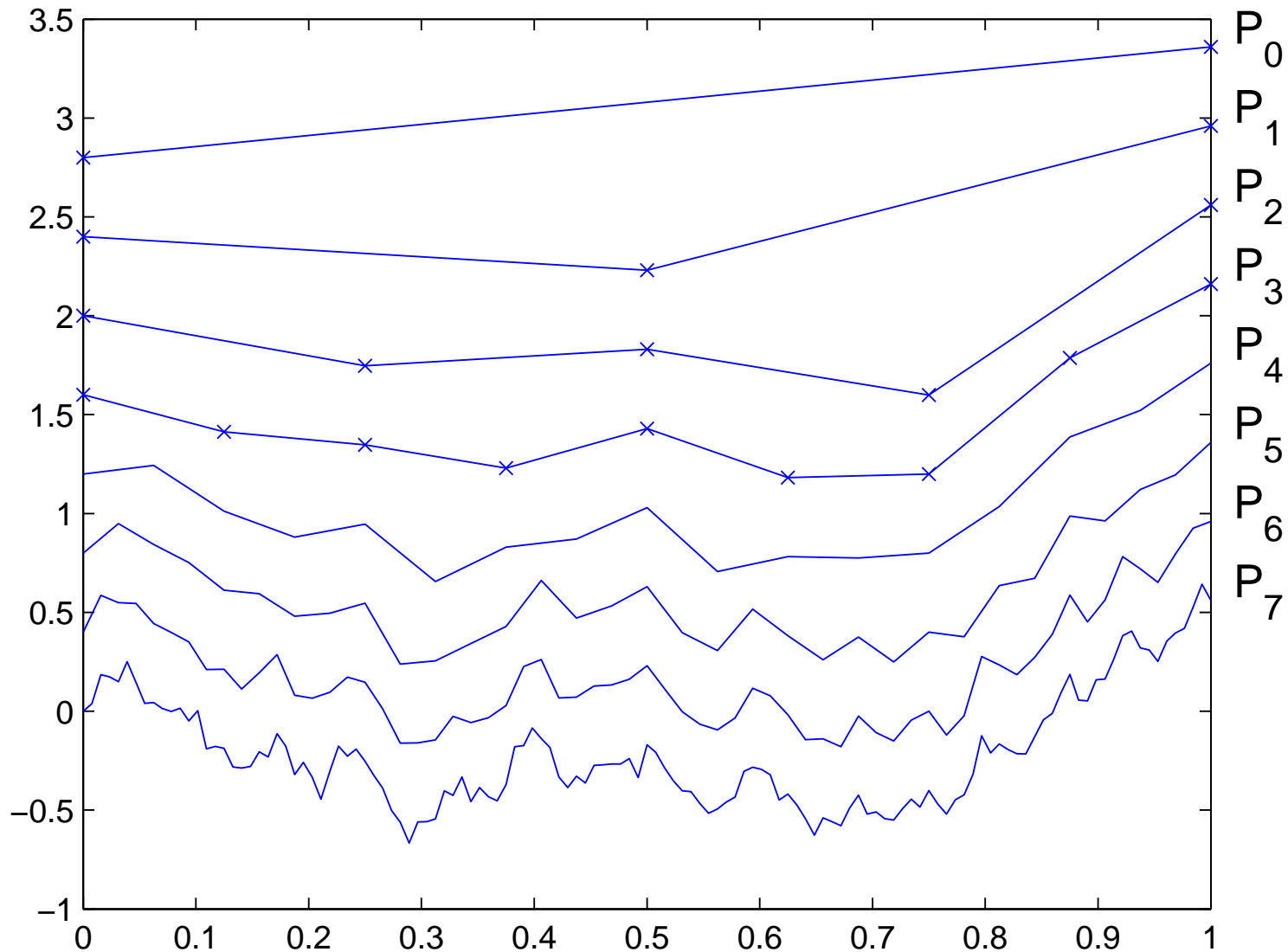
Expected value is same – aim is to reduce variance of estimator for a fixed computational cost.

Key point: approximate $\mathbb{E}[\hat{P}_l - \hat{P}_{l-1}]$ using N_l simulations with \hat{P}_l and \hat{P}_{l-1} obtained using same Brownian path.

$$\hat{Y}_l = N_l^{-1} \sum_{i=1}^{N_l} \left(\hat{P}_l^{(i)} - \hat{P}_{l-1}^{(i)} \right)$$

Multilevel MC Approach

Discrete Brownian path at different levels



Multilevel MC Approach

- each level adds more detail to Brownian path
- $\mathbb{E}[\hat{P}_l - \hat{P}_{l-1}]$ reflects impact of that extra detail on the payoff
- different timescales handled by different levels
 - similar to different wavelengths being handled by different grids in multigrid

Multilevel MC Approach

Using independent paths for each level, the variance of the combined estimator is

$$\mathbb{V} \left[\sum_{l=0}^L \hat{Y}_l \right] = \sum_{l=0}^L N_l^{-1} V_l, \quad V_l \equiv \mathbb{V}[\hat{P}_l - \hat{P}_{l-1}],$$

and the computational cost is proportional to $\sum_{l=0}^L N_l h_l^{-1}$.

Hence, the variance is minimised for a fixed computational cost by choosing N_l to be proportional to $\sqrt{V_l h_l}$.

The constant of proportionality can be chosen so that the combined variance is $O(\varepsilon^2)$.

Multilevel MC Approach

For the Euler discretisation and the Lipschitz payoff function

$$\mathbb{V}[\hat{P}_l - P] = O(h_l) \quad \Longrightarrow \quad \mathbb{V}[\hat{P}_l - \hat{P}_{l-1}] = O(h_l)$$

and the optimal N_l is asymptotically proportional to h_l .

To make the combined variance $O(\varepsilon^2)$ requires

$$N_l = O(\varepsilon^{-2} L h_l).$$

To make the bias $O(\varepsilon)$ requires

$$L = \log_2 \varepsilon^{-1} + O(1) \quad \Longrightarrow \quad h_L = O(\varepsilon).$$

Hence, we obtain an $O(\varepsilon^2)$ MSE for a computational cost which is $O(\varepsilon^{-2} L^2) = O(\varepsilon^{-2} (\log \varepsilon)^2)$.

Results

Geometric Brownian motion:

$$dS = r S dt + \sigma S dW, \quad 0 < t < T,$$

$$T = 1, \quad S(0) = 1, \quad r = 0.05, \quad \sigma = 0.2$$

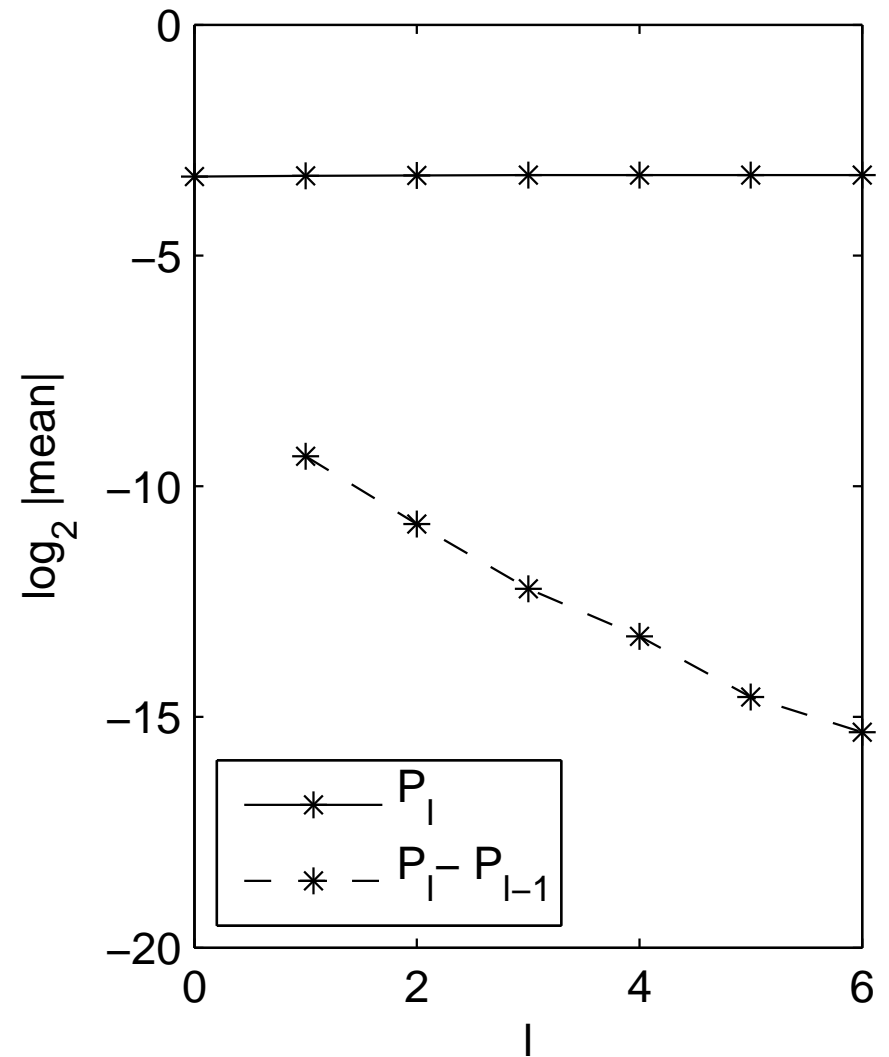
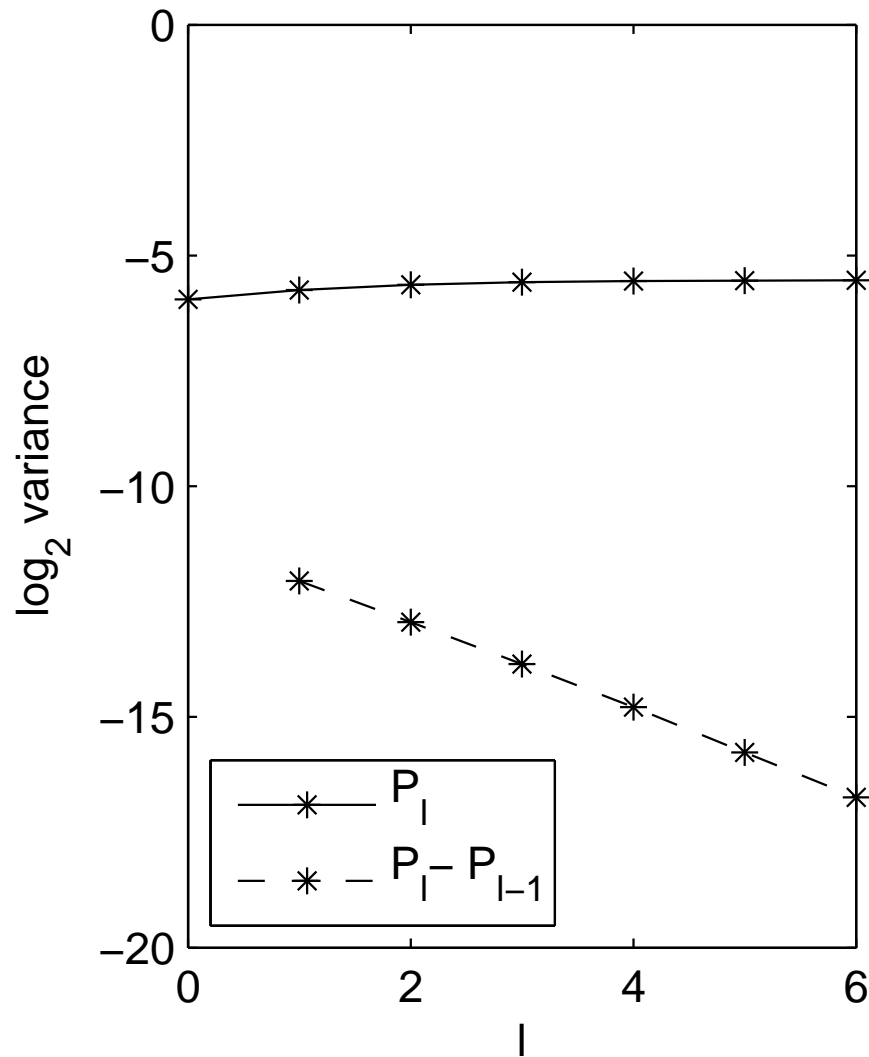
European call option with discounted payoff

$$\exp(-rT) \max(S(T) - K, 0)$$

with $K = 1$.

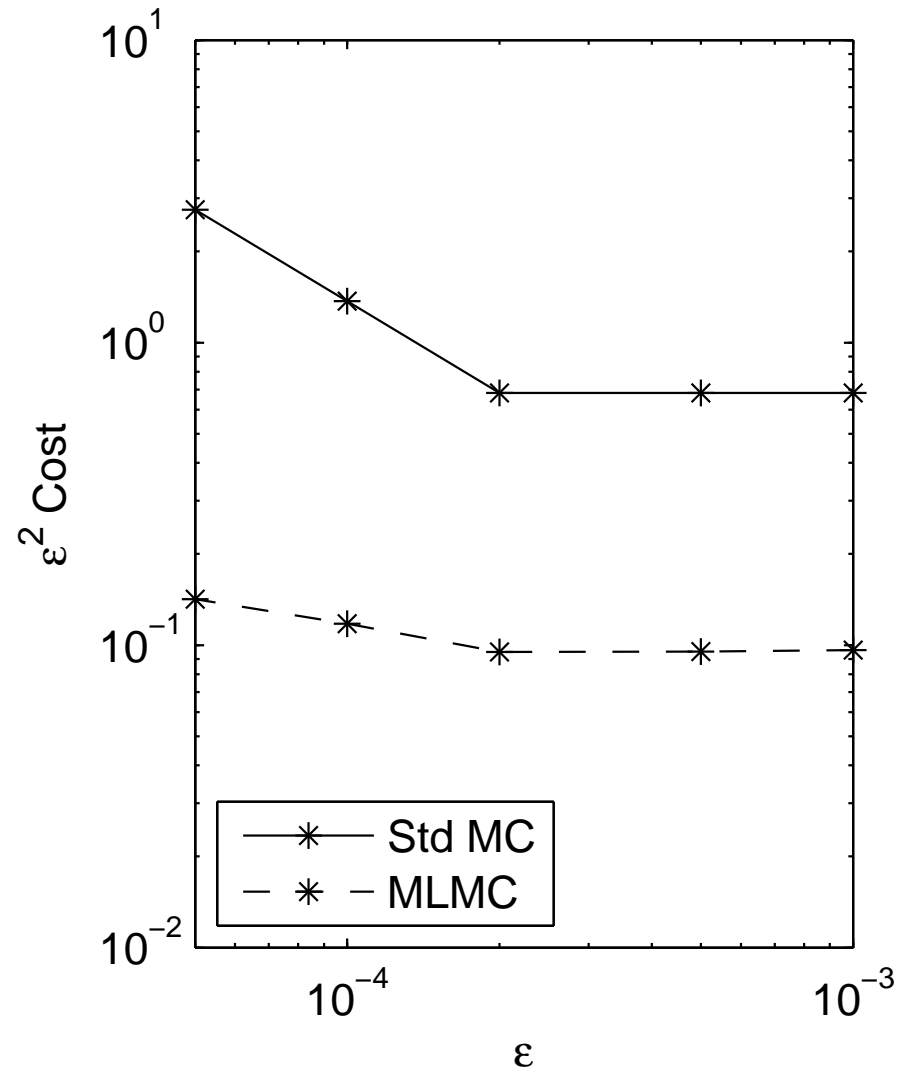
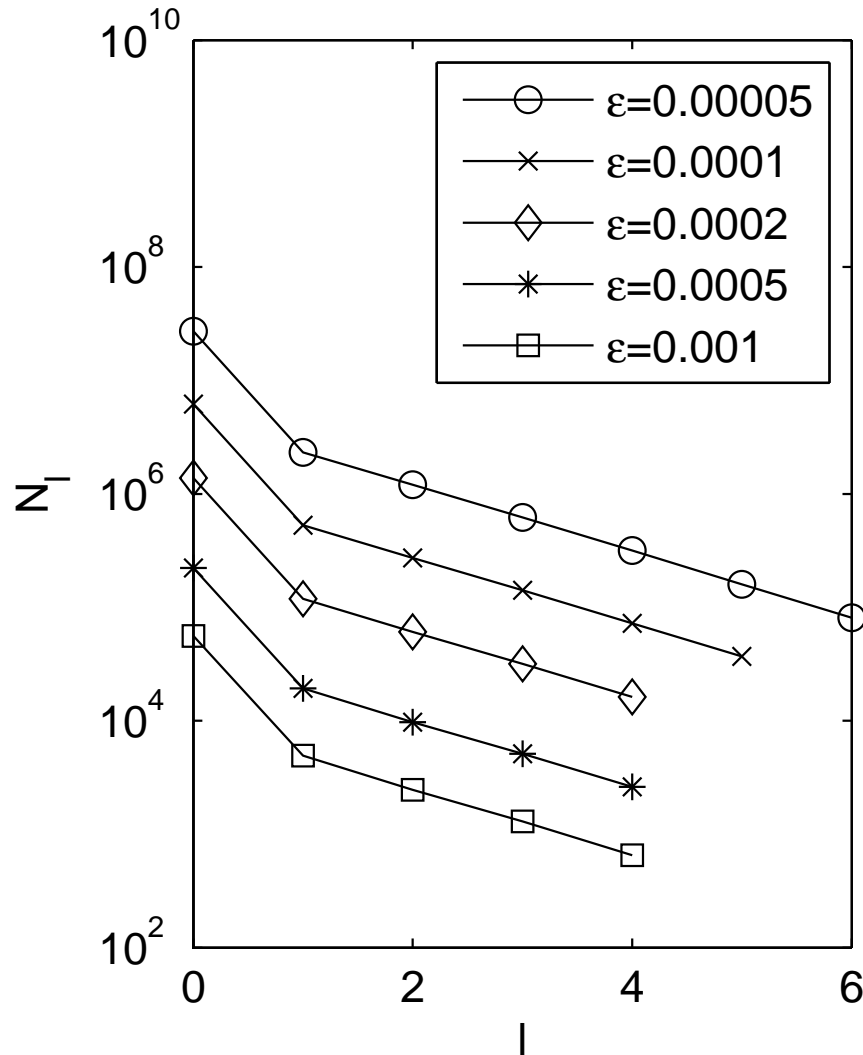
MLMC Results

GBM: European call, $\exp(-rT) \max(S(T) - K, 0)$



MLMC Results

GBM: European call, $\exp(-rT) \max(S(T) - K, 0)$



Multilevel MC Approach

Theorem: Let P be a functional of the solution of a stochastic o.d.e., and \widehat{P}_l the discrete approximation using a timestep $h_l = M^{-l} T$.

If there exist independent estimators \widehat{Y}_l based on N_l Monte Carlo samples, and positive constants $\alpha \geq \frac{1}{2}$, β , c_1 , c_2 , c_3 such that

$$i) \quad \left| \mathbb{E}[\widehat{P}_l - P] \right| \leq c_1 h_l^\alpha$$

$$ii) \quad \mathbb{E}[\widehat{Y}_l] = \begin{cases} \mathbb{E}[\widehat{P}_0], & l = 0 \\ \mathbb{E}[\widehat{P}_l - \widehat{P}_{l-1}], & l > 0 \end{cases}$$

$$iii) \quad \mathbb{V}[\widehat{Y}_l] \leq c_2 N_l^{-1} h_l^\beta$$

iv) C_l , the computational complexity of \widehat{Y}_l , is bounded by

$$C_l \leq c_3 N_l h_l^{-1}$$

Multilevel MC Approach

then there exists a positive constant c_4 such that for any $\varepsilon < e^{-1}$ there are values L and N_l for which the multi-level estimator

$$\hat{Y} = \sum_{l=0}^L \hat{Y}_l,$$

has Mean Square Error $MSE \equiv E \left[\left(\hat{Y} - E[P] \right)^2 \right] < \varepsilon^2$

with a computational complexity C with bound

$$C \leq \begin{cases} c_4 \varepsilon^{-2}, & \beta > 1, \\ c_4 \varepsilon^{-2} (\log \varepsilon)^2, & \beta = 1, \\ c_4 \varepsilon^{-2 - (1-\beta)/\alpha}, & 0 < \beta < 1. \end{cases}$$

Milstein Scheme

The theorem suggests use of Milstein approximation
– better strong convergence, same weak convergence

Generic scalar SDE:

$$dS(t) = a(S, t) dt + b(S, t) dW(t), \quad 0 < t < T.$$

Milstein scheme:

$$\hat{S}_{n+1} = \hat{S}_n + a h + b \Delta W_n + \frac{1}{2} b' b \left((\Delta W_n)^2 - h \right).$$

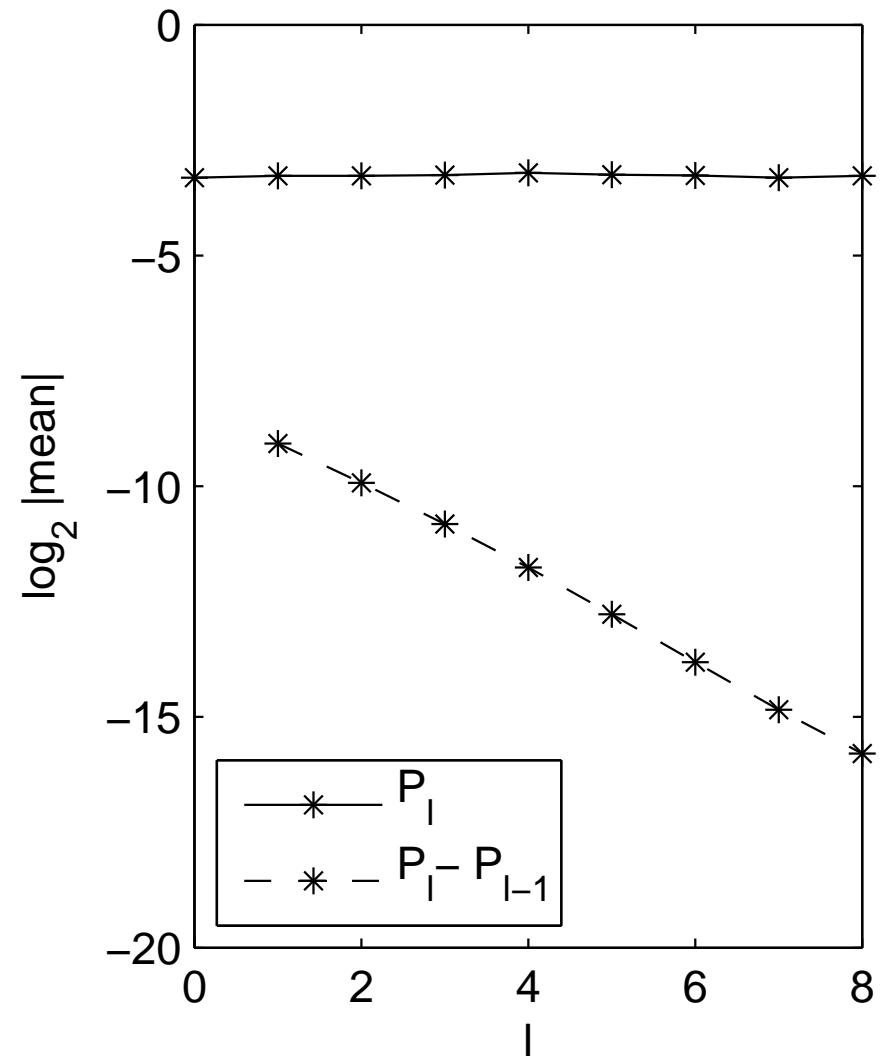
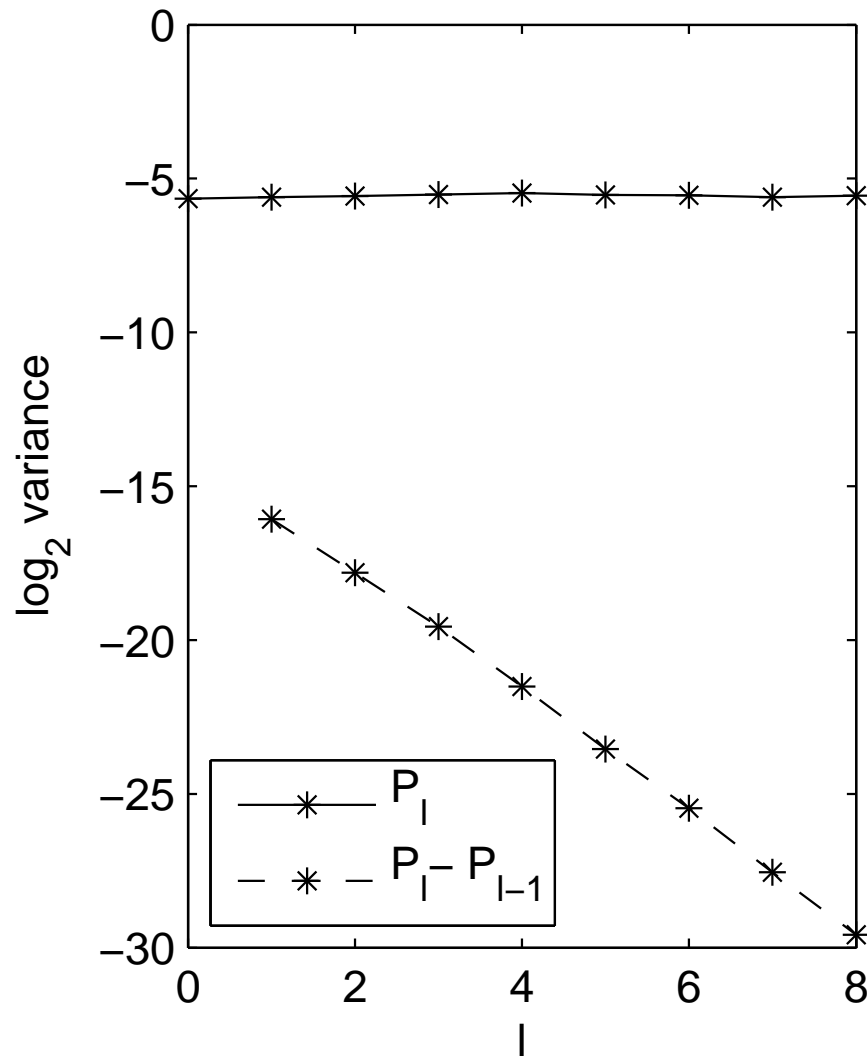
Milstein Scheme

In scalar case:

- $O(h)$ strong convergence
- $O(\varepsilon^{-2})$ complexity for Lipschitz payoffs – trivial
- $O(\varepsilon^{-2})$ complexity for more complex cases using carefully constructed estimators based on Brownian interpolation or extrapolation
 - digital, with discontinuous payoff
 - Asian, based on average
 - lookback and barrier, based on min/max

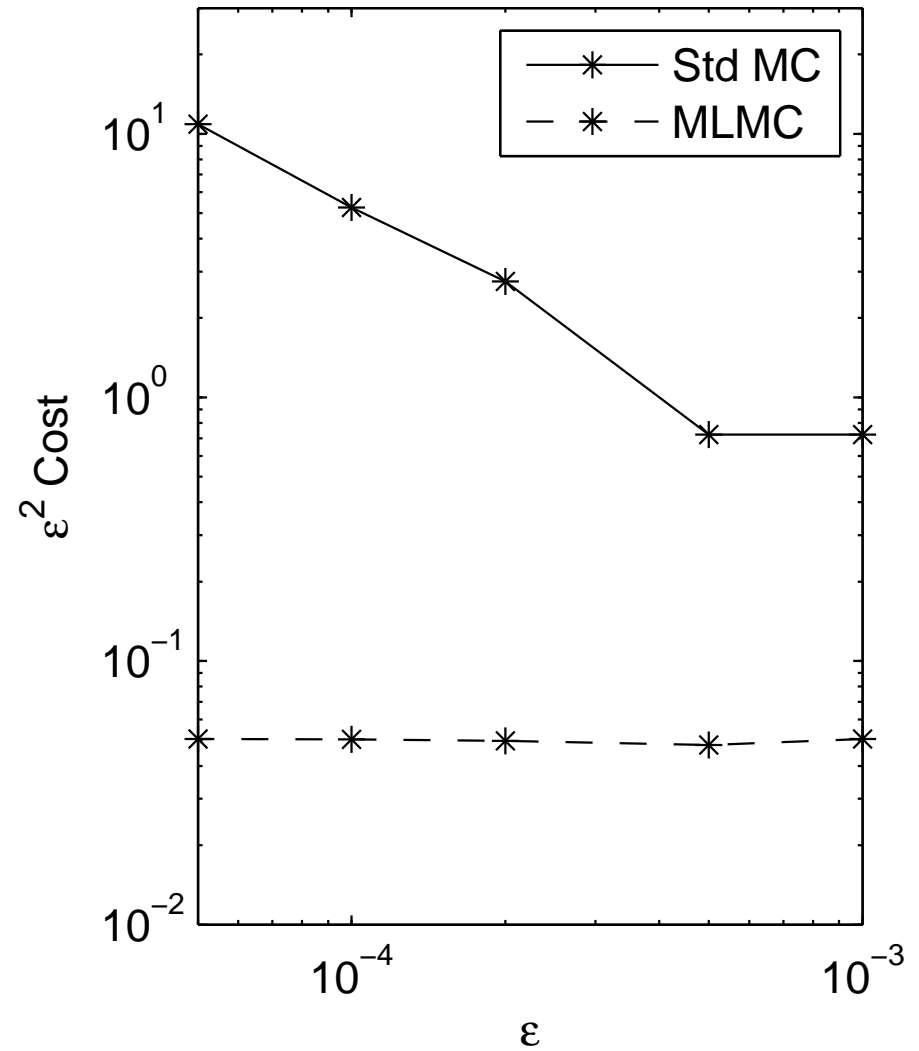
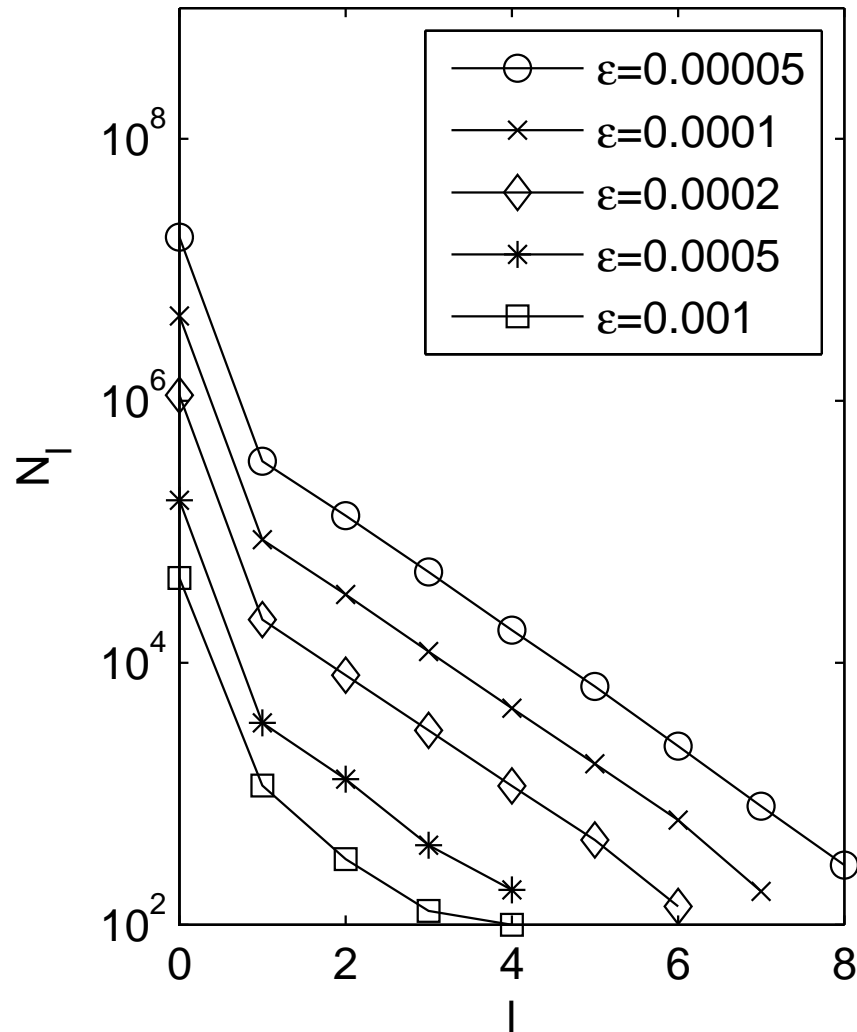
MLMC Results

GBM: European call, $\exp(-rT) \max(S(T) - K, 0)$



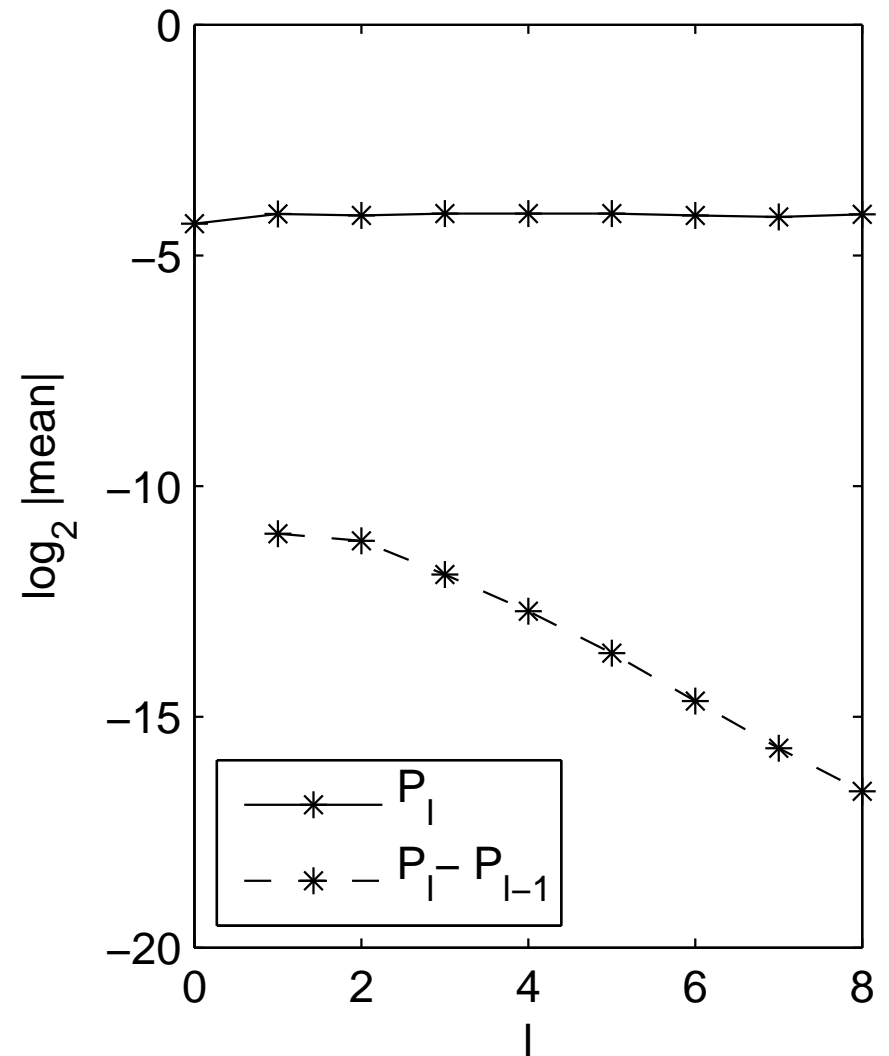
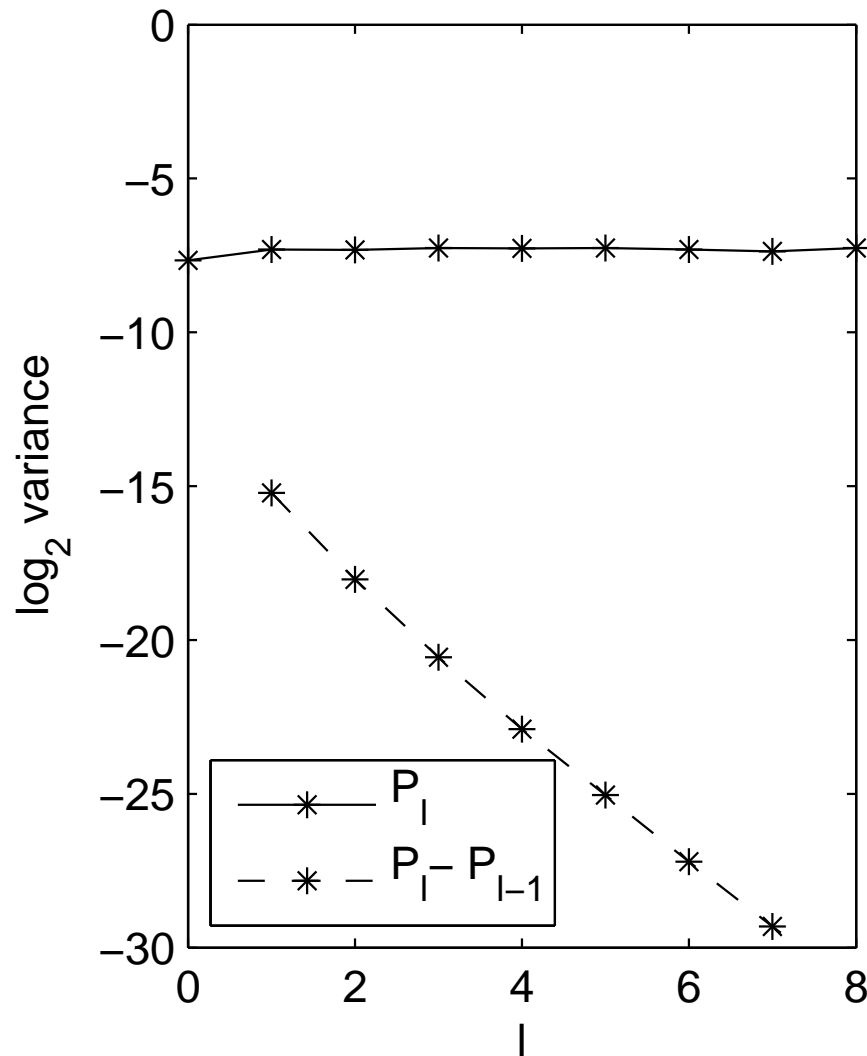
MLMC Results

GBM: European call, $\exp(-rT) \max(S(T) - K, 0)$



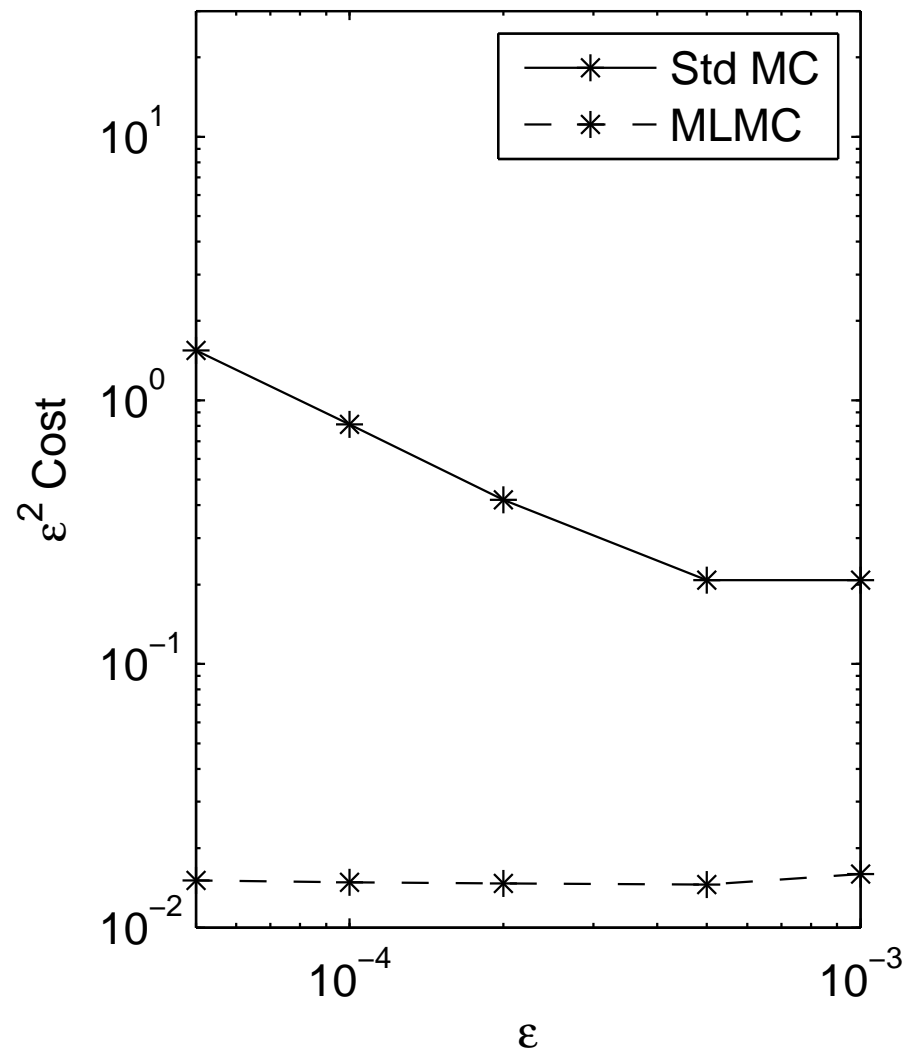
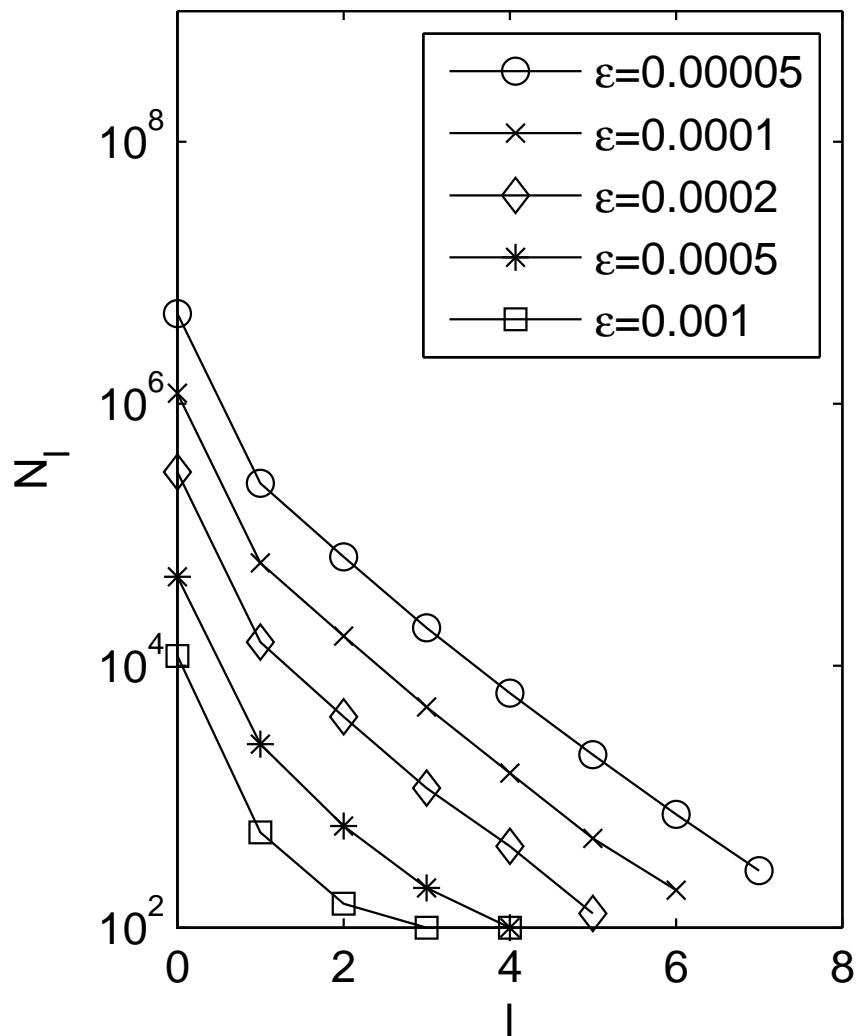
MLMC Results

GBM: Asian option, $\exp(-rT) \max(T^{-1} \int_0^T S(t) dt - 1, 0)$



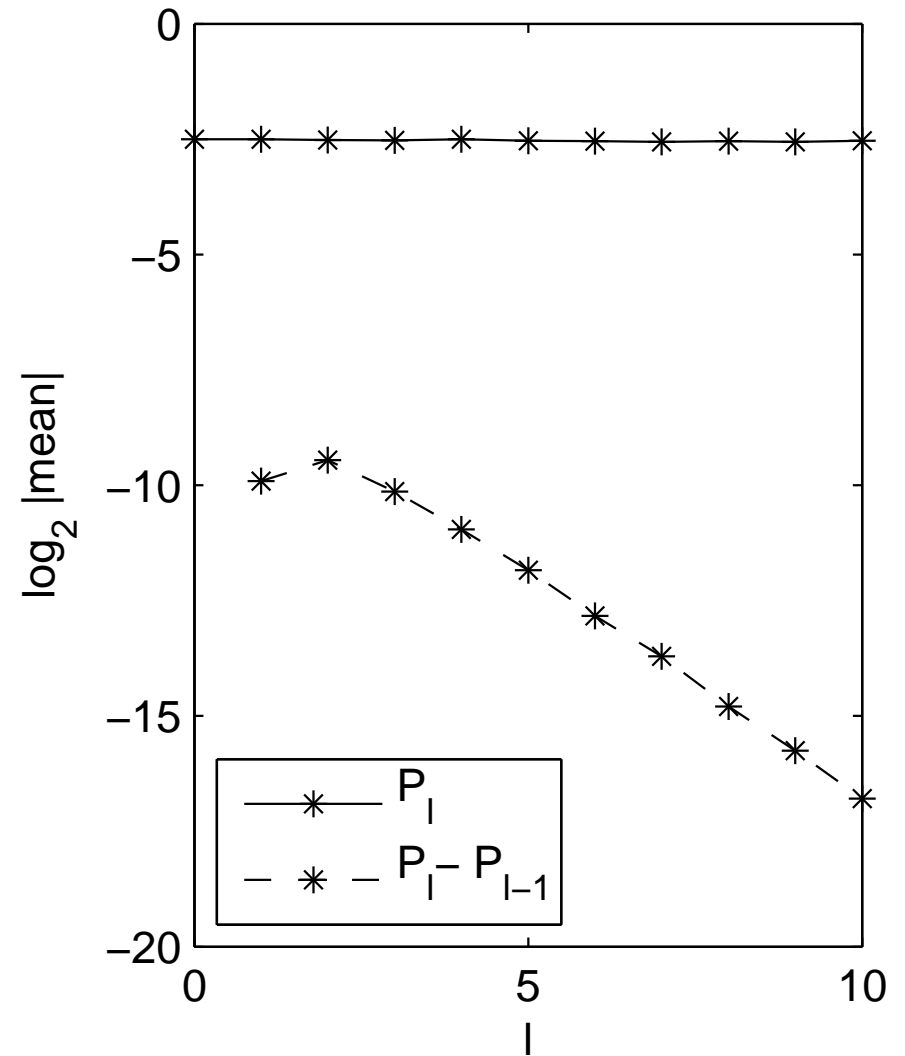
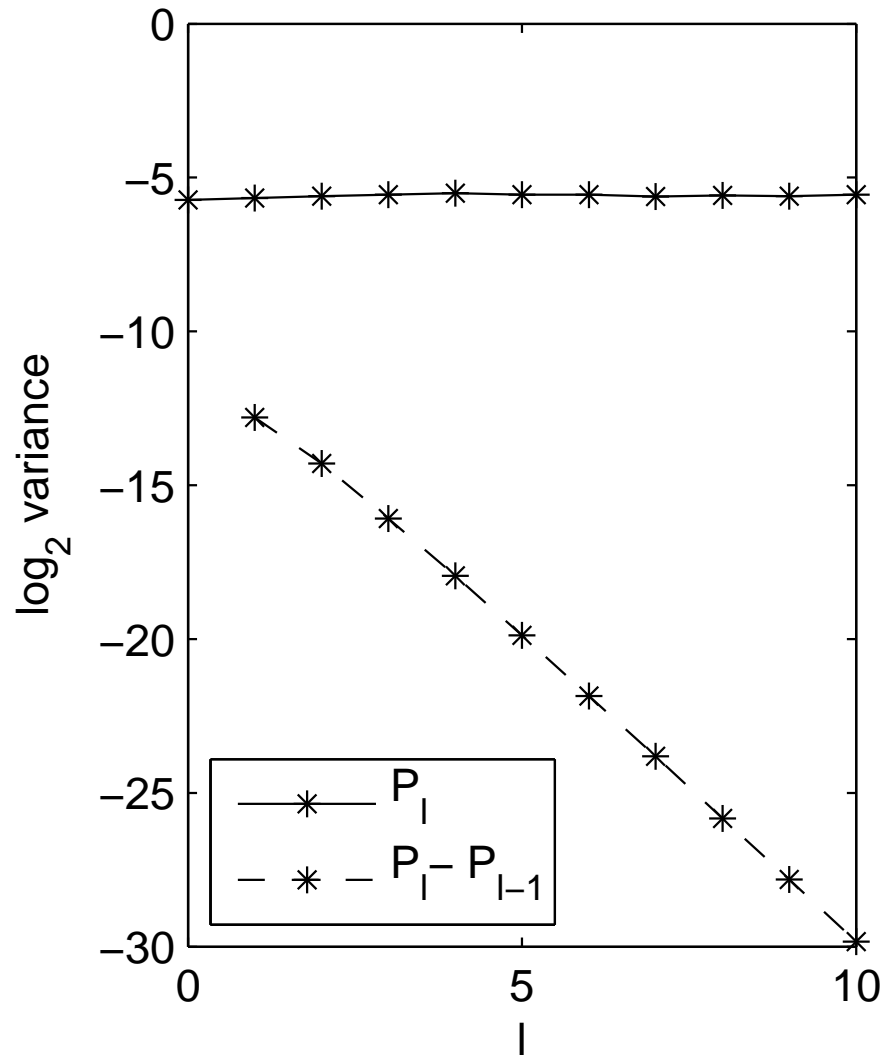
MLMC Results

GBM: Asian option, $\exp(-rT) \max(T^{-1} \int_0^T S(t) dt - 1, 0)$



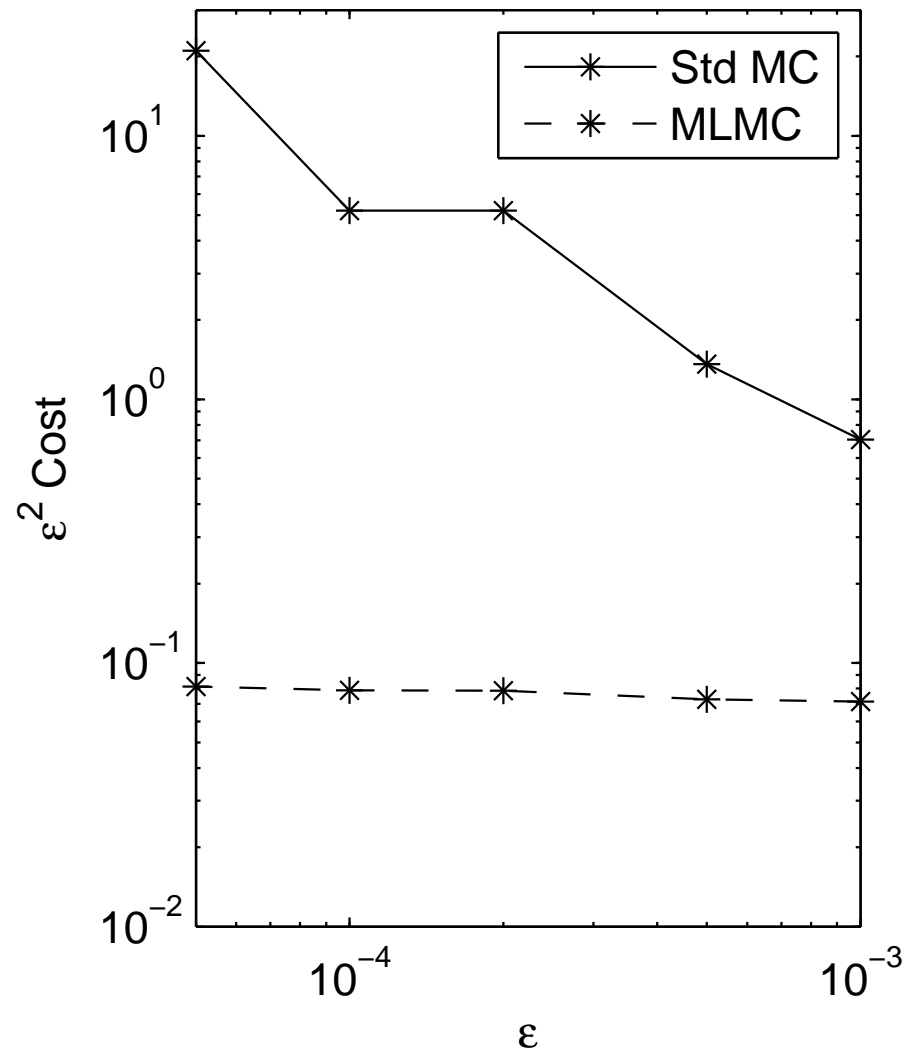
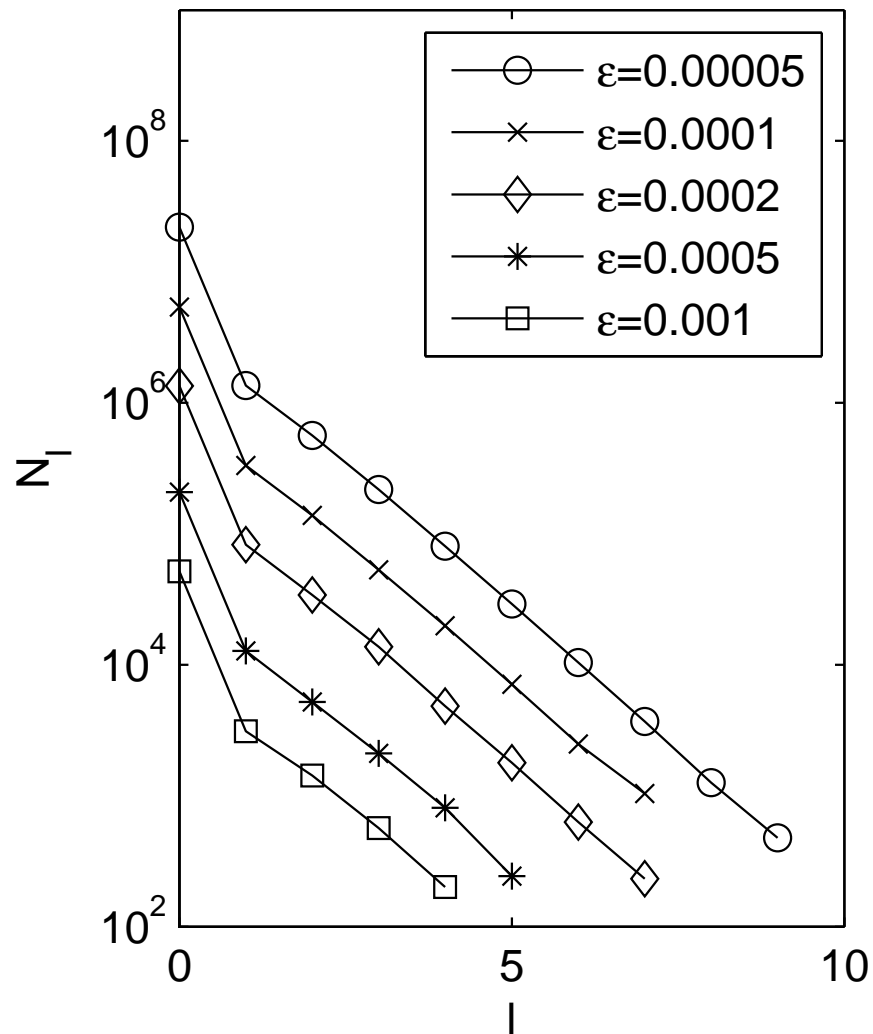
MLMC Results

GBM: lookback option, $\exp(-rT) (S(T) - \min_{0 < t < T} S(t))$



MLMC Results

GBM: lookback option, $\exp(-rT) (S(T) - \min_{0 < t < T} S(t))$



Milstein Scheme

Generic vector SDE:

$$dS(t) = a(S, t) dt + b(S, t) dW(t), \quad 0 < t < T,$$

with correlation matrix $\Omega(S, t)$ between elements of $dW(t)$.

Milstein scheme:

$$\begin{aligned} \widehat{S}_{i,n+1} &= \widehat{S}_{i,n} + a_i h + b_{ij} \Delta W_{j,n} \\ &\quad + \frac{1}{2} \frac{\partial b_{ij}}{\partial S_l} b_{lk} \left(\Delta W_{j,n} \Delta W_{k,n} - h \Omega_{jk} - A_{jk,n} \right) \end{aligned}$$

with implied summation, and Lévy areas defined as

$$A_{jk,n} = \int_{t_n}^{t_{n+1}} (W_j(t) - W_j(t_n)) dW_k - (W_k(t) - W_k(t_n)) dW_j.$$

Milstein Scheme

In vector case:

- $O(h)$ strong convergence if Lévy areas are simulated correctly – expensive
- $O(h^{1/2})$ strong convergence in general if Lévy areas are omitted, except if a certain commutativity condition is satisfied (useful for a number of real cases)
- Lipschitz payoffs can be handled well using antithetic variables
- Other cases may require approximate simulation of Lévy areas

Results

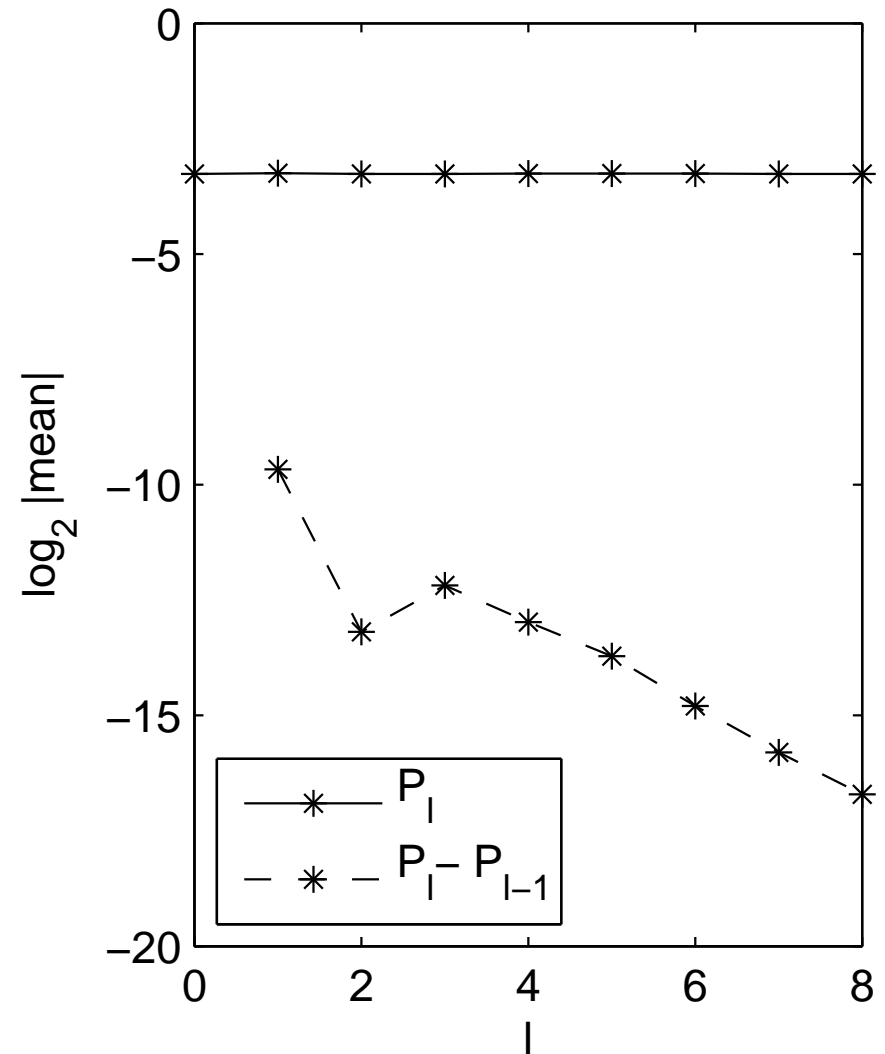
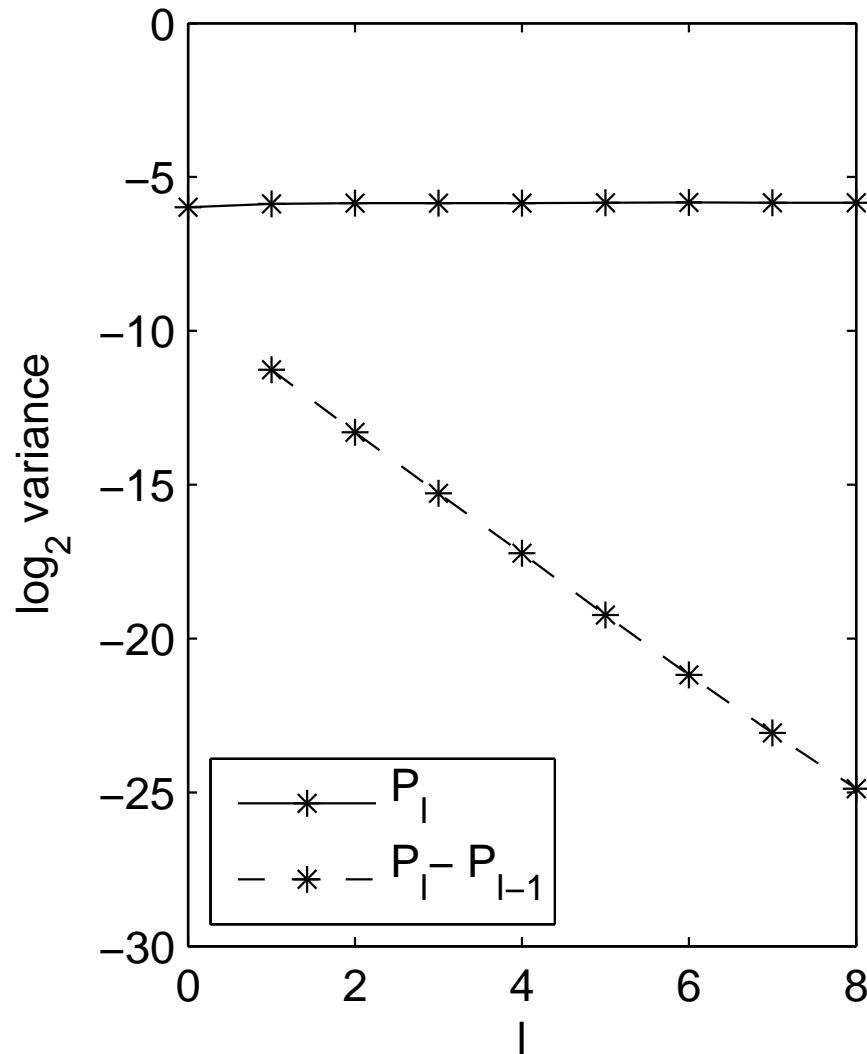
Heston model:

$$\begin{aligned}dS &= r S dt + \sqrt{V} S dW_1, & 0 < t < T \\dV &= \lambda (\sigma^2 - V) dt + \xi \sqrt{V} dW_2,\end{aligned}$$

$$\begin{aligned}T &= 1, & S(0) &= 1, & V(0) &= 0.04, & r &= 0.05, \\ \sigma &= 0.2, & \lambda &= 5, & \xi &= 0.25, & \rho &= -0.5\end{aligned}$$

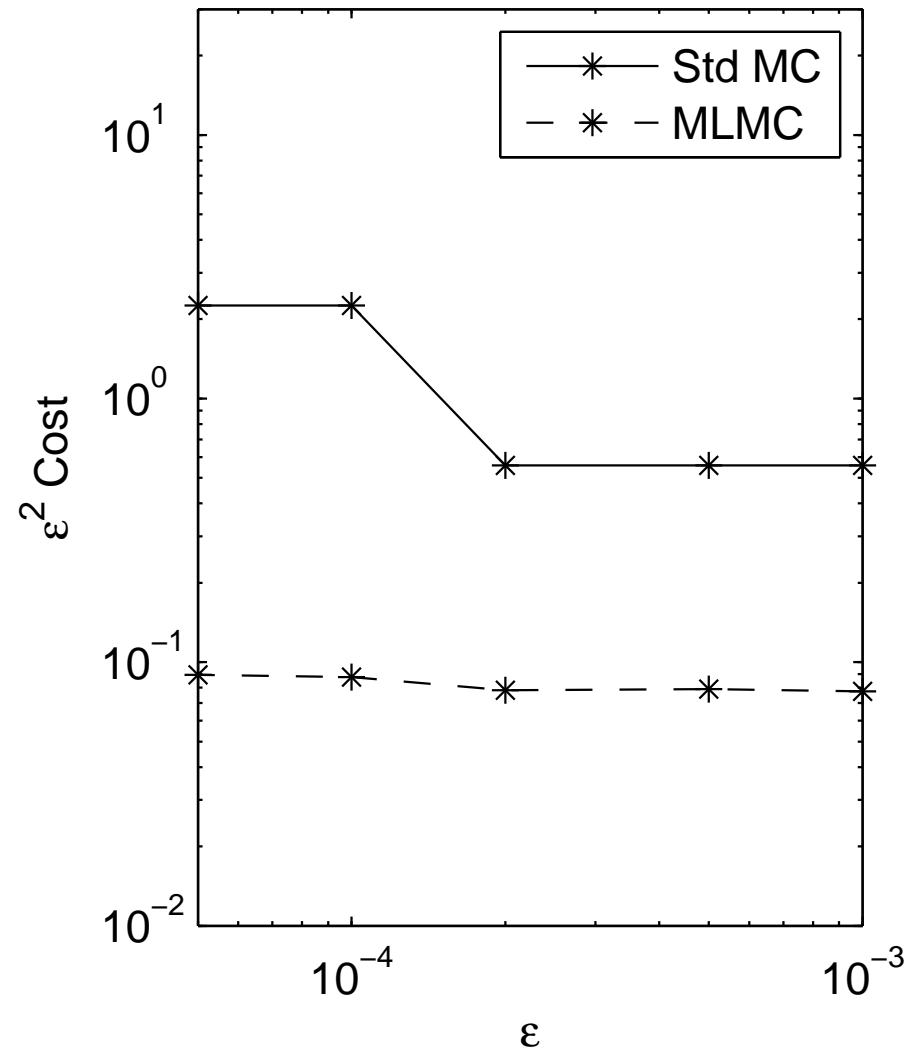
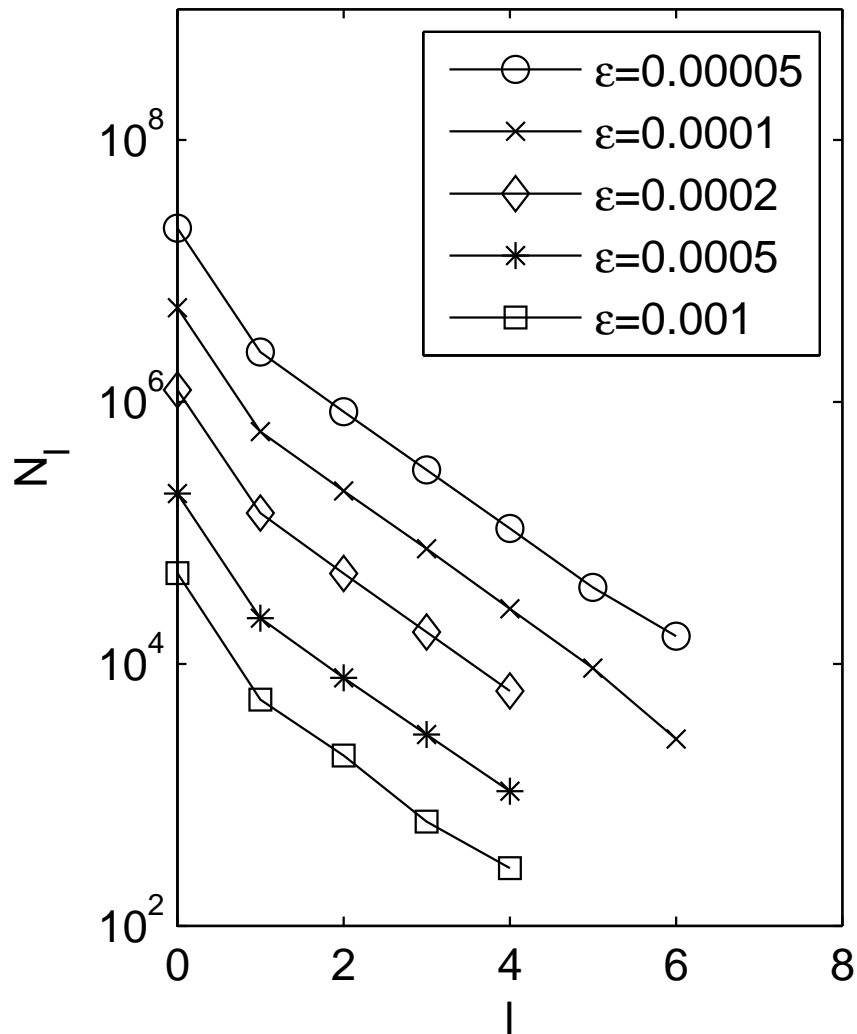
MLMC Results

Heston model: European call



MLMC Results

Heston model: European call



Extensions

1) Quasi-Monte Carlo

- standard Monte Carlo has a random sampling error proportional to $N^{-1/2}$
- Quasi-Monte Carlo uses a deterministic choice of sample “points” to achieve an error which is nearly $O(N^{-1})$ in the best cases
- Not much applicable theory because financial payoffs don't have required smoothness
- In practice, get great results using rank-1 lattice rules developed by Ian Sloan's group at UNSW
- Haven't yet tried Sobol sequences

Extensions

2) Numerical Analysis

option	Euler		Milstein	
	numerics	analysis	numerics	analysis
Lipschitz	$O(h)$	$O(h)$	$O(h^2)$	$O(h^2)$
Asian	$O(h)$	$O(h)$	$O(h^2)$	$O(h^2)$
lookback	$O(h)$	$O(h)$	$O(h^2)$	$o(h^{2-\delta})$
barrier	$O(h^{1/2})$	$o(h^{1/2-\delta})$	$O(h^{3/2})$	$o(h^{3/2-\delta})$
digital	$O(h^{1/2})$	$O(h^{1/2} \log h)$	$O(h^{3/2})$	$o(h^{3/2-\delta})$

Table: V_l convergence observed numerically (for GBM) and proved analytically (for more general SDEs) for both the Euler and Milstein discretisations. δ can be any strictly positive constant.

Extensions

Analysis for Euler discretisations:

- lookback and barrier options: Giles, Higham & Mao (*Finance & Stochastics*, 2009)
 - lookback analysis follows from strong convergence
 - barrier analysis shows dominant contribution comes from paths which are near the barrier; uses asymptotic analysis, first proving that “extreme” paths have negligible contribution
 - similar analysis for digital options gives $O(h^{1/2-\delta})$ bound instead of $O(h^{1/2} \log h)$
- digital options: Avikainen (*Finance & Stochastics*, 2009)
 - method of analysis is quite different

Extensions

Analysis for Milstein discretisations:

- work in progress by Giles, Debrabant & Rößler
- uses boundedness of all moments to bound the contribution to V_l from “extreme” paths
(e.g. for which $\max_n |\Delta W_n| > h^{1/2-\delta}$ for some $\delta > 0$)
- uses asymptotic analysis to bound the contribution from paths which are not “extreme”

Extensions

3) “Greeks”

- this is the name given to derivatives such as $\frac{\partial}{\partial S_0} \mathbb{E}[P]$
- under certain circumstance, this is equal to $\mathbb{E} \left[\frac{\partial P}{\partial S_0} \right]$
 - this leads to the pathwise differentiation approach
- the multilevel approach should again work well but not tried yet
- can also incorporate the adjoint approach developed with Paul Glasserman – more efficient when many Greeks are wanted for one payoff function

Extensions

4) multivariate discontinuous payoffs

- problem with discontinuous payoffs is that small changes in path can lead to a big change in the payoff
- so far, have treated digital options using a “trick” in Paul Glasserman’s book, taking the conditional expectation one timestep before maturity, which effectively smooths the payoff
- the “vibrato” Monte Carlo idea generalises this to multivariate cases in which the conditional expectation is not known in closed form
- alternative is to use “splitting”, with multiple independent simulations of final timestep of each path

Extensions

5) American options

- with European options, the buyer can only exercise the option at maturity, the final time T
- with American options, the buyer can exercise at any time, leading to an optimal control problem
- in PDE approaches, this is solved using a linear complementarity approach which marches backwards in time
- modifying Monte Carlo methods is much harder – an active research topic
- I have some ideas on how to incorporate the multilevel approach – hope to start a project on this soon

Extensions

6) CUDA implementation on NVIDIA graphics cards

- advances in computer hardware/software are important as well as advances in mathematics
- graphics cards are very powerful parallel processors, with up to 240 cores per graphics chip (GPU)
- 2 years ago, NVIDIA introduced the CUDA development environment which uses minor extension to C/C++
- in 2007, a visiting student Xiaoke Su, achieved $100\times$ speedup on a LIBOR application using a 128-core GPU
- more recently, have been working with NAG to develop a random number generation library – freely available to academics

Extensions

Key issue in uniform random number generation:

- when generating 10M random numbers, might have 5000 threads and want each one to compute 2000 random numbers
- need a “skip-ahead” capability so that thread n can jump to the start of its “block” efficiently (usually $\log N$ cost to jump N elements)

Extensions

mrg32k3a (Pierre l'Ecuyer, '99, '02)

- popular generator in Intel MKL and ACML libraries
- pseudo-uniform output is $(x_{n,1} - x_{n,2} \bmod m_1) / m_1$ where integers $x_{n,1}, x_{n,2}$ are defined by

$$x_{n,1} = a_1 x_{n-2,1} - b_1 x_{n-3,1} \bmod m_1$$

$$x_{n,2} = a_2 x_{n-1,2} - b_2 x_{n-3,2} \bmod m_2$$

$$a_1 = 1403580, \quad b_1 = 810728, \quad m_1 = 2^{32} - 209,$$

$$a_2 = 527612, \quad b_2 = 1370589, \quad m_2 = 2^{32} - 22853.$$

Extensions

- Both recurrences are of the form

$$y_n = A y_{n-1} \pmod{m}$$

where y_n is a vector $y_n = (x_n, x_{n-1}, x_{n-2})^T$ and A is a 3×3 matrix. Hence

$$y_{n+2^k} = A^{2^k} y_n \pmod{m} = A_k y_n \pmod{m}$$

where A_k is defined by repeated squaring as

$$A_{k+1} = A_k A_k \pmod{m}, \quad A_0 \equiv A.$$

Can generalise this to jump N places in $O(\log N)$ operations.

Extensions

- output distributions:
 - uniform
 - exponential: trivial
 - Normal: Box-Muller or inverse CDF
 - Gamma: using “rejection” methods which require a varying number of uniforms and Normals to generate 1 Gamma variable
- producing Normals with **mrg32k3a**:
 - 2400M values/sec on a 216-core GTX260
 - 70M values/sec on a Xeon using Intel’s VSL
- have also implemented a **Sobol** generator to produce quasi-random numbers

Conclusions

Multilevel Monte Carlo method has already achieved

- improved order of complexity
- significant benefits for model problems

but more research is needed, both theoretical and applied.

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