A Hypergraph Regularity Method for Generalized Turán Problems*

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Received 14 July 2007; accepted 26 May 2008; received in final form 19 September 2008
Published online in Wiley InterScience (www.interscience.wiley.com).
DOI 10.1002/rsa.20249

ABSTRACT: We describe a method that we believe may be foundational for a comprehensive theory of generalized Turán problems. The cornerstone of our approach is a quasirandom counting lemma for quasirandom hypergraphs, which extends the standard counting lemma by not only counting copies of a particular configuration but also showing that these copies are evenly distributed. We demonstrate the power of the method by proving a conjecture of Mubayi on the codegree threshold of the Fano plane, that is, any 3-graph on \( n \) vertices for which every pair of vertices is contained in more than \( n/2 \) edges must contain a Fano plane, for \( n \) sufficiently large. For projective planes over fields of odd size \( q \) we show that the codegree threshold is between \( n/2 - q + 1 \) and \( n/2 \), but for \( PG_2(4) \) we find the somewhat surprising phenomenon that the threshold is less than \( (1/2 - \epsilon)n \) for some small \( \epsilon > 0 \). We conclude by setting out a program for future developments of this method to tackle other problems.

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Keywords: hypergraphs; extremal combinatorics; hypergraph regularity

1. INTRODUCTION

A famous unsolved question of Turán asks for the maximum size of a 3-graph\(^1\) on \( n \) vertices that does not contain a tetrahedron \( K_4^3 \), i.e. four vertices on which every triple is present. Despite the superficial similarity to the analogous easy question for graphs (the maximum size of a graph with no triangle) this problem has evaded even an asymptotic solution for over

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*Supported by NSF (DMS-0555755).

\(^1\)A \( k \)-graph \( G \) consists of a vertex set \( V(G) \) and an edge set \( E(G) \), each edge being some \( k \)-tuple of vertices.

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60 years. It may be considered a test case for the general Turán problem, that of determining
the maximum size of a $k$-graph on $n$ vertices that does not contain some fixed $k$-graph $F$.
This maximum size is called the Turán number of $F$, denoted $\text{ex}(n, F)$. It is not hard
to show that the limit $\pi(F) = \lim_{n \to \infty} \text{ex}(n, F)/\binom{n}{k}$ exists. As a first step to understanding the
Turán number we may ask to determine this limit, the Turán density, which describes the
asymptotic behavior of the Turán number. There are very few known results even in this
weaker form, and no general principles have been developed, even conjecturally.

For $0 \leq s \leq k$ we may define a generalized Turán number $\text{ex}_s(n, F)$ as the largest
number $m$ such that there is a $k$-graph $H$ on $n$ vertices that does not contain $F$ and has
minimum $s$-degree $\delta_s(H) \geq m$. Note that we recover $\text{ex}(n, F)$ in the case $s = 0$, and the case
$s = k$ is trivial. The cases $s = 0$ and $s = 1$ are essentially equivalent, via a well-
known induction argument, and so there is no new theory here for graphs. However, for
general hypergraphs we obtain a rich source of new problems, and it is not apparent how they relate to each other. There has been much recent interest in the case $s = k - 1$, which
were called codegree problems in [20]. (See also [17, 22] for similar questions involving
structures that are spanning rather than fixed.) We may define generalized Turán densities
as $\pi_s(F) = \lim_{n \to \infty} \text{ex}_s(n, F)/\binom{n}{k}$. A simple averaging argument shows that if we define
the normalized minimum $s$-degrees of $H$ as $\hat{\delta}_s(H) = \delta_s(H)/\binom{n}{k}$ then we have a hierarchy
$\hat{\delta}_0(H) \geq \hat{\delta}_1(H) \geq \cdots \geq \hat{\delta}_{k-1}(H)$, so $\pi_{k-1}(F) \leq \cdots \leq \pi_0(F)$.

Projective geometries $\text{PG}_m(q)$ provide examples of configurations $F$ that are surprisingly
tractable for these problems. For the Fano plane ($m = q = 2$) the exact Turán number for
$n$ sufficiently large was determined independently and simultaneously by Keevash and
Sudakov [13] and Füredi and Simonovits [6]. They also characterized the unique maximizing
configuration: a balanced complete bipartite $3$-graph. Earlier, de Caen and Füredi [4] had
obtained the Turán density $\pi(\text{PG}_2(2)) = 3/4$. On the other hand, Mubayi [19] showed that the
codegree density of the Fano plane is $\pi_2(\text{PG}_2(2)) = 1/2$. He conjectured that the exact
codegree threshold satisfies $\text{ex}_2(n, \text{PG}_2(2)) \leq n/2$. The following result establishes this and
characterizes the case of equality.

Theorem 1.1. If $n$ is sufficiently large and $H$ is a $3$-graph on $n$ vertices with minimum
2-degree at least $n/2$ that does not contain a Fano plane then $n$ is even and $H$ is a balanced
complete bipartite $3$-graph.

General projective geometries have been studied in [11] (the Turán problem) and [15]
(the codegree problem). A general bound $\pi_q(\text{PG}_m(q)) \leq 1 - 1/m$ was obtained in [15], and
it was shown that equality holds whenever $m = 2$ and $q$ is $2$ or odd, and whenever $m = 3$
and $q$ is $2$ or $3$. We prove the following results, which give quite precise information about

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2Given $S \subseteq V(H)$ the neighborhood of $S$ in $H$ is $N_H(S) = \{T \subseteq V(H) \setminus S : S \cup T \in E(H)\}$ and the degree of $S$ is $|N_H(S)|$. The minimum $s$-degree $\delta_s(H) = \min_{|S|=s} |N_H(S)|$ is the minimum of $|N_H(S)|$ over all subsets $S$ of size $s$.

3In the codegree case $s = k - 1$ this limit was shown to exist in [20]. In general the existence may be deduced from a very general theory of Razborov [21]. This is perhaps using a sledgehammer to crack a nut, and in fact
the method of [20] can be extended using martingale estimates in place of hypergeometric estimates. We will elaborate slightly on the martingale aspect in the final section, but a detailed treatment is beyond the scope we have set for this article.

4A $k$-graph $H$ is bipartite if there is a partition $V(H) = A \cup B$ so that there are no edges of $H$ lying entirely within $A$ or entirely within $B$. A complete bipartite $k$-graph contains all edges that intersect both $A$ and $B$. It is balanced if $|A| - |B| \leq 1$.

Random Structures and Algorithms DOI 10.1002/rsa
the codegree threshold for planes over a field of odd size and demonstrate a surprisingly different behavior for $PG_2(4)$.

**Theorem 1.2.** Suppose $q$ is an odd prime power. Then $\lfloor n/2 \rfloor - q + 1 \leq ex_q(n, PG_2(q)) \leq n/2$. In the case $q = 3$ and $n$ even we have $ex_3(n, PG_2(3)) = n/2 - 1$.

**Theorem 1.3.** There is $\epsilon > 0$ for which $\pi_4(PG_2(4)) < 1/2 - \epsilon$.

The main idea in our arguments is a quasirandom counting lemma that extends the (usual) counting lemma for quasirandom hypergraphs. We adopt the Gowers framework as being most compatible with our argument (there are other approaches to this theory, see Rödl et al. (e.g. [24, 25]) and Tao [29]). We will give precise definitions later, and for now describe our result on an intuitive level. Hypergraph regularity theory gives a method of decomposing a hypergraph into a bounded number of pieces, each of which behaves in an approximately random fashion. The number of pieces depends only on the degree of approximation and is independent of the size of the hypergraph. In order for such a decomposition to be useful, the notion of random behavior should be sufficiently powerful for applications, and the general criterion that has been used is that there should be a counting lemma, meaning a result that a sufficiently random hypergraph contains many copies of any small fixed configuration. Our quasirandom counting lemma will state that not only are there many copies, but that they are uniformly distributed within the hypergraph.

We also make use of the idea of stability, a phenomenon which was originally discovered by Erdős and Simonovits in the 60s in the context of graphs with excluded subgraphs, but has only been systematically explored relatively recently, as researchers have realized the importance and applications of such results in hypergraph Turán theory [13, 14], enumeration of discrete structures and extremal set theory (see [12] as a recent example and for many further references).

The rest of this article is organized as follows. The next section is expository in nature: it introduces the theory needed in later sections for the special case of graphs, where it will be mostly familiar to many readers (although our quasirandom counting lemma is new even for graphs). Then in Section 3 we introduce the Gowers quasirandomness framework for 3-graphs and present a case of our quasirandom counting lemma that we will need to prove Theorem 1.1. Section 4 contains the proof of Theorem 1.1, using the quasirandomness theory from Section 3 and also the method of “stability,” or approximate structure. In Section 5 we present the general theory of quasirandomness hypergraphs and the full form of our quasirandom counting lemma: this is the engine behind our entire approach. We also give an application to generalized Turán problem for configurations that have a certain special form. This general theory is applied in Section 6 to the study of codegree problems in projective planes, where we prove the other theorems stated above. The final section sets out a program for future developments of this method to other generalized Turán problems. Since our formulation of the Gowers quasirandomness framework uses some non-trivial variations on the original framework, we give justifications for these variations in the Appendix of this article.

**Notation.** Write $[n] = \{1, \ldots, n\}$. If $X$ is a set and $k$ is a number then $\binom{X}{k} = \{Y \subseteq X : |Y| = k\}$, $\binom{X}{<k} = \cup_{i<k} \binom{X}{i}$ and $\binom{X}{\leq k} = \cup_{i\leq k} \binom{X}{i}$. $a \pm b$ denotes an unspecified real number in the interval $[a - b, a + b]$. It is convenient to regard a finite set $X$ as being equipped with the uniform probability measure $P(\{x\}) = 1/|X|$, so that we can express the average.
of a function \( f \) defined on \( X \) as \( \mathbb{E}_{x \in X} f(x) \). A \( k \)-graph \( H \) consists of a vertex set \( V(H) \) and an edge set \( E(H) \), each edge being some \( k \)-tuple of vertices. We often identify \( H \) with \( E(H) \), thus \( |H| \) is the number of edges in \( H \). Given \( S \subseteq V(H) \) the neighborhood of \( S \) in \( H \) is \( N_H(S) = \{ T \subseteq V(H) \setminus S : S \cup T \in E(H) \} \) and the degree of \( S \) is \( |N_H(S)| \). The minimum \( s \)-degree \( \delta_s(H) = \min_{|S|=s} |N_H(S)| \) is the minimum of \( |N_H(S)| \) over all subsets \( S \) of size \( s \). Given \( X \subseteq V(H) \) the restriction \( H[X] \) is a \( k \)-graph with vertex set \( X \) and edge set equal to all those edges of \( H \) that are contained in \( X \). Suppose \( F \) and \( H \) are \( k \)-graphs. The homomorphism density \( d_F(H) \) is the probability that a randomly chosen map \( \phi : V(F) \to V(H) \) is a homomorphism, i.e. \( \phi(e) \) is an edge of \( H \) for every edge \( e \) of \( F \) (which we also write as \( \phi(F) \subseteq H \)). We also use the same notation in a “partite setting” (this will be explained when it occurs). When \( F = e \) consists of just a single edge we write \( d(H) = d_e(H) = k!|E(H)||V(H)|^{-k} \), and call this the density of \( H \). We use the notation \( 0 < \alpha \ll \beta \) to mean that there is an increasing function \( f(x) \) so that the ensuing argument is valid for \( 0 < \alpha < f(\beta) \).

2. GRAPHS: REGULARITY AND COUNTING, QUASIRANDOMNESS, AND QUASIRANDOM COUNTING

The purpose of this section is expository: we introduce the theory needed in later sections for the special case of graphs, where it is considerably simpler, and partly familiar to many readers.

In the first subsection we describe Szemerédi’s regularity lemma [27], one of the most powerful tools in modern graph theory. Roughly speaking, it says that any graph can be approximated by an average with respect to a partition of its vertex set into a bounded number of classes, the number of classes depending only on the accuracy of the desired approximation, and not on the number of vertices in the graph. Each pair of classes span a bipartite subgraph that is “regular,” meaning that the proportion of edges in any large bipartite subgraph is close to the proportion of edges in the pair as a whole. A key property of this approximation is that it leads to a “counting lemma,” allowing an accurate prediction of the number of copies of any small fixed graph spanned by some specified classes of the partition. We refer the reader to [16] for a survey of the regularity lemma and its applications.

The second subsection discusses quasirandomness of graphs, a concept introduced by Chung, Graham and Wilson [3] (see also Thomason [31] for a similar notion). There are many ways of describing this concept, all of which are broadly equivalent (up to renaming constants); in fact, it is also equivalent to regularity (as described in the first subsection). A particularly simple formulation is to call a bipartite graph quasirandom if the number of 4-cycles is close to what would be expected in a random graph with the same edge density. A closely related formulation that forms the basis for the Gowers approach to quasirandomness in hypergraphs is to say that if we count 4-cycles weighted by the “balanced function” of the graph then the result is small. Our discussion in this subsection is based on Section 3 of [7] (we are more brief on those points discussed there, but we also provide some additional arguments that are omitted there).

In the third subsection we introduce the graph case of our quasirandom counting lemma, an extension of the counting lemma discussed in the first subsection, saying that copies of any small fixed graph are well-distributed in the graph. This is a new result even in the special case of graphs, and has consequences that are somewhat surprising at first sight.
2.1. Regularity and Counting

We start by describing the notion of “regularity” for bipartite graphs. The density of a bipartite graph $G = (A, B)$ with vertex classes $A$ and $B$ is defined to be

$$d_G(A, B) := \frac{e_G(A, B)}{|A||B|}.$$ 

We often write $d(A, B)$ if this is unambiguous. Given $\epsilon > 0$, we say that $G$ is $\epsilon$-regular if for all subsets $X \subseteq A$ and $Y \subseteq B$ with $|X| > \epsilon|A|$ and $|Y| > \epsilon|B|$ we have that $|d(X, Y) - d(A, B)| < \epsilon$.

The regularity lemma says that any graph can be partitioned into a bounded number of regular pairs and a few leftover edges. Formally:

**Theorem 2.1.** For every real $\epsilon > 0$ and number $m_0 \geq 1$ there are numbers $m, n_0 \geq 1$ so that for any graph $G$ on $n \geq n_0$ vertices we can partition its vertices as $V(G) = V_0 \cup V_1 \cup \cdots \cup V_k$ so that

- $m_0 \leq k \leq m$,
- $|V_0| < \epsilon n$,
- $|V_1| = |V_2| = \cdots = |V_k|$, and
- $(V_i, V_j)$ spans an $\epsilon$-regular bipartite subgraph of $G$ for all but at most $\epsilon k^2$ pairs $1 \leq i < j \leq k$.

**Remarks.** In applications one takes $k \geq m_0 \geq \epsilon^{-1}$ so that the number of edges within any $V_i$ is negligible. An “exceptional class” $V_0$ is allowed so that the remaining partition can be “equitable,” i.e. $|V_1| = |V_2| = \cdots = |V_k|$. If one prefers not to have an exceptional class then its vertices may be distributed among the other classes to obtain a partition with the same regularity properties for a slightly larger $\epsilon$ with the class sizes differing by at most one.

We refer the reader to Section 1 of [16] for further discussion of variants of the regularity lemma.

A key property of regularity is that we can accurately count copies of any small fixed graph. For the purpose of exposition we state two simple cases, and then the general “counting lemma.” Throughout we assume a hierarchy $0 < 1/n \ll \epsilon \ll d$, that every density we consider is at least $d$ and each part in our graphs contains at least $n$ vertices. The following well-known statements can be proved using similar arguments to that given for Lemma 2.1 in [16].

**Triangles.** Suppose $G$ is a tripartite graph with parts $V_1$, $V_2$, $V_3$ and each pair $(V_i, V_j)$ spans an $\epsilon$-regular bipartite graph of density $d_{ij}$. Let $\Delta(G)$ be the set of triangles in $G$.

Then we can estimate the “triangle density” in $G$ as $d_{\Delta}(G) = |\Delta(G)|/|V_1||V_2||V_3| = d_{12}d_{13}d_{23} \pm 8\epsilon$.

**Four-cycles.** Suppose $G = (X, Y)$ is an $\epsilon$-regular bipartite graph with density $d$. Let $C_4(G)$ be the number of labelled 4-cycles in $G$, i.e. quadruples $(x_1, x_2, y_1, y_2)$ with $x_1 \neq x_2 \in X$, $y_1 \neq y_2 \in Y$ such that $x_1y_1$, $x_1y_2$, $x_2y_1$ and $x_2y_2$ are all edges of $G$. We may define

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5The constant 8 is not best possible, but we only care that the error should tend to zero as $\epsilon$ tends to 0.
Suppose $G$ is an $r$-partite graph on $V = V_1 \cup \cdots \cup V_r$ and each pair $(V_i, V_j), 1 \leq i < j \leq r$ spans an $\epsilon$-regular bipartite graph with density $d_{ij}(G)$. Suppose $H$ is an $r$-partite graph on $Y = Y_1 \cup \cdots \cup Y_r$. Let $\Phi(Y, V)$ be the set of $r$-partite maps from $Y$ to $V$, i.e. maps $\phi : Y \to V$ with $\phi(Y_i) \subseteq V_i$ for $1 \leq i \leq r$. Define the $r$-partite homomorphism density of $H$ in $G$ as $d_H(G) = \mathbb{P}_{\phi \in \Phi(Y, V)}[\phi(H) \subseteq G]$. Then $d_H(G) = \prod_{e \in E(H)} d_e(G) + O_H(\epsilon)$, where $d_e(G)$ means that density $d_{ij}(G)$ for which $e \in H[Y_i, Y_j]$ and the $O_H$-notation indicates that the implied constant depends only on $H$.

### 2.2. Quasirandomness

The regularity property discussed in the previous subsection turns out to be characterized by the counting lemma; in fact, somewhat surprisingly, it is characterized just by counting 4-cycles. To be precise, if $0 < 1/n \ll \epsilon \ll \epsilon' \ll d \leq 1$ and $G = (X, Y)$ is a bipartite graph with $|X|, |Y| \geq n$, density $d$ and $C_4$-density $d_{C_4}(G) < d^4 + \epsilon$ then $G$ is $\epsilon'$-regular.

To illuminate some later more general arguments we will prove this fact here, via a closely related characterization of Gowers (see Section 3 of [7]) that forms the basis for the Gowers approach to quasirandomness in hypergraphs.

Suppose $G$ is a bipartite graph with parts $X$ and $Y$. We can identify $G$ with its characteristic function $G(x, y)$, which for $x \in X$ and $y \in Y$ is defined to be 1 if $xy$ is an edge of $G$, or 0 otherwise. Define the balanced function $\overline{G}(x, y) = G(x, y) - d$, where $d = d_G(X, Y) = |E_G(X, Y)|/|X||Y|$ is the density of $G$, i.e. $\overline{G}(x, y)$ is $1 - d$ if $xy$ is an edge of $G$, or $-d$ otherwise. Note that $\sum_{x \in X, y \in Y} \overline{G}(x, y) = 0$. For any function $f$ defined on $X \times Y$ we define

$$C_4(f) = |X|^2|Y|^2 \sum_{x_1, x_2 \in X} \sum_{y_1, y_2 \in Y} f(x_1, y_1) f(x_1, y_2) f(x_2, y_1) f(x_2, y_2).$$

We can rephrase this by recalling the setup in the previous subsection, where we had a 4-cycle $C_4 = (A, B)$, considered as a bipartite graph $K_{2,2}$ on $A = \{a_1, a_2\}$ and $B = \{b_1, b_2\}$, and let $\Phi$ denote the set of all bipartite maps $\phi$ from $A \cup B$ to $X \cup Y$. Then

$$C_4(f) = \mathbb{E}_{\phi \in \Phi} \prod_{e \in E(C_4)} f(\phi(e)).$$

In particular $C_4(G) = d_{C_4}(G)$, where the first instance of $G$ is to be understood as the characteristic function $G(x, y)$.

We say that a function $f : X \times Y \to [-1, 1]$ is $\eta$-quasirandom if $C_4(f) < \eta$, and we say that $G$ is $\eta$-quasirandom if its balanced function $\overline{G}$ is $\eta$-quasirandom.

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6By $\phi(C_4) \subseteq G$ we mean $\phi(e) \in E(G)$ for each edge $e$ of $C_4$. 

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Theorem 2.2. Suppose $G$ is a bipartite graph of density $d$ with parts $X$ and $Y$ of size at least $n$. The following are equivalent, in the sense that each implication $(i) \Rightarrow (j)$ is true for $0 < 1/n \ll \epsilon_i \ll \epsilon_j \ll d \leq 1$:

1. $G$ is $\epsilon_1$-regular.
2. $d_{C_4}(G) = d^4 \pm \epsilon_2$.
3. $C_4(G) < \epsilon_3$.

Before giving the proof we quote a simple version of the “second moment method,” Lemma 6.5 in [7]: if $0 < \alpha, d < 1$, $\sum_{i=1}^{n} a_i \geq (d - \alpha)n$ and $\sum_{i=1}^{n} a_i^2 \leq (d^2 + \alpha)n$ then $a_i = d \pm \alpha^{1/2}$ for all but at most $3\alpha^{1/2}$ values of $i$. Indeed, $\sum_{i=1}^{n} (a_i - d)^2 = \sum a_i^2 - 2d \sum a_i + d^2 n \leq (d^2 + \alpha - 2d(d-\alpha) + d^2)n < 3an$. Also, we often use the Cauchy-Schwartz inequality in the form $0 \leq \mathbb{E}(Z - \mathbb{E}Z)^2 = \mathbb{E}Z^2 - (\mathbb{E}Z)^2$, for a random variable $Z$.

Proof. (1) $\Rightarrow$ (2): This implication is given by the counting lemma quoted in the previous subsection, but for completeness we give a proof here. Consider a random map $\phi \in \Phi$, i.e. a random bipartite map from $A \cup B$ to $X \cup Y$, where as before we consider $C_4$ as a bipartite graph with parts $A = \{a_1, a_2\}$ and $B = \{b_1, b_2\}$. Let $E_1$ be the event that $x_1 = \phi(a_1)$ has $(d \pm \epsilon_1)|Y|$ neighbors in $Y$. By definition of $\epsilon_1$-regularity we have $\mathbb{P}(E_1) > 1 - 2\epsilon_1$ (there are at most $\epsilon_1|X|$ vertices with more than $(d + \epsilon_1)|Y|$ neighbors and at most $\epsilon_1|X|$ vertices with less than $(d - \epsilon_1)|Y|$ neighbors). Let $E_2$ be the event that $x_2 = \phi(a_2)$ has $(d \pm \epsilon_1)|Y|$ neighbors in $N(x_1)$. Again, $\epsilon_1$-regularity gives $\mathbb{P}(E_2|E_1) > 1 - 2\epsilon_1$. Now the event $\{\phi(C_4) \subseteq G\}$ occurs if and only if $y_1 = \phi(b_1)$ and $y_2 = \phi(b_2)$ lie in $N(x_1) \cap N(x_2)$, so we have

$$d_{C_4}(G) = \mathbb{P}_{\phi \in \Phi}[\phi(C_4) \subseteq G] = \mathbb{P}_{\phi \in \Phi}[[\phi(C_4) \subseteq G] \cap E_1 \cap E_2] \pm 4\epsilon_1 = d^4 \pm 10\epsilon_1.$$

This proves the implication with $\epsilon_2 = 10\epsilon_1$.

(2) $\Rightarrow$ (3): Suppose that $d_{C_4}(G) = d^4 \pm \epsilon_2$. Our first step is to show that we can also count any subgraph of $C_4$, in that $d_{P_2}(G) = d^2 \pm \epsilon_2$ and $d_{P_3}(G) = d^3 \pm 4\epsilon_2^{1/4}$, where $P_1$ is the path with $i$ edges (we already know that $d_{P_1}(G) = d|X - Y| = d$, and it is immediate that for a matching $M_2$ of two edges we have $d_{M_2}(G) = d^2$). We start with the homomorphism density of $P_2$, say with the central vertex being mapped to $X$ and the two outer vertices to $Y$ (the same bound will hold vice versa). We can write

$$d_{P_2}(G) = \mathbb{E}_{x \in X, y_1 \in Y, y_2 \in Y} G(x, y_1)G(x, y_2) = \mathbb{E}_{x \in X} (\mathbb{E}_{y \in Y} G(x, y))^2,$$

which by Cauchy-Schwartz is at least $(\mathbb{E}_{x \in X} \mathbb{E}_{y \in Y} G(x, y))^2 = d^2$. On the other hand, we can again apply Cauchy-Schwartz to get

$$d_{P_2}(G)^2 = (\mathbb{E}_{x_1, y_1} G(x, y_1)G(x, y_2))^2 \leq \mathbb{E}_{x_1, y_2} (\mathbb{E}_x G(x, y_1)G(x, y_2))^2 = \mathbb{E}_{x_1, y_2} \mathbb{E}_{x_2, y_2} G(x_1, y_1)G(x_1, y_2)G(x_2, y_1)G(x_2, y_2) = d_{C_4}(G) < d^4 + \epsilon_2.$$

This gives $d_{P_2}(G) = d^2 \pm \epsilon_2$, and moreover, applying the second moment method quoted before the proof, for a random $x \in X$, with probability at least $1 - 3\epsilon_2^{1/2}$ we have $\mathbb{E}_{x_1, x_2} G(x, y_1)G(x, y_2) = d^2 \pm \epsilon_2^{1/4}$, so $|N(x)|/|Y| = \mathbb{E}_x G(x, y) = d \pm \epsilon_2^{1/4}$. This allows us to estimate

$$d_{P_3}(G) = \mathbb{E}_{x_1, x_2} \mathbb{E}_{y_1, y_2} G(x_1, y_1)G(x_1, y_2)G(x_2, y_2) = (d^2 \pm \epsilon_2^{1/4})(d \pm \epsilon_2^{1/4}) \pm 6\epsilon_2^{1/2} = d^3 \pm 4\epsilon_2^{1/4},$$

Random Structures and Algorithms DOI 10.1002/rsa
where the main term gives the contribution when $x_1$ and $y_2$ have typical neighborhoods and $6\varepsilon_2^{1/2}$ bounds the error coming from atypical $x_1$ and $y_2$.

Now we can estimate $C_4(G)$. Write $\overline{G} = f_0 - f_1$, where $f_0(x, y) = G(x, y)$ and $f_1(x, y) = d$ (a constant function). Then

$$C_4(\overline{G}) = \mathbb{E}_{x_1, y_1, y_2 \in \overline{G}} \overline{G}(x_1, y_1) \overline{G}(x_1, y_2) \overline{G}(x_2, y_1) \overline{G}(x_2, y_2)$$

$$= \sum_{M} \sum_{-\mathbb{E}_{x_1, y_1, y_2 \in \overline{G}} M_{ij}^1(x_1, y_1) M_{ij}^2(x_1, y_2) M_{ij}^2(x_2, y_1) M_{ij}^2(x_2, y_2),$$

where $M = (M_{ij})_{i,j \in [1, 2]}$ ranges over $2 \times 2$ matrices with $\{0, 1\}$-entries and $\sum M$ is $\sum_{i,j} M_{ij}$. For each $M$ we can estimate the summand corresponding to $M$, using the estimate for the homomorphism density of the subgraph of $C_4$ corresponding to the $0$-entries of $M$. For example, if $M_{11} = M_{12} = M_{21} = 0$ and $M_{22} = 1$, then $M$ corresponds to a $P_3$, and the corresponding summand is $-d \cdot d_{P_3}(G) = -d^4 + 4\varepsilon_2^{1/4}$. Thus we obtain $8$ summands in the range $d^4 + 4\varepsilon_2^{1/4}$ and $8$ summands in the range $-d^4 + 4\varepsilon_2^{1/4}$, giving a total of at most $64\varepsilon_2^{1/4}$.

This proves the implication with $\varepsilon_3 = 64\varepsilon_2^{1/4}$.

(3) $\Rightarrow$ (1): Suppose $C_4(\overline{G}) < \varepsilon_3$. Consider $X' \subseteq X$ with $|X'| > \varepsilon_1 |X|$ and $Y' \subseteq Y$ with $|Y'| > \varepsilon_1 |Y|$. Write $X'(x)$ for the characteristic function of $X'$, i.e. $X'(x) = 1$ if $x \in X'$, otherwise it is 0; similarly, let $Y'(y)$ be the characteristic function of $Y'$. Then

$$d_G(X', Y') = d = \mathbb{E}_{x \in X', y \in Y'} \overline{G}(x, y) = \frac{|X||Y|}{|X'||Y'|} \mathbb{E}_{x \in X', y \in Y'} X'(x) Y'(y) \overline{G}(x, y)$$

and by Cauchy-Schwartz (in the first and third inequalities below) we have

$$(\mathbb{E}_{x \in X, y \in Y} X'(x) Y'(y) \overline{G}(x, y))^4 \leq (\mathbb{E}_{x \in X} (\mathbb{E}_{y \in Y} X'(x) Y'(y) \overline{G}(x, y))^2)^2$$

$$\leq (\mathbb{E}_{x \in X} (\mathbb{E}_{y \in Y} Y'(y) \overline{G}(x, y))^2)^2$$

$$= (\mathbb{E}_{x \in X} Y_{y_1, y_2 \in Y} Y'(y_1) Y'(y_2) \overline{G}(x, y_1) \overline{G}(x, y_2))^2$$

$$\leq \mathbb{E}_{y_1, y_2 \in Y} Y'(y_1) Y'(y_2) \overline{G}(x, y_1) \overline{G}(x, y_2))^2$$

$$\leq (\mathbb{E}_{x_1, x_2 \in X} (\mathbb{E}_{x \in X} \overline{G}(x, y_1) \overline{G}(x_1, y_1) \overline{G}(x, y_2))^2$$

$$= C_4(\overline{G}) < \varepsilon_3.$$

This gives $d_G(X', Y') = d = \varepsilon_1^{-2} \varepsilon_3^{1/4}$, which proves the implication with $\varepsilon_1 = \varepsilon_3^{1/12}$.

### 2.3. Quasirandom Counting

Now we will introduce the graph case of our quasirandom counting lemma, an extension of the counting lemma discussed in the first subsection, saying that copies of any small fixed graph are well-distributed in the graph. As before, for the purpose of exposition we lead up to the general case through two illustrative cases.

**Triangles.** Suppose $G$ is a tripartite graph with parts $V_1$, $V_2$, $V_3$ and each pair $(V_i, V_j)$ spans an $\varepsilon$-regular bipartite graph of density $d_{ij}$. We remarked before that the “triangle density” in $G$ can be estimated as $d_d(G) = d_{12}d_{13}d_{23} \pm 8\varepsilon$. Moreover, it is easy to see that the triangles of $G$ are “well-distributed,” in that any sufficiently large subsets

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$V'_i \subseteq V_i, 1 \leq i \leq 3$ induce a subgraph of $G$ that also has triangle density about $d_{12}d_{13}d_{23}$. Indeed, if $V'_i \subseteq V_i$ with $|V'_i| > \epsilon^{1/2}|V_i|, 1 \leq i \leq 3$, then each $(V'_i, V'_j)$ induces an $\epsilon^{1/2}$-regular subgraph of $G$ with density $d_0 \pm \epsilon$, so we can apply triangle counting directly to see that $(V'_1, V'_2, V'_3)$ has triangle density $(d_{12} \pm \epsilon)(d_{13} \pm \epsilon)(d_{23} \pm \epsilon) \pm 8\epsilon^{1/2} = d_{12}d_{13}d_{23} \pm 15\epsilon^{1/2}$ (say).

**Four-cycles.** Suppose $G = (X, Y)$ is an $\epsilon$-regular bipartite graph with density $d$. We saw before that we can estimate the $C_4$-density of $G$ as $d_{C_4}(G) = d^4 \pm 10\epsilon$. A similar argument to that just given for triangles allows us to estimate the $C_4$-density of $G$ restricted to sufficient large subsets $X_0 \subseteq X$ and $Y_0 \subseteq Y$. But in fact, we can make a stronger claim: given any sufficiently dense graphs $H_X$ on $X$ and $H_Y$ on $Y$ we can estimate the density of 4-cycles $\{x_1y_1, x_1y_2, x_2y_1, x_2y_2\}$ of $G$ in which $x_1, x_2$ is an edge of $H_X$ and $y_1, y_2$ is an edge of $H_Y$.

To do this, first recall our definition of $d_{C_4}(G)$ as $\mathbb{P}_{\Phi \in \Phi}[\phi(C_4) \subseteq G]$, where we fix a 4-cycle $C_4 = (A, B)$, considered as a bipartite graph on $A = \{a_1, a_2\}$ and $B = \{b_1, b_2\}$ and let $\Phi$ be the set of all bipartite maps $\phi$ from $A \cup B$ to $X \cup Y$. Consider an auxiliary “$C_4$-homomorphism” graph $G'$, defined as follows. $G'$ is a bipartite graph with parts $X'$ and $Y'$, where $X'$ consists of all maps $\phi_1 : A \to X$ and $Y'$ consists of all maps $\phi_2 : B \to Y$. Given $\phi_1 \in X'$ and $\phi_2 \in Y'$ we can construct a bipartite map $\phi = (\phi_1, \phi_2) \in \Phi$ in an obvious manner: $\phi$ restricts to $\phi_1$ on $A$ and $\phi_2$ on $B$. We say that $\phi_1, \phi_2$ is an edge of $G'$ if $\phi = (\phi_1, \phi_2)$ is a homomorphism from $C_4$ to $G$, i.e. $\phi(a_i, b_j) \in E(G)$ for $i, j \in \{1, 2\}$.

We claim that if $0 < \epsilon \ll \epsilon'$ and $G$ is $\epsilon$-regular then $G'$ is $\epsilon'$-regular. To see this we use Theorem 2.2. By construction $G'$ has density $d_{G'}(X', Y') = d_{C_4}(G) = d^4 \pm 10\epsilon$. Also we can calculate $d_{C_4}(G') = d_{K_{4, 4}}(G)$, where $K_{4, 4}$ is the complete bipartite graph with 4 vertices in each part. To see this, label the parts of the $K_{4, 4}$ as $A^1 \cup A^2$ and $B^1 \cup B^2$, where $A^1 = \{a_1^1, a_2^1\}$, $A^2 = \{a_1^2, a_2^2\}$, $B^1 = \{b_1^1, b_2^1\}$, $B^2 = \{b_1^2, b_2^2\}$. Then we can identify a bipartite map $\psi'$ from $A \cup B$ to $X' \cup Y'$ with a bipartite map $\psi$ from $(A^1 \cup A^2) \cup (B^1 \cup B^2)$ to $X \cup Y$: define $\psi(a_i^j) = \psi'(a_i)(a_j)$ and $\psi(b_i^j) = \psi'(b_i)(b_j)$.

Now $\psi'$ is a homomorphism from $C_4$ to $G'$ if and only if $\psi$ is a homomorphism from $K_{4, 4}$ to $G$: we have $\{\psi'(a_i), \psi'(b_k)\} \in G' \forall i, k \in \{1, 2\} \iff \{\psi'(a_i)(a_j), \psi'(b_k)(b_j)\} \in G \forall i, j, k, \ell \in \{1, 2\} \iff \{\psi'(a_i^j), \psi'(b_{k\ell}^j)\} \in G \forall i, j, k, \ell \in \{1, 2\}$. Therefore $d_{C_4}(G') = d_{K_{4, 4}}(G)$. But we can estimate $d_{K_{4, 4}}(G) = d^{16} \pm O(\epsilon)$ by the counting lemma, so $d_{C_4}(G') = d_{C_4}(X', Y')^4 \pm O(\epsilon)$. Now Theorem 2.2 tells us that $G'$ is $\epsilon'$-regular.

Finally, suppose that we have graphs $H_X$ on $X$ and $H_Y$ on $Y$ with $|E(H_X)| > \epsilon'|X|^2$ and $|E(H_Y)| > \epsilon'|Y|^2$. Applying the definition of $\epsilon'$-regularity, we see that if we choose a random edge $x_1x_2$ of $H_X$ and (independently) a random edge $y_1y_2$ of $H_Y$ then $\{x_1y_1, x_1y_2, x_2y_1, x_2y_2\}$ is a 4-cycle in $G$ with probability $d^4 \pm 2\epsilon'$.

**General graphs.** Suppose $G$ is an $r$-partite graph on $V = V_1 \cup \cdots \cup V_r$ and each pair $(V_i, V_j), 1 \leq i < j \leq r$ spans an $\epsilon$-regular bipartite graph with density $d_\epsilon(G)$. Suppose $H$ is an $r$-partite graph on $Y = Y_1 \cup \cdots \cup Y_r$. Let $\Phi(Y, V)$ be the set of $r$-partite maps from $Y$ to $V$, i.e. maps $\phi : Y \to V$ with $\phi(Y_i) \subseteq V_i$ for $1 \leq i \leq r$. Recall that the $r$-partite homomorphism density of $H$ in $G$ is $d_H(G) = \mathbb{P}_{\phi \in \Phi(Y, V)}[\phi(H) \subseteq G]$. The counting lemma says that $d_H(G) = \prod_{\epsilon \in E(H)} d_\epsilon(G) \pm O(H)$, where $d_\epsilon(G)$ means that density $d_\epsilon(G)$ for which $\epsilon \in H[Y_i, Y_j]$ and the $O_H$-notation indicates that the implied constant depends only on $H$.

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7To interpret this notation, observe that, e.g., $\psi'(a_1) = \omega \in X'$ is a map from $A$ to $X$, so $\psi(a_1^2) = \omega(a_2) \in X$.
Moreover, copies of \( H \) are well-distributed in \( G \) in the following sense. Consider the auxiliary “\( H \)-homomorphism” graph \( G' \), defined as follows. \( G' \) is an \( r \)-partite graph on \( V' = V_1' \cup \cdots \cup V_r' \), where \( V_i' \) consists of all maps \( \phi_i : Y_i \to V_i \). Given \( \phi_i \in V_i' \) and \( \phi_j \in V_j' \), for some \( 1 \leq i < j \leq r \) we say that \( \phi_i \phi_j \) is an edge of \( G' \) if \( \phi_i(y_i) \phi_j(y_j) \) is an edge of \( G \) for every \( y_i, y_j \) of \( H \) with \( y_i \in Y_i \) and \( y_j \in Y_j \). Then copies of the complete graph \( K_r \) in \( G' \) correspond to homomorphisms from \( H \) to \( G \).

Suppose we have a constant hierarchy \( 0 < \epsilon \ll \epsilon' \ll \epsilon'' < 1 \). A similar argument to that used in the case of 4-cycles shows that for every \( 1 \leq i < j \leq r \), \((V_i', V_j')\) induces a subgraph of \( G' \) with density \( d_{\epsilon'' H}^{Y_i,Y_j} + O(\epsilon) \) that is \( \epsilon' \)-regular. Suppose also that we have a \(|Y_i|\)-graph \( J_i \) on \( V_i \) with \( |E(J_i)| > \epsilon'' |V_i|^2 \) for \( 1 \leq i \leq r \). Then by \( \epsilon' \)-regularity of \( G' \) and the counting lemma for \( K_r \), we see that if choose independent random edges \( e_i \) in \( J_i \) for \( 1 \leq i \leq r \) then \( \{e_1, \ldots, e_r\} \) spans a copy of \( H \) in \( G \) with probability \( \prod_{e \in E(H)} d_e(G) \pm \epsilon'' \).

3. QUASIRANDOM 3-GRAPHS

In this section we discuss the Gowers approach to quasirandomness in 3-graphs: our exposition will be quite condensed, and for more details we refer the reader to that given in [7]. Although we will later repeat this discussion for general \( k \)-graphs, we feel it is helpful to first present the case \( k = 3 \), which is simpler to grasp for a reader new to the subject. We conclude this section by describing our quasirandom counting lemma for the special case of the distribution of octahedra in 3-graphs, which is analogous to the distribution of 4-cycles in graphs described above; our solution of Mubayi’s conjecture, Theorem 1.1, will only make use of this case.

3.1. Two Cautionary Examples

When considering how to generalise regularity from graphs to 3-graphs, a natural first attempt is to take a 3-graph \( H \) and partition its vertices as \( V(H) = V_1 \cup \cdots \cup V_k \), for some \( k(\epsilon) \), so that all but at most \( \epsilon k^3 \) triples \( (V_a, V_b, V_c) \) span a tripartite 3-graph that is \( \epsilon \)-\( \text{vertex-regular} \), meaning that for any \( V'_{a} \subseteq V_a, V'_{b} \subseteq V_b, V'_{c} \subseteq V_c \) with \( |V'_a| \geq \epsilon |V_a|, |V'_b| \geq \epsilon |V_b|, |V'_c| \leq \epsilon |V_c| \) we have \( \mathbb{E}_{x \in V'_a, y \in V'_b, z \in V'_c} H(x,y,z) = \mathbb{E}_{x \in V_a, y \in V_b, z \in V_c} H(x,y,z) \pm \epsilon \). This is indeed possible, as shown by Chung and Graham [2], with a proof closely modelled on that of the graph regularity lemma. This result is often known as the “weak hypergraph regularity lemma,” although it does have some applications, the following example of Rödl demonstrates quite dramatically.

**Example.** Take a set \( X = X_1 \cup X_2 \cup X_3 \cup X_4 \), where \( X_i, 1 \leq i \leq 4 \) are pairwise disjoint sets of size \( n \). Consider a random orientation of the complete 4-partite graph on \( X \), i.e., for every \( x_i \in X_i, x_j \in X_j, 1 \leq i < j \leq 4 \) we choose the arc \( x_i x_j \) or the arc \( x_j x_i \), each choice having probability 1/2, all choices being independent. Define a 3-graph \( H \) on \( X \) to consist of all triples \( x_i x_j x_k \) that induce a cyclic triangle in the orientation. Then with high probability each triple \( (X_i, X_j, X_k) \) spans an \( \epsilon \)-\( \text{vertex-regular} \) triangle when \( n \gg \epsilon^{-1} \), and this will remain true even if we partition \( H \) into \( k(\epsilon) \) parts. However, it is easy to see that \( H \) does not contain any copy of the tetrahedron \( K_4^3 \).

*Random Structures and Algorithms* DOI 10.1002/rsa
In the light of this example, we might informally say that vertex-regularity uses a random model of a 3-graph in which triples are randomly chosen as edges with some probability \( p \), but that a counting lemma fails because of potential correlations between edges that share a pair of vertices. Our second example (adapted from [7]) points the way to a better model.

**Example.** Take a set \( X = X_1 \cup X_2 \cup X_3 \), where \( X_i, 1 \leq i \leq 3 \) are pairwise disjoint sets of size \( n \). For each pair \( 1 \leq i < j \leq 3 \) define a random bipartite graph \( G_{ij} \), in which each \( x_i x_j, x_i \in X_i, x_j \in X_j \) is an edge with probability \( d_{ij} \), all choices being independent. Consider \( \Delta = \Delta(G_{12} \cup G_{13} \cup G_{23}) \), the triangles spanned by these graphs. Define a random 3-graph \( H \) by taking each triangle of \( \Delta \) to be an edge with probability \( d_{123} \). With high probability \( H \) has density \( d(H) = |E(H)|/n^3 = d_{12}d_{13}d_{23} + o(1) \). Now consider the number of copies of some fixed 3-graph \( F \) in \( H \), or for a better parallel with our analysis of quasirandom graphs consider the “tripartite \( F \)-homomorphism density” \( d_F(H) = \mathbb{P}_{\phi \in \Phi} [\phi(F) \subseteq H] \), where \( \Phi \) is the set of tripartite maps from \( F \) to \( H \). We can no longer estimate this by \( d(H)^{|F|} \), as might at first be expected from the analysis for graphs. For a simple example, suppose that \( F \) consists of two edges sharing a pair of vertices; say \( F \) has vertex set \( A = A_1 \cup A_2 \cup A_3 \) with \( A_1 = \{a_1\}, A_2 = \{a_2\}, A_3 = \{a_1^2, a_2^2\} \) and edges \( a_1a_2a_1^2, a_1a_2a_2^2 \). Then with high probability \( d_F(H) = d_{12}d_{13}d_{23}d_{123}^2 + o(1) \).

### 3.2. Quasirandom Complexes

The second example above shows that a counting lemma for 3-graphs must take account of densities of pairs, as well as densities of triples. Thus we are led to define quasirandomness not for a 3-graph in isolation, but for a simplicial complex consisting of a 3-graph together with all subsets of its edges. We say that \( H \) is a tripartite 3-complex\(^8\) on \( X = X_1 \cup X_2 \cup X_3 \) if we have \( H = \bigcup_{i \subseteq \{1,2,3\}} H_i \), where \( H_\emptyset = \{\emptyset\} \), \( H_{\{i\}} \) is a subset of \( X_i \) for \( 1 \leq i \leq 3 \), \( H_{\{ij\}} \) is a bipartite graph with parts \( H_{\{i\}}, H_{\{j\}} \) for \( 1 \leq i < j \leq 3 \), and \( H_{\{1,2,3\}} \) is a 3-graph contained in the set of triangles spanned by \( H_{\{1,2\}}, H_{\{1,3\}} \) and \( H_{\{2,3\}} \). Of course, the interesting part of this structure is the 3-graph together with its underlying graphs: we usually take \( H_{\{i\}} = X_i \) for \( 1 \leq i \leq 3 \), and we only include \( H_\emptyset \) so that we are formally correct when referring to \( H \) as a (simplicial) complex. Sometimes we allow \( H_{\{i\}} \) to be a strict subset of \( X_i \), but then we can reduce to the usual case by redefining the ground set as \( X' = X'_1 \cup X'_2 \cup X'_3 \), where \( X'_i = H_{\{i\}}, 1 \leq i \leq 3 \).

As usual we identify each \( H_i \) with its characteristic function, for example \( H_{123}(x_1, x_2, x_3) \) is 1 if \( x_1x_2x_3 \) is an edge of \( H_{123} \), otherwise 0 (henceforth we write 123 instead of \( \{1,2,3\} \), etc. for more compact notation). We let \( H_{123}^* \) denote the set of triangles spanned by \( H_{\{1,2\}}, H_{\{1,3\}} \) and \( H_{\{2,3\}} \), and also the characteristic function of this set. Then we define the relative density by

\[
d_{123}(H) = |H_{123}|/|H_{123}^*|.
\]

In words, it is the proportion of graph triangles that are triples of the complex. Note that we can describe the densities of the bipartite graph \( H_{ij} \) with a similar notation: we let \( H_{ij}^* \) denote the pairs spanned by \( H_i \) and \( H_j \), i.e. the complete bipartite graph with parts \( X_i \) and \( X_j \), and then \( H_{ij} \) has density \( d_{ij}(H) = |H_{ij}|/|H_{ij}^*| \).

We have seen that vertex-regularity is not the correct notion for defining quasirandomness in 3-graphs; it turns out that the other two properties described in Theorem 2.2 can be

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\( ^8 \)We prefer this term to “chain,” which is used in [7].

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generalized to a useful concept for 3-graphs. These properties were defined using 4-cycles for graphs; for 3-graphs they are defined in terms of the octahedron $O_3$. This is a tripartite 3-graph with vertex set $A = A_1 \cup A_2 \cup A_3$, where $A_i = \{a_i^0, a_i^1\}$, $1 \leq i \leq 3$ and 8 edges $a_i^e, a_i^f, a_i^g$, where $e = (e_1, e_2, e_3)$ ranges over vectors in $\{0, 1\}^3$. One could define quasirandomness of a complex in terms of the density of octahedra (by analogy with property 2 in Theorem 2.2), but we will follow Gowers and consider “octahedra with respect to the balanced function” (by analogy with property 3 in Theorem 2.2). The balanced function is defined as

$$\overline{H}_{123}(x_1, x_2, x_3) = H_{123}(x_1, x_2, x_3) - d_{123}(H)H_{123}(x_1, x_2, x_3),$$

i.e. for $x_1 \in X_1, x_2 \in X_2, x_3 \in X_3$, $\overline{H}_{123}(x_1, x_2, x_3)$ equals $1 - d_{123}(H)$ if $x_1x_2x_3$ is an edge of $H_{123}$, equals $-d_{123}(H)$ if $x_1x_2x_3$ is a triangle in $H_{123}$ but not an edge of $H_{123}$, or equals 0 otherwise. As before we note that $\mathbb{E}_{x_1 \in X_1, x_2 \in X_2, x_3 \in X_3} \overline{H}_{123}(x_1, x_2, x_3) = 0$. Let $\Phi$ be the set of tripartite maps from $A = A_1 \cup A_2 \cup A_3$ to $X = X_1 \times X_2 \times X_3$. For any function $f : X_1 \times X_2 \times X_3 \rightarrow [-1, 1]$ we define

$$O_3(f) = \mathbb{E}_{\phi \in \Phi} \prod_{e \in E(O_3)} f(\phi(e)).$$

We say that $f$ is $\eta$-quasirandom (with respect to $H$) if $O_3(f) < \eta(d_{12d_1d_2}d_3)^4$ (note that we are assuming that $d_1 = d_2 = d_3 = 1$). One can think of this as saying that $O_3(f)$ is small compared to the density of the “graph octahedron,” i.e. the complete tripartite graph $K(2, 2, 2)$, as this has density about $(d_{12d_1d_2}d_3)^4$ when the graphs $H_{ij}$ are quasirandom (as they will be). Then we say that $H_{123}$ is $\eta$-quasirandom (with respect to $H$) if its balanced function $\overline{H}_{123}$ is $\eta$-quasirandom.

### 3.3. Quasirandom Decomposition and the Counting Lemma

Next we will see why the notion of quasirandomness given in the previous subsection is useful: there is a decomposition theorem and a counting lemma.

To describe the decomposition theorem we need to first think about the kind of structure that will arise in a decomposition into simplicial complexes. It will not simply be a disjoint union of simplicial complexes: just as in a graph decomposition the various bipartite graphs that will arise in a decomposition into simplicial complexes. It will not simply be a disjoint useful: there is a decomposition theorem and a counting lemma.

Next we will see why the notion of quasirandomness given in the previous subsection is useful: there is a decomposition theorem and a counting lemma.

To describe the decomposition theorem we need to first think about the kind of structure that will arise in a decomposition into simplicial complexes. It will not simply be a disjoint union of simplicial complexes: just as in a graph decomposition the various bipartite graphs share vertices, in a 3-graph decomposition the various tripartite complexes may share vertices and pairs. Suppose $H$ is an $r$-partite 3-graph on $X = X_1 \cup \cdots \cup X_r$. A decomposition of $H$ is described by partitions $X_i = X_i^1 \cup \cdots \cup X_i^{n_i}$ of each part $X_i$ into subsets, and partitions $K(X_i, X_j) = G_{ij}^1 \cup \cdots \cup G_{ij}^{n_{ij}}$ of each complete bipartite graph with parts $X_i, X_j$ into bipartite subgraphs. Let $P$ denote this “partition system.” For each edge $e$ of $H$ there is an induced complex $H(e, P)$ defined as follows. The vertex set of $H(e, P)$ is $X_i^{a_{ij}} \cup X_j^{b_{ij}} \cup X_k^{c_{ij}}$, where $e = x_i x_j x_k$ with $x_i \in X_i^{a_{ij}}, x_j \in X_j^{b_{ij}}, x_k \in X_k^{c_{ij}}$. The pairs in $H(e, P)$ are given by the restriction of $G_{ij}^{a_{ij}} \cup G_{ik}^{b_{ik}} \cup G_{jk}^{c_{jk}}$ to the vertex set of $H(e, P)$, where $x_i x_j \in G_{ij}^{a_{ij}}, x_i x_k \in G_{ik}^{b_{ik}}, x_j x_k \in G_{jk}^{c_{jk}}$. The triples in $H(e, P)$ are those edges of $H$ that also form triangles in the pairs of $H(e, P)$.

The following decomposition is a variant of Theorem 8.10 in [7]. (We defer a justification of the differences here and elsewhere until the appendix.)

**Theorem 3.1.** Suppose $0 < 1/n \ll d_1 \ll \eta_2 \ll d_2 \ll \eta_3 \ll 1/r, \epsilon < 1$ and $H$ is an $r$-partite 3-graph on $X = X_1 \cup \cdots \cup X_r$, where each $|X_i| \geq n$. Then there is a partition system $P$ consisting of partitions $X_i = X_i^0 \cup X_i^1 \cup \cdots \cup X_i^{n_i}$ with each $n_i \leq 1/d_1$ and partitions $K(X_i, X_j) = G_{ij}^0 \cup \cdots \cup G_{ij}^{n_{ij}}$ with each $n_{ij} \leq d_2^{-1/2}$, such that.

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• each exceptional class \(X^0_i\) has size at most \(\epsilon |X_i|\), for \(1 \leq i \leq r\),
• the classes \(X^1_i, 1 \leq i \leq r, 1 \leq a_i \leq n_i\) all have the same size, and
• if \(e\) is a randomly chosen edge of \(H\), then with probability at least \(1 - \epsilon\), in the induced complex \(H(e, P)\), the graphs are \(\eta_2\)-quasirandom, and the 3-graph is \(\eta_3\)-quasirandom with respect to \(H(e, P)\).

To apply this decomposition we need a counting lemma, estimating the number of homomorphisms from a fixed \(r\)-partite 3-complex \(F\) on \(Y = Y_1 \cup \cdots \cup Y_r\) to a quasirandom \(r\)-partite 3-complex \(H\) on \(X = X_1 \cup \cdots \cup X_r\). Let \(\Phi(\gamma, \chi)\) denote the set of \(r\)-partite maps \(\phi : Y \to X\).

We define the partite homomorphism density of \(F\) in \(H\) by

\[
d_F(H) = \mathbb{P}_{\phi \in \Phi(\gamma, \chi)} \prod_{e \in F} H(\phi(e)).
\]

The following theorem is a variant of Corollary 5.2 in [8].

**Theorem 3.2.** Suppose \(0 \leq \eta_2 \ll d_2 \ll \eta_3 \ll d_3, \epsilon, 1/|F| < 1\), that \(H\) and \(F\) are \(r\)-partite 3-complexes on \(X = X_1 \cup \cdots \cup X_r\) and \(Y = Y_1 \cup \cdots \cup Y_r\) respectively, every graph \(H_{ij}\) is \(\eta_2\)-quasirandom with density at least \(d_2\), and that every 3-graph \(H_{ijk}\) is \(\eta_3\)-quasirandom with relative density at least \(d_3\) with respect to \(H\). Then

\[
d_F(H) = \mathbb{P}_{\phi \in \Phi(\gamma, \chi)} \prod_{e \in F} H(\phi(e)) = (1 + \epsilon) \prod_{e \in F} d_e(H).
\]

**Remarks.**

1. The special case when \(F\) is a single edge shows that the absolute density of a triple in \(H\) is well-approximated by the product of its relative densities, viz.

\[
d(H_{ijk}) = |H_{ijk}|/|X_i||X_j||X_k| = (1 + \epsilon)d_1d_2d_3d_{ijk}.
\]

2. It is important to note the hierarchy of the parameters, as this is where some of the technical difficulties in hypergraph regularity lie. In the decomposition the number of graphs \(G_{ij}\) in the partition may be much larger than \(\eta_1^{-1}\), so the counting lemma has to cope with graphs \(H_{ij}\) with density much smaller than \(\eta_3\), the quasirandomness parameter for triples; fortunately it can be arranged that \(\eta_2\), the quasirandomness parameter for pairs, is smaller still, and this is sufficient for the counting lemma.

**3.4. Counting Homomorphisms to Functions and Uniform Edge-Distribution**

We will also need the following more general theorem, a variant of Theorem 5.1 in [8], which allows estimation of the number of \(F\)-homomorphisms with respect to functions supported on \(H\); the generalization is similar to that which we saw earlier from counting 4-cycles in quasirandom graphs to estimating \(C_4(f)\) for a function \(f\). Intuitively it expresses the fact that, in a quasirandom 3-complex, sets of size \(k \leq 3\) are almost uncorrelated with

---

9Note that we use the same notation for the normal homomorphism density: it will be clear from the context which is intended, and in any case they are roughly equivalent for our purposes in that they only differ by a constant factor.

10Recall that \(d_e(H)\) means the density \(d_1(H)\) for which \(e \in H_t\).
sets of size less than $k$. First we need some notation. Suppose $X = X_1 \cup \cdots \cup X_r$ is a set partitioned into $r$ parts. For $A \subseteq [r]$ let $K_A(X)$ be the set of all $A$-tuples with one point in each $X_i$, $i \in A$. We say that $e \in K_A(X)$ has index $A$ and we sometimes abuse notation and use $e$ instead of $A$; thus, if $H$ is an $r$-partite 3-complex on $X$ and $e \in H_A$ for some $A \subseteq [r]$, $|A| \leq 3$ then $H_e = H_A$.

**Theorem 3.3.** Suppose $0 \leq \eta_2 \ll d_2 \ll \eta_3 \ll d_3, \epsilon, 1/|F| < 1$, that $H$ and $F$ are $r$-partite 3-complexes on $X = X_1 \cup \cdots \cup X_r$ and $Y = Y_1 \cup \cdots \cup Y_r$, respectively, every graph $H_{ij}$ is $\eta_2$-quasirandom with density at least $d_2$, and that every 3-graph $H_{ijk}$ is $\eta_3$-quasirandom with relative density at least $d_3$ with respect to $H$. Suppose $F_0$ is a subcomplex of $F$, for each $e \in F$ we have a function $f_e$ on $K_e(X)$ with $f_e = f_e H_e$, and $f_e = H_e$ for all $e \in F \setminus F_0$. Then

$$
\mathbb{E}_{\phi \in \Phi(Y,X)} \left[ \prod_{e \in F} f_e(\phi(e)) \right] = \prod_{e \in F \setminus F_0} d_e(H) \cdot \mathbb{E}_{\phi \in \Phi(Y,X)} \left[ \prod_{e \in F_0} f_e(\phi(e)) \right] \pm \epsilon \prod_{e \in F} d_e(H).
$$

Note that Theorem 3.2 is the case of Theorem 3.3 when $F_0 = \emptyset$. The following corollary shows that quasirandomness implies vertex-uniformity, the analogue of graph regularity that we mentioned earlier as being too weak for a useful hypergraph regularity theory.

**Corollary 3.4.** Suppose $0 \leq \eta_2 \ll d_2 \ll \eta_3 \ll d_3, \epsilon < 1$, that $H$ is a tripartite 3-complex on $X = X_1 \cup X_2 \cup X_3$ with $H_i = X_i$, $1 \leq i \leq 3$, that $H_{12}, H_{13}$ and $H_{23}$ are $\eta_2$-quasirandom with density at least $d_2$, and that $H_{123}$ is $\eta_3$-quasirandom with relative density at least $d_3$ with respect to $H$. Suppose $W_i \subseteq X_i$, $1 \leq i \leq 3$ and let $H_{123}[W]$ be the restriction of $H_{123}$ to $W = W_1 \cup W_2 \cup W_3$. Then

$$
\frac{|H_{123}[W]|}{|X_1||X_2||X_3|} = \frac{|H_{123}|}{|X_1||X_2||X_3|} \cdot \left( \frac{|W_1||W_2||W_3|}{|X_1||X_2||X_3|} \pm \epsilon \right).
$$

**Proof:** Set $Y_i = \{i\}$ for $1 \leq i \leq 3$, $F$ equal to all subsets of $Y = Y_1 \cup Y_2 \cup Y_3$ and $F_0 \subseteq F$ equal to the subsets of size at most 1. Apply Theorem 3.3 with $f_i(x_i) = 1_{x_i \in W_i}$ and $f_e = H_e$ for $e \in F \setminus F_0$. Since $H_i = X_i$, $1 \leq i \leq 3$ we have

$$
\prod_{e \in F \setminus F_0} d_e(H) = \prod_{e \in F} d_e(H). \quad \text{Also} \quad \mathbb{E}_{\phi \in \Phi(Y,X)} \prod_{e \in F_0} f_e(\phi(e)) = \prod_{i=1}^3 |W_i|/|X_i|. \quad \text{Now} \quad \mathbb{E}_{\phi \in \Phi(Y,X)} \left[ \prod_{e \in F} f_e(\phi(e)) \right] = |H_{123}[W]|/|X_1||X_2||X_3| \text{ so Theorem 3.3 gives the result.} \quad \square
$$

### 3.5. Uniform Distribution of Octahedra in Quasirandom 3-Complexes

Now we will establish a further special case of our quasirandom counting lemma, the uniform distribution of octahedra in quasirandom 3-complexes, which is the heart of our solution of Mubayi’s conjecture. Recall that the octahedron $O_3$ is a tripartite 3-graph with vertex set $A = A_1 \cup A_2 \cup A_3$, where $A_i = \{a_i^0, a_i^1\}$, $1 \leq i \leq 3$ and 8 edges $a_i^1 a_j^2 a_k^3$, where $e = (e_1, e_2, e_3)$ ranges over vectors in $\{0, 1\}^3$. We write $O_3^c$ for the tripartite 3-complex generated by $O_3$ (i.e. it contains all subsets of the edges of $O_3$) and $O_3^c = O_3^c \setminus O_3$ for the tripartite 2-complex of strict subsets of edges of $O_3$.

\[\text{Here we are using the notation described before the theorem: if } e \text{ has index some } A \subseteq [r], |A| \leq 3, \text{ then } K_e(X) = K_A(X) = \prod_{i \in A} X_i. \text{ The equation } f_e = f_e H_e \text{ is a concise way of saying that } f_e \text{ is supported on } H_e = H_A.\]
Suppose $H$ is a tripartite 3-complex on $X = X_1 \cup X_2 \cup X_3$, where $H_i = X_i$, $1 \leq i \leq 3$. We define an “auxiliary” tripartite 3-complex $H'$ on $X' = X'_1 \cup X'_2 \cup X'_3$, where $X'_i$ consists of all maps $\phi : A_i \to X_i$, $1 \leq i \leq 3$ as follows. The triples $H'_{123}$ consist of all $\phi_1 \phi_2 \phi_3$ for which $\phi = (\phi_1, \phi_2, \phi_3)$ is a homomorphism from $O_3$ to $H$; the pairs $H'_ij$ consist of all $\phi_i j$ such that $(\phi_i, \phi_j)$ is a homomorphism from $(O_3)^{\leq i}_j$ to $H_j$, for $1 \leq i < j \leq 3$.

**Theorem 3.5.** Suppose $0 \leq 1/n \leq \eta_2 \leq \eta_2' \leq \eta_3 \leq \eta_3' \leq \epsilon, \eta_3 < 1$, that $H$ is a tripartite 3-complex on $X = X_1 \cup X_2 \cup X_3$, where $H_i = X_i$ with $|X_i| \geq n$, $1 \leq i \leq 3$, that $H_{12}, H_{13}$ and $H_{23}$ are $\eta_2$-quasirandom with density at least $d_2$, and that $H_{123}$ is $\eta_3$-quasirandom with relative density at least $d_3$ with respect to $H$. Then, in the auxiliary complex $H'$, each $H'_{ij}$ is $\eta_2'$-quasirandom with density $(1 \pm \epsilon_2) d_{ij}(H)^4$, and $H'_{123}$ is $\eta_3'$-quasirandom with relative density $(1 \pm \epsilon_3) d_{123}(H)^8$ with respect to $H$.

Furthermore, suppose that $G_i$ is a graph on $X_i$ with $|E(G_i)| > \epsilon |X_i|^2$, $1 \leq i \leq 3$ and we randomly and independently select edges $x_i x'_i$ of $G_i$, $1 \leq i \leq 3$. Then $(x_1 x'_1, x_2 x'_2, x_3 x'_3)$ spans an octahedron in $H$ with probability $(1 \pm \epsilon) d_{12}(H)^4 d_{13}(H)^4 d_{23}(H)^4 d_{123}(H)^8$.

**Proof.** It may help the reader to note that the prime symbol $'$ is a visual cue, distinguishing quantities referring to $H'$ from those referring to $H$. We write $d_i = d_i(H)$ and $d'_i = d_i(H')$ for $I \subseteq \{1, 2, 3\}, |I| = 2, 3$.

The quasirandomness of each $H'_{ij}$ follows from the discussion of quasirandom counting of 4-cycles in Subsection 2.3, so it remains to establish the claim for $H'_{123}$. The final statement will then follow from Corollary 3.4 and Eq. (1).

First we estimate the relative density of $H'$ using Eq. (1): we have (using $\epsilon_3/3$ instead of $\epsilon$)

$$d(H'_{123}) = |H'_{123}| / |X_1|^2 |X_2|^2 |X_3|^2 = \mathbb{P}_{\phi \in \Phi(A,X)}[\phi(O_3^+) \subseteq H] = (1 + \epsilon_3/3) d_{12} d_{13} d_{23} d_{123}^4 d_{123}^8,$$

and

$$d(H''_{123}) = |H''_{123}| / |X_1|^2 |X_2|^2 |X_3|^2 = \mathbb{P}_{\phi \in \Phi(A,X)}[\phi(O_3^+) \subseteq H] = (1 + \epsilon_3/3) d_{12} d_{13} d_{23} d_{123},$$

so $d_{123}(H') = |H'_{123}| / |H''_{123}| = (1 \pm \epsilon_3) d_{123}^8$.

Next we need to estimate $O_3(H'_{123})$. Write $f'_0 = H'_{123}$ (the characteristic function) and $f_1 = d_{123} H'_{123}$, so that $H'_{123} = f'_0 - f'_1$. Then

$$O_3(H'_{123}) = \mathbb{P}_{\phi \in \Phi(A,X)} \prod_{e \in O_3} H'_{123}(\phi'(e)) = \sum_{j} (-1)^{j} \sum_{\epsilon \in O_3} \prod_{j \in e} f'_j(\phi'(e)),$$

where $j$ ranges over all $\{0, 1\}$-sequences $(j_e)_{e \in O_3}$ and $\sum_j = \sum_{e \in O_3} j_e$. We interpret each summand as counting homomorphisms to functions from $O_3^+ O_3$, by which we mean (with unusual notation!) the complete tripartite 3-graph with 4 vertices in each class (cf. our earlier analysis of $C_4$ via $K_{4,4}$.) To accomplish this we consider a complete tripartite 3-graph, which we call $O_3^+ O_3$, with 3 parts $A_{i,j} \cup A_{i.j}$, $1 \leq i \leq 3$, where $A_{i,j} = \{a_{i,j}^0, a_{i,j}^1\}$, $1 \leq i \leq 3, j = 0, 1$. We think of the edges of $O_3^+ O_3$ as a union of 8 copies of $O_3$, indexed by $O_3$, thus $E(O_3^+ O_3) = \bigcup_{e \in O_3} O_3^e$, where, writing $e = \{a_{i,j}^0, a_{i,j}^1, a_{i,j}^2, a_{i,j}^3\}$, $O_3^e$ is the copy of $O_3$ with parts $\{a_{i,j}^0, a_{i,j}^1\}$, $1 \leq i \leq 3$ (see Fig. 1).\(^{13}\)

\(^{12}\)To avoid confusion, note that $d_2$, $d_3$ are lower bounds for $d_{ij}$, $d_{123}$ respectively, not relative densities of $H_2$, $H_3$.

\(^{13}\)Some readers may like to think of this construction as a “tripartite tensor product.”
We identify a tripartite map $\phi'$ from $A = A_1 \cup A_2 \cup A_3$ to $X' = X'_1 \cup X'_2 \cup X'_3$ with a tripartite map $\phi$ from $AA = (A_{1,0} \cup A_{1,1}) \cup (A_{2,0} \cup A_{2,1}) \cup (A_{3,0} \cup A_{3,1})$ to $\tilde{X}$ by defining $\phi(a_{ij}^k) = \phi'(a_{ij}^k)$. Thus for $e = \{a_{i1}^1, a_{i2}^2, a_{i3}^3\} \in O_3$, $\phi^e = (\phi'(a_{i1}^1), \phi'(a_{i2}^2), \phi'(a_{i3}^3))$ acts on $O'_3$. We can rewrite the summands above using

$$\prod_{e \in O_3} f'_e(\phi'(e)) = \prod_{e \in O_3} \prod_{\hat{e} \in O'_3} f'_e(\phi'(\hat{e})) = \prod_{E \in O_3 O_3} f_{(E)}(\phi(E)),$$

where we write $e(E)$ for that $e \in O_3$ such that $E \in O'_3$, $f_0 = H_{123}$ (the characteristic function) and $f_1 = (d'_{123})^{1/8} H_{123} = (1 \pm \epsilon^1_3/4) d_{123} H_{123}$.

Now we can estimate the contribution from the summand corresponding to $j = (j_e)_{e \in O_3}$ using the counting lemma for homomorphisms to $H$ from a tripartite 3-complex $K_j$ on $AA$, in which the pairs form the complete tripartite graph on $AA$ and the triples are $\cup_{e_j = 0} O'_3$. We have

$$E_{\phi' \in \Phi(AX)} \prod_{e \in O_3} f'_e(\phi'(e)) = E_{\phi \in \Phi(AX)} \prod_{E \in O_3 O_3} f_{(E)}(\phi(E))$$

$$= (d'_{123})^j \sum_{\phi \in \Phi(AX)} [\phi(K_j) \subseteq H]$$

$$= ((1 \pm \epsilon_3) d_{12}^{16} d_{13}^{16} d_{23}^{16})^j (1 + \epsilon_3) d_{12}^{16} d_{13}^{16} d_{23}^{16} - 8 \sum_{j}$$

$$= (1 + O(\epsilon^3)) (d'_{12} d'_{13} d'_{23})^j (d'_{123})^8.$$
where the implicit constant in the $O(\cdot)$ notation is absolute, say 100. Therefore

$$O_3(\overline{H}_{123}) = \sum_j (-1)^{j_1}(1 \pm O(\epsilon_3))(d'_{12}d'_{13}d'_{23})^4(d'_{123})^8$$
$$= \sum_j O(\epsilon_3)(d'_{12}d'_{13}d'_{23})^4(d'_{123})^8 < \eta'_3(d'_{12}d'_{13}d'_{23})^4,$$

since $\epsilon_3 \ll \eta'_3$. This proves that $H'_{123}$ is $\eta'_3$-quasirandom with relative density $(1 \pm \epsilon)d_{123}(H)^8$ with respect to $H$. The third statement of the theorem follows from Corollary 3.4, so we are done.

4. THE FANO PLANE

In this section we prove Theorem 1.1, thus establishing Mubayi’s conjecture. Our key tool will be Theorem 3.5, our result above on uniform distribution of octahedra in quasirandom 3-complexes. We also make use of the idea of stability, or approximate structure, which can be traced back to work of Erdős and Simonovits in the 60s in extremal graph theory. Informally stated, a stability result tells us about the structure of configurations that are close to optimal in an extremal problem: for example, a triangle-free graph with $n^2/4 - o(n^2)$ edges differs from a complete bipartite graph by $o(n^2)$ edges. Such a result is interesting in its own right, but somewhat surprisingly it is often a useful stepping stone in proving an exact result. Indeed, it was developed by Erdős and Simonovits to determine the exact Turán number for $k$-critical graphs.

To further explain our method we first need some definitions. The Fano plane $\text{PG}_2(2)$ is the projective plane over $\mathbb{F}_2 = \{0, 1\}$, the field with two elements: it is a 3-graph in which the vertex set consists of the 7 non-zero vectors of length 3 over $\mathbb{F}_2$, and the 7 edges are those triples of vectors $abc$ with $a + b = c$ (see Fig. 2). A natural construction of a 3-graph not containing the Fano plane is a balanced complete bipartite 3-graph $H_2(n)$, which we recall is obtained by partitioning a set into two parts $X = X_1 \cup X_2$ with sizes as equal as possible and taking as edges all triples that are not contained in either part (i.e. that intersect both parts): it is easy to verify that the Fano plane is not bipartite, so is not contained in $H_2(n)$.

Random Structures and Algorithms DOI 10.1002/rsa
In 1976 Sós [26] conjectured that $H_2(n)$ gives the exact value of \( \text{ex}(n, PG_2(2)) \). This was established for large \( n \), together with the characterization of \( H_2(n) \) as the unique maximizing configuration, independently and simultaneously by Keevash and Sudakov [13] and Füredi and Simonovits [6]. In [13] the following stability method was used: the first step was to show that a Fano-free 3-graph with \( (3/4 - o(1))^{n/3} \) edges differs from a complete bipartite 3-graph by \( o(n^3) \) edges; then, the second step was to examine the possible imperfections in structure in a 3-graph that is close to being bipartite, showing that they exclude more edges in a Fano-free 3-graph than a complete bipartite 3-graph with no imperfections.

In considering the codegree problem, Mubayi [19] noted that the same construction \( H_2(n) \) gives a lower bound \( \text{ex}_2(n, PG_2(2)) \geq \lceil n/2 \rceil \) and proved that \( \pi_2(PG_2(2)) = 1/2 \), so this bound is asymptotically best possible. In our solution of the exact problem, the first step will be to use Theorem 3.5 to obtain some partial structure of Fano-free 3-graphs with nearly maximal minimum codegree, namely that a Fano-free 3-graph \( H \) on \( n \) vertices with minimum 2-degree \( (1/2 - o(1))n \) has a sparse set of size \( (1/2 - o(1))n \). The second step (fairly straightforward) is to complete the approximate structure, showing that \( H \) differs from a complete bipartite 3-graph by \( o(n^3) \) edges. Then the final step is to examine the possible imperfections in structure to show that the largest minimum codegree is attained with no imperfections: this is similar in spirit to the analysis of the normal Turán problem, although there are some additional technical difficulties for the codegree problem. We will divide the three steps of the proof into subsections of this section.

### 4.1. The Sparse Set

The first step of the proof is the following application of Theorem 3.5.

**Theorem 4.1.** Suppose \( 0 < 1/n \ll \theta \ll \epsilon \ll 1 \), \( H \) is a 3-graph on a set \( X \) of \( n \) vertices that does not contain a Fano plane, and \( \delta_H(H) > (1/2 - \theta)n \). Then there is a set \( Z \subseteq X \) with \( |Z| > (1/2 - \epsilon)n \) such that \( d(H[Z]) < \epsilon \), i.e. \( Z \) spans at most \( \epsilon |Z|^3/6 \) edges of \( H \).

The idea of the proof is as follows. First we apply Theorem 3.1 to decompose \( H \) into induced 3-complexes, most of which are quasirandom. Then a simple double-counting gives us a vertex \( x \) such that there are about half of the vertex classes in the partition on which \( x \) has a dense neighborhood graph. For each triple of such classes, each quasirandom induced 3-complex on these classes must be very sparse; otherwise Theorem 3.5 gives us an octahedron in which the pairs from each class are in the neighborhood graph of \( x \), and this structure contains a Fano plane. This implies that we have about half of the vertex classes in the partition inducing a sparse subhypergraph of \( H \).

**Proof of Theorem 4.1.** We operate with a constant hierarchy \( 0 < 1/n \ll d_1 \ll \eta_2 \ll d_2 \ll \eta_3 \ll d_1, \gamma \ll \theta \ll \epsilon \ll 1 \). Set \( r = \lceil \theta^{-1} \rceil \) and let \( X = X_1 \cup \cdots \cup X_r \) be an arbitrary partition with \( ||X_i| - n/r| < 1 \) for \( 1 \leq i \leq r \). Let \( H_1 \) be the edges of \( H \) that respect this partition, i.e. have at most one point in each \( X_i \). Applying Theorem 3.1 to the generated \( r \)-partite 3-complex \( H_1^r \), using \( \gamma \) instead of \( \epsilon \), we obtain a partition system \( P \) consisting of partitions \( X_i = X_i^0 \cup X_i^1 \cup \cdots \cup X_i^{|X_i|/r} \) with each \( n_i \leq 1/d_1 \) and partitions \( K(X_i, X_j) = G_i^j \cup \cdots \cup G_j^{n_j} \) with each \( n_j \leq d_2^{-1/2} \), such that,

- each exceptional class \( X_i^0 \) has size at most \( \gamma |X_i| \), for \( 1 \leq i \leq r \),
- the classes \( X_i^{n_i}, 1 \leq i \leq r, 1 \leq a_i \leq n_i \) have the same size, say \( m \), and
• if $e$ is a randomly chosen edge of $H_1$, then with probability at least $1 - \gamma$, in the induced complex $H^\varepsilon_1(e, P)$, the graphs are $\eta_2$-quasirandom, and the 3-graph is $\eta_3$-quasirandom with respect to $H^\varepsilon_1(e, P)$.

Next we find a vertex $x$ as described in the sketch before the proof. Consider the set of all pairs contained in the vertex classes of $P$: denote it $Y = \cup_{(i, i_t) \in I} Y_{i, i_t}$, where $I = \{(i, t_i) : 1 \leq i \leq r, 1 \leq t_i \leq n_i\}$ and $Y_{i, i_t} = (\binom{x_i}{2})$. (Note that we are not considering any exceptional classes here.) Double-counting gives

$$\sum_{x \in X} |N_H(x) \cap Y| = \sum_{y \in Y} |N_H(Y) \cap X| > |Y|(1/2 - \theta)n,$$

so we can choose $x$ with $|N_H(x) \cap Y| > (1/2 - \theta)|Y|$. Let $C = \{(i, t_i) : |N_H(x) \cap Y_{i, t_i}| > \theta|Y_{i, t_i}|\}$. Then, since $|X^{i_t}| = m$ for every $(i, t_i) \in I$,

$$(1/2 - \theta)|I|\binom{m}{2} = (1/2 - \theta)|Y| < |N_H(x) \cap Y| < |C|\binom{m}{2} + |I \setminus C|\theta\binom{m}{2},$$

so $|C| > (1/2 - 2\theta)|I|$. Let $Z = \cup_{(i, t_i) \in C} X_{i, t_i}$. Then $|Z| = |C|m > (1/2 - 3\theta)n$.

Consider any induced tripartite 3-complex $H^\varepsilon_1(e, P)$ in which the three vertex classes are $X^{i_t}_a$, $1 \leq a \leq 3$ with $(i_a, t_a) \in C$, $1 \leq a \leq 3$ and $i_1, i_2, i_3$ are distinct. We claim that it cannot be that all three graphs in $H^\varepsilon_1(e, P)$ are $\eta_2$-quasirandom with density at least $d_2$ and the 3-graph in $H^\varepsilon_1(e, P)$ is $\eta_3$-quasirandom with density at least $d_3$. For then, since $\gamma \ll \theta$, Theorem 3.5 gives us an octahedron (in fact many!) with parts $\{x_a, x'_a\}$, $1 \leq a \leq 3$ in which $x_a, x'_a$ are an edge of $N_H(x) \cap Y_{i_a, t_a}$, $1 \leq a \leq 3$. However such a structure clearly contains a Fano plane, which contradicts our assumptions.

It follows that at least one quasirandomness condition or density condition fails for each induced tripartite 3-complex $H^\varepsilon_1(e, P)$ with vertex classes in $Z$. Now the choice of partition system $P$ guarantees that at most $\gamma n^3$ edges of $H_1$ belong to induced complexes which fail a quasirandomness condition. Also, at most $d_2 n^3$ edges of $H_1$ belong to an induced complex for which the 3-graph has relative density at most $d_3$, and at most $3d_2^{1/2} n^3$ edges of $H_1$ belong to an induced complex in which one of the 2-graphs has relative density at most $d_2$ (since $n_{ij} \leq d_2^{1/2}$ for every $i, j$). Finally, we can estimate the number of edges of $H$ that do not respect the partition $X = X_1 \cup \cdots \cup X_r$, as $|H \setminus H_1| < n^3/r$. We conclude that $d(H[Z]) < \epsilon$, as required.

4.2. Stability

The remainder of the proof makes no further use of quasirandomness. The next step is to use the sparse set to deduce a stability result, i.e. that a Fano-free 3-graph with minimum codegree about $n/2$ is approximately bipartite.

**Theorem 4.2.** Suppose $0 < 1/n \ll \theta \ll \epsilon \ll 1$, $H$ is a 3-graph on a set $X$ of $n$ vertices that does not contain a Fano plane, and $\delta_2(H) > (1/2 - \theta)n$. Then there is a partition $X = A \cup B$ such that at most $\epsilon n^3$ edges are contained entirely within either $A$ or $B$.

*Random Structures and Algorithms* DOI 10.1002/rsa
Proof: It is convenient to introduce a constant $\phi$ with $\theta \ll \phi \ll \epsilon$. By Theorem 4.1 we can find $A \subseteq X$ with $|A| > (1/2 - \phi)n$ and $d(H[A]) < \phi$, i.e. $|H[A]| < \phi |A|^3 / 6$. Let $B = X \setminus A$. Then

$$3|H[A]| + \sum_{p \in \binom{X}{2}} |N_H(p) \cap B| = \sum_{p \in \binom{X}{2}} |N_H(p)| > \binom{|A|}{2} (1/2 - \theta)n,$$

so

$$\sum_{b \in B} |N_H(b) \cap \binom{A}{2}| = \sum_{p \in \binom{X}{2}} |N_H(p) \cap B| > (1/2 - \theta)n \binom{|A|}{2} - |H[A]^3|/2 > (1/2 - 2\phi)n \binom{|A|}{2}.$$

This implies that $|B| > (1/2 - 2\phi)n$, so both $|A|$ and $|B|$ are $(1/2 \pm 2\phi)n$. Let

$$B_0 = \left\{ b \in B : |N_H(b) \cap \binom{A}{2}| < (1 - \phi^{1/2}) \binom{|A|}{2} \right\}.$$

Then

$$(1/2 - 2\phi)n \binom{|A|}{2} < \sum_{b \in B} |N_H(b) \cap \binom{A}{2}| < |B_0|(1 - \phi^{1/2}) \binom{|A|}{2} + |B \setminus B_0| \binom{|A|}{2}$$

$$= |B| \binom{|A|}{2} - |B_0| \cdot \phi^{1/2} \binom{|A|}{2},$$

so $|B_0| < \phi^{-1/2}(|B| - (1/2 - 2\phi)n) < 3\phi^{1/2}n$.

We claim that there is no edge $e \in H[B \setminus B_0]$. Otherwise we would have

$$\left| \bigcap_{b \in e} N_H(b) \cap \binom{A}{2} \right| > (1 - 3\phi^{1/2}) \binom{|A|}{2}.$$

Then, since $\phi \ll 1$, $\bigcap_{b \in e} N_H(b) \cap \binom{A}{2}$ contains a copy of $K_4$, the complete graph on 4 vertices; we can see this by Turán’s theorem, which gives $\pi(K_4) = 2/3$, or even just a simple averaging argument which gives $\pi(K_4) \leq 5/6$. However, this $K_4$ and $e$ span a copy of the Fano plane, contradiction. Therefore every edge of $B$ contains a point of $B_0$, so $|H[B]| < |B_0||B|^2 < 3\phi^{1/2}n^3$. Since $|H[A]| < \phi |A|^3 / 6$ there are less than $\epsilon n^3$ edges contained entirely within either $A$ or $B$.

4.3. The Exact Codegree Result for the Fano Plane

Finally, we use the stability result above to prove Theorem 1.1, which states that if $n$ is sufficiently large and $H$ is a 3-graph on $n$ vertices with minimum 2-degree at least $n/2$ that does not contain a Fano plane, then $n$ is even and $H$ is a balanced complete bipartite 3-graph.

Proof of Theorem 1.1. Let $\epsilon$ be sufficiently small and $n > n_0(\epsilon)$ sufficiently large. By Theorem 4.2 we have a partition $X = X_0 \cup X_1$ so that $|H[X_0]| + |H[X_1]| < \epsilon n^3$. Choose the partition that minimises $|H[X_0]| + |H[X_1]|$. We will show that this partition satisfies the
conclusion of the theorem. Note first that the same argument used in the proof of Theorem 4.2 shows that $|X_0| > 0$ and $|X_1|$ are $(1/2 \pm 2\epsilon)n$.

Next we show that there is no vertex $x \in X_0$ with degree at least $\epsilon^{1/4}n^{3/2}$ in $H[X_0]$. For suppose there is such a vertex $x$. By choice of partition we have $|N_H(x) \cap \binom{X_0}{2}| \geq |N_H(x) \cap \binom{X_2}{2}| \geq \epsilon^{1/4} n^2$, or we could reduce $|H[X_0]| + |H[X_1]|$ by moving $x$ to $X_1$. We can choose matchings $M = \{x_1^1x_2^1, \ldots, x_n^1x_n^2\}$ in $N_H(x) \cap \binom{X_0}{2}$ and $M' = \{y_1^1y_2^1, \ldots, y_m^1y_m^2\}$ in $N_H(x) \cap \binom{X_2}{2}$, with $m = \epsilon^{1/4}/n/2$; indeed, it is a well-known observation that any maximal matchings will be at least this large, as the vertex set of a maximal matching in a graph covers all of its edges. Now we use an averaging argument to find a pair of edges of $M$ that are not “traversed” by any edges of $H$, in the following sense. We have

$$\sum_{I=\{i,j\} \in \binom{[m]}{2}} \sum_{J=\{i,j\} \in \binom{[2]}{2}} |N_{H[X_0]}(\{x_i^1, x_j^2\})| < 3|H[X_0]| < 3\epsilon n^3,$$

so we can choose $I \in \binom{[m]}{2}$ such that

$$\sum_{J=\{i,j\} \in \binom{[2]}{2}} |N_{H[X_0]}(\{x_i^1, x_j^2\})| < \left(\frac{m}{2}\right)^{-1} 3\epsilon n^3 < 50\epsilon^{1/2}n.$$

Also, for every $1 \leq k \leq m$, $H$ cannot contain all 8 edges $x_i^jx_{i+k}^j$, with $j_1, j_2, j_3$ in $\{1, 2\}$, as then together with $x$ we would have a copy of the Fano plane. By double-counting we deduce that there is some pair $p = x_i^1x_j^2$ such that there are at least $m/4$ vertices $y_k^3$ that do not belong to $N_H(p)$. This gives $|N_H(p)| < 50\epsilon^{1/2}n + |X_0| - \epsilon^{1/4}n/8 < (1/2 - \epsilon)n$, which contradicts our assumptions. We deduce that there is no vertex in $X_0$ with degree at least $\epsilon^{1/4}n^2$ in $H[X_0]$. Similarly there is no vertex in $X_1$ with degree at least $\epsilon^{1/4}n^2$ in $H[X_1]$.

Write $|X_0| = n/2 + t$ and $|X_1| = n/2 - t$, where without loss of generality $0 \leq t \leq 2\epsilon n$. Suppose for a contradiction that there exists $t > 0$ or $t = 0$ and there is an edge in $H[X_0]$ or $H[X_1]$. Note that for every pair $p \in \binom{X_0}{2}$ we have $|N_H(p) \cap X_0| \geq |N_H(p)| - |X_1| \geq t$. Thus we can assume there is at least one edge in $X_0$ (since the case $t = 0$ is symmetrical). Let $Y_t \subseteq X_t$ be minimum size transversals of $H[X_t]$ (i.e., $|e \cap Y_t| \neq 0$ for every $e$ in $H[X_t]$). Then $Y_0 \neq \emptyset$. Also, by the previous paragraph

$$|Y_0|\epsilon^{1/4}n^2 > \sum_{x \in Y_0} |N_{H[X_0]}(x)| \geq |H[X_0]| = \frac{1}{3} \sum_{p \in \binom{X_0}{2}} |N_{H[X_0]}(p)| \geq \frac{t}{3} \binom{|X_0|}{2},$$

so $|Y_0| > \epsilon^{-1/4}t^{-2}n^{-2}(n/2+1) > 50t$, say, since $\epsilon$ is small.

For each edge $e \in X_0$ consider all possible ways to extend it to a copy of the Fano plane using some $F \in \binom{X_1}{3}$. Since $H$ does not contain a Fano plane there is some triple with 2 points in $F$ and one point in $e$ that is not an edge. We count such triples at most $|X_0|^{3/2}$ times, so we get a set of at least $\left(\frac{|X_1|^{3/2}}{2}\right)^{-1} \binom{|X_1|}{3}$ distinct triples. Thus there is some point $x$ in $e$ for which we have a set $M_x$ of at least $\left(\frac{|X_1|^{3/2}}{2}\right)^{-1} \binom{|X_1|}{3} > \frac{1}{20} \binom{|X_1|}{3}$ “missing” triples involving $x$ and a pair in $\binom{X_1}{2}$ not belonging to $N_H(x)$. Now varying $e$ over all edges in $X_0$, we get sets $M_x, x \in T$ of at least $\frac{1}{20} \binom{|X_1|}{3}$ “missing” triples involving $x$, for some transversal $T$. Since $Y_0$ is a minimum size transversal, $M = \cup_{x \in T} M_x$ has size at least $|M| \geq |Y_0| \cdot \frac{1}{20} \binom{|X_1|}{3}$.

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Now, for each pair \( p \in \binom{X_i}{2} \) we have

\[
n/2 \leq |N_H(p)| = |N_H(p) \cap X_0| + |N_H(p) \cap X_1| \leq |X_0| - |N_H(p)| + |N_H(p) \cap X_1|,
\]

so

\[
3|H[X_1]| = \sum_{p \in \binom{X_i}{2}} |N_H(p) \cap X_1| \geq \sum_{p \in \binom{X_i}{2}} (|N_H(Q)| - t) = |M| - t\left(\frac{|X_i|}{2}\right).
\]

Since \(|M| \geq \frac{1}{20}|Y_0|\binom{|X_0|}{2}\) and \(\frac{1}{20}|Y_0| - t > \frac{1}{30}|Y_0|\) we have \(|H[X_1]| > \frac{1}{120}|Y_0|\binom{|X_0|}{2}\). Also \(|Y_1|e^{1/4}n^2 > \sum_{x \in Y_1} |N_H[X_1](x)| > |H[X_1]|\), so \(|Y_1| > e^{-1/4}\frac{1}{120}|Y_0|n^2\left(n^{2/3}\right) > 2|Y_0|\), say, since \(\epsilon\) is small.

Finally, we can apply the argument of the previous paragraph interchanging \(X_0\) and \(X_1\). We get a set \(M'\) of at least \(|Y_1| \cdot \frac{1}{20}\binom{|X_0|}{2}\) distinct triples that are not edges, each having 2 points in \(X_0\) and 1 point in \(X_1\). For each pair \(p \subseteq X_0\) we now have \(|N_H(p) \cap X_0| \geq |N_M'(p)| + t\), so \(|H[X_0]| > \frac{1}{3} \cdot (|Y_1|/20 + t)\binom{|X_0|}{2}\) and \(|Y_0| > e^{-1/4}n^{-2}|H[X_0]| > 2|Y_1|\). This contradiction completes the proof.

\[\Box\]

5. QUASIRANDOM HYPERGRAPHS

We now return to the theory of quasirandom hypergraphs: we will discuss the Gowers approach in full generality, together with some variants that we need for our arguments, the general form of our quasirandom counting lemma, and its application to generalized Turán problems. The essential ideas are already present in the discussion above for graphs and 3-graphs, so we will be fairly brief. We refer the reader to [8] for full details of the Gowers theory, with the proviso that we have also adopted some notation and terminology from [9] and made other changes for consistency of notation. This section is divided into four subsections: the first extends the notation and definitions introduced above to general \(k\)-graphs, the second contains the general forms of the decomposition theorem and counting lemma, the third the general form of our quasirandom counting lemma, and the fourth a generalized form of Theorem 4.1.

5.1. Definitions

Consider a set \(X = X_1 \cup \cdots \cup X_r\) partitioned into \(r\) parts. For brevity call this an \(r\)-partite set. \(A \subseteq X\) is an \(r\)-partite subset if it contains at most one point from each \(X_i\). An \(r\)-partite \(k\)-graph on \(X\) is a \(k\)-graph on \(X\) consisting of \(r\)-partite subsets. An \(r\)-partite \(k\)-complex \(H\) is a collection of \(r\)-partite subsets of \(X\) of size at most \(k\) that forms a simplicial complex, i.e. if \(S \in H\) and \(T \subseteq S\) then \(T \in H\). If \(H\) is a \(k\)-graph then we can generate a \(k\)-complex \(H^{-} = \{T : 3S \in H, T \subseteq S\}\) and a \((k - 1)\)-complex \(H= = \{T : 3S \in H, T \subseteq S\}\) (strict subsets). If \(H\) is an \(r\)-partite \(k\)-complex on \(X = X_1 \cup \cdots \cup X_r\) and \(A \in \binom{X}{\leq k}\) we define an \(|A|\)-complex \(H_A^{-} = \cup_{S \subseteq A} H_S\) and an \((|A| - 1)\)-complex \(H_A= = \cup_{B \subseteq A} H_B\).

The index of an \(r\)-partite subset \(A\) of \(X\) is \(i(A) = \{i \in [r] : A \cap X_i \neq \emptyset\}\). Let \(K_A(X) = \{S \subseteq X : i(S) = i(A)\}\). If \(H\) is a \(k\)-graph or \(k\)-complex we write \(H_A\) for the collection of sets in \(H\) of index \(A\). In particular, if \(S\) is an \(r\)-partite subset of \(X\) and \(A \subseteq i(S)\) then \(S_A = S \cap \cup_{i \in A} X_i\). We also use \(H_A : K_A(X) \to \{0, 1\}\) to denote the characteristic function of this set, i.e. \(H_A(S) = 1\) if \(S \in H_A\) and 0 otherwise.

*Random Structures and Algorithms* DOI 10.1002/rsa
Write $H_A^*$ for the collection of sets $S$ of index $A$ such that all proper subsets of $S$ belong to $H$. (Note that $H_A \subseteq H_A^*$.) We also use $H_A^*$ to denote the characteristic function of this set. The relative $A$-density of $H$ is $d_A(H) = |H_A|/|H_A^*|$. We shorten this to $d_A$ if $H$ is clear from the context. In particular we have $d_0 = 1$, since $H_0 = H_0^* = \emptyset$. We often assume that $H_{[i]} = X_i$, so that $d_{[i]} = 1$ for $1 \leq i \leq r$: in general we can apply a result obtained under this assumption by replacing $X = X_1 \cup \cdots \cup X_r$ with $X' = X'_1 \cup \cdots \cup X'_r$, where $X'_r = H_{[i]}$. We also use an unscripted $d$ to denote (absolute) density, e.g. $d(H_A) = |H_A|/|K_A(X)|$.

Suppose $Y = Y_1 \cup \cdots \cup Y_s$ is another $r$-partite set. We let $\Phi(Y, X)$ denote the set of all $r$-partite maps from $Y$ to $X$: these are maps $\phi : Y \to X$ such that $\phi(Y_i) \subseteq X_i$ for each $i$. If $J$ is a $k$-graph or $k$-complex on $Y$ and $H$ a $k$-graph or $k$-complex on $X$ we say that $\phi$ is a homomorphism if $\phi(J) \subseteq H$.

For each number $i$ let $U_i = \{u_i^0, u_i^1\}$ be a set of size 2 and let $U = \cup_i U_i$. For a set of numbers $A$ let $U_A = \cup_{i \in A} U_i$. The $A$-octahedron on $U$ is $O_A = \{B \subseteq U_A : |B \cap U_i| = 1, i \in A\}$. Suppose that for each $B \in O_A$ we have functions $f_B$ defined on $K_A(X)$. Their Gowers inner product is $\langle \{f_B\}_{B \in O_A} \rangle_{\infty} = \mathbb{E}_{B \in O_A} f_B(\omega(B))$. (Effectively we are averaging over $\omega$ in $\Phi(U_A, X_A)$, but extending the function does not affect the average.) Given a function $f$ defined on $K_A(X)$ let $\text{Oct}(f) = \langle \{f_B\}_{B \in O_A} \rangle_{\infty}$, where $f_B = f$ for every $B \in O_A$. Say that $f$ is $\eta$-quasirandom relative to $H$ if $\text{Oct}(f) \leq \eta \prod_{B \in O_A} d_B(H)$.\footnote{Recall that $O_A^*$ is a $(k-1)$-complex consisting of strict subsets of the sets in $O_A$, and write $d_B(H) = d_B(\{B\})$ for less cumbersome notation. This will be a general rule: a subscript $A$ should be understood as $i(A)$ where appropriate.}

If $H$ is a $k$-complex on $X$ and $A \in {\binom{[r]}{\leq k}}$ and $B \subseteq [r], |B| \leq k$, the balanced function of $H_A$ is $\overline{H_A} = (H_A - d_A)H_A^* = H_A - d_AH_A^*$. We say that $H_A$ is $\eta$-quasirandom if its balanced function $\overline{H_A}$ is $\eta$-quasirandom. If $F$ is a fixed $r$-partite $k$-complex we say that $H$ is $(\epsilon, F, k)$-quasirandom if $H_A$ is $\eta_{|A|}$-quasirandom for every $A \subseteq F$, where the “hidden parameters” $\eta_2, \ldots, \eta_k$ are defined recursively by $\epsilon_k = \epsilon$, and $\eta_s = \frac{1}{k}(\epsilon, \prod_{A \in F, |A| \geq s} d_A)^2$, $\epsilon_{r-1} = \eta_r |F|^{-1} \prod_{s=2}^k 2^{-s}$ for $2 \leq s \leq k$. This terminology is a convenient way of expressing a sufficient condition for approximate counting of homomorphisms from $F$ to $H$. It is sometimes helpful to have a notation including the hidden parameters, thus we say that $H$ is $\overline{\eta}$-quasirandom, where $\overline{\eta} = (\eta_1, \ldots, \eta_k)$, if $H_A$ is $\eta_{|A|}$-quasirandom with respect to $H$ for every $A \in {\binom{[r]}{\leq k}}$ (we can take $\eta_1 = 0$), and we say that $H$ is $\overline{d}$-dense, where $\overline{d} = (d_1, \ldots, d_k)$. If $H_A$ has relative density at least $d_{|A|}$ with respect to $H$ for every $A \in {\binom{[r]}{\leq k}}$ we also say that $H$ is $(F, \overline{d})$-dense, if $H_A$ has relative density at least $d_{|A|}$ with respect to $H$ for every $A \subseteq F$. There is a natural hierarchy $0 \leq \eta_1 \ll d_1 \ll \eta_2 \ll \epsilon_2 \ll d_2 \ll \cdots \ll \eta_k \ll \epsilon_k \ll d_k$, $1/|F| \leq 1$ that arises in applications, and assuming that the functions implicit in the $\ll$-notation decay sufficiently quickly, if $H$ is $\overline{\eta}$-quasirandom and $(F, \overline{d})$-dense then $H_A \leq (\epsilon_{|A|}, F_{A^{\leq k}}, |A|)$-quasirandom for every $A \in {\binom{[r]}{\leq k}}$.

**Comments.** We prefer the term $k$-complex ($k$-chain is used in [8]), as we are dealing with simplicial complexes. We are using the letter $d$ for densities ($\delta$ is used in [8]) to reserve the letter $\delta$ for minimum degree or density. We should emphasise that $d$ with a set subscript indicates relative density with respect to that set, whereas $d$ with no subscript indicates (absolute) density (which will be generally well-approximated by a product of relative densities). For the sake of consistency we always use set subscripts to indicate some kind of restriction to that set. The functions defined on $K_A(X)$ may be regarded as $A$-functions as defined in [8]. The Gowers inner product was formalized in [9] in a slightly restricted context.
and generally in [30], where it was called the cube inner product. Here we are departing from [8], where a counting rather than averaging convention is used, as we find it more convenient not to have to keep track of normalizing factors. Although we do not need these facts, it may aid the reader to know that for \( d \geq 2 \) the operation \( f \mapsto \|f\|_{\ell^d}^2 = \text{Oct}(f)^2 - d \) defines a norm, and there is a Gowers-Cauchy-Schwartz inequality \( |\langle f_B \rangle_{B \in O_A} \|_{\ell^d}^2 | \leq \prod_{B \in O_A} \|f_B\|_{\ell^d}^2 \).

### 5.2. Quasirandom Decomposition and the Counting Lemma

As for 3-graphs, the above notion of quasirandomness admits a decomposition theorem and a counting lemma.

Suppose as before, that \( X = X_1 \cup \cdots \cup X_r \) is an \( r \)-partite set. A partition \( k \)-system \( P \) is a collection of partitions \( P_A \) of \( K_A(X) \) for every \( A \in \binom{X}{k} \). For each \( A \) there is a natural refinement of \( P_A \), called strong equivalence in [8], where sets \( S, S' \) in \( K_A(X) \) are strongly equivalent if \( S_B \) and \( S'_B \) belong to the same class of \( P_B \) for every \( B \subseteq A \).\(^{15}\) Given \( x \in K_{\{r\}}(X) \) the induced complex \( P(x) \) has maximal edges equal to all sets strongly equivalent to some \( x_A = \{x_i : i \in A\} \) with \( |A| = k \). The main case of interest is when \( H \) is a \( k \)-graph or \( k \)-complex and \( P_A = (H_A, K_A(X) \setminus H_A) \).

The following decomposition theorem is a variant of Theorem 7.3 in [8], with two key differences: (i) we have an equitable partition of \( X \), together with an exceptional class, and (ii) we want some flexibility in the choice of the hidden parameters in quasirandomness, rather than the specific choice inherent in the definition of \((\epsilon, F, k)\)-quasirandomness. The proof is very similar to that given in [8], so we will defer a sketch of the necessary modifications that give this version to appendix A.

**Theorem 5.1.** Suppose \( 0 < 1/n \ll d_1 \ll d_2 \ll d_3 \ll \cdots \ll d_{k-1} \ll d_k \ll \epsilon, 1/r, 1/m < 1 \), that \( X = X_1 \cup \cdots \cup X_r \) is an \( r \)-partite set with each \( |X_i| \geq n \), and \( P \) is a partition \( k \)-system on \( X \) such that each such that each \( P_A \) partitions \( K_A(X) \) into at most \( m \) sets. Then there is a partition \( k \)-system \( Q \) refining \( P \) such that

- \( Q_A = P_A \) when \( |A| = k \),
- \( |Q_A| \leq d_{|A|}^{\frac{1}{|A|}} \) when \( |A| < k \),
- \( Q_i = Q_{\{i\}}, 1 \leq i \leq r \) are of the form \( X_{i}^{0,1} \cup \cdots \cup X_i^{0,0,1} \cup X_i^{1,1} \cup \cdots \cup X_i^{m_0} \), where \( X_i^{0,1} = X_i^{0,0,1} \cup \cdots \cup X_i^{0,0,0} \) are exceptional classes of size \( |X_i^{0,0}| < \epsilon |X_i| \), and \( X_i^t \) \( 1 \leq i \leq r, t \neq 0 \) are all of equal size, and
- if \( x = (x_1, \cdots, x_r) \) is chosen uniformly at random from \( K_{\{r\}}(X) \), then with probability at least \( 1 - \epsilon \) the induced complex \( Q(x) \) is \( \overline{\eta} \)-quasirandom and \( \overline{d} \)-dense, for some \( 0 < d_k \leq 1 \).

Next, we have the following generalized counting lemma, the adaptation of Theorem 5.1 in [8] to the setting of \( \overline{\eta} \)-quasirandom \( \overline{d} \)-dense complexes. (The proof is identical.)

**Theorem 5.2.** Suppose \( 0 < 1/n \ll d_1 \ll d_2 \ll d_3 \ll \cdots \ll d_{k-1} \ll d_k \ll \epsilon, d_k, 1/|F| < 1 \), \( H \) and \( F \) are \( r \)-partite \( k \)-complexes on \( X = X_1 \cup \cdots \cup X_r \) and \( Y = Y_1 \cup \cdots \cup Y_r \) respectively, and \( H \) is \( \overline{\eta} \)-quasirandom and \((F, \overline{d})\)-dense. Suppose \( F_0 \) is a

---

\(^{15}\)Recall that set subscripts indicate an appropriate restriction: \( S_B = S \cap \cup_{i \in B} X_i \).
subcomplex of $F$, for each $e \in F$ we have a function $f_e$ on $K_v(X)$ with $f_e = f_e H_e$, and $f_e = H_e$ for all $e \in F \setminus F_0$. Then

$$
\mathbb{E}_{\phi \in \Phi(Y, X)} \left[ \prod_{e \in F} f_e(\phi(e)) \right] = \prod_{e \in F \setminus F_0} d_e(H) \cdot \mathbb{E}_{\phi \in \Phi(Y, X)} \left[ \prod_{e \in F_0} f_e(\phi(e)) \right] \pm \epsilon \prod_{e \in F} d_e(H).
$$

As for 3-graphs, there are two special cases of Theorem 5.2 that are particularly useful. Firstly, when $F_0 = \emptyset$ we get a counting lemma for the partite homomorphism density of $F$ in $H$:

$$
d_F(H) = \mathbb{E}_{\phi \in \Phi(Y, X)} \left[ \prod_{e \in F} H_e(\phi(e)) \right] = (1 \pm \epsilon) \prod_{e \in F} d_e(H).
$$

The further special case when $F$ is a single edge gives an approximation of absolute densities in terms of relative densities:

$$
d(H_A) = (1 \pm \epsilon) \prod_{B \subseteq A} d_B(H).
$$

Secondly, we see that quasirandomness implies vertex-uniformity (the proof is almost identical to that given for 3-graphs, so we omit it):

**Corollary 5.3.** Suppose $0 < 1/n \ll d_1 \ll \eta_2 \ll d_2 \ll \eta_3 \ll \cdots \ll d_{k-1} \ll \eta_k \ll \epsilon, d_k < 1$ and $H$ is an $\eta$-quasirandom $\bar{d}$-dense $k$-partite $k$-complex on $X = X_1 \cup X_2 \cup X_3$, with $H_i = X_i$, $1 \leq i \leq k$. Suppose $W_i \subseteq X_i$, $1 \leq i \leq r$ and let $H_{[k]}[W]$ be the restriction of $H_{[k]}$ to $W = W_1 \cup \cdots \cup W_k$. Then

$$
\frac{|H_{[k]}[W]|}{|X_1| \cdots |X_k|} = \frac{|H_{[k]}|}{|X_1| \cdots |X_k|} \left( \frac{|W_1| \cdots |W_k|}{|X_1| \cdots |X_k|} \pm \epsilon \right).
$$

### 5.3. The Homomorphism Complex: Quasirandom Counting

Now we will present our quasirandom version of the counting lemma in full generality. Suppose $J$ is a $t$-partite $k$-complex with vertex set $E = E_1 \cup \cdots \cup E_t$ and $G$ is an $t$-partite $k$-complex with vertex set $Y = Y_1 \cup \cdots \cup Y_t$. We define the homomorphism complex $J \to G$, which we also denote $G'$ for the sake of compact notation. The vertex set is $Y' = Y_1' \cup \cdots \cup Y_t'$, where $Y_i'$ is the set of all maps $\phi_i : E_i \to Y_i$. For each $A \subseteq [t]$, the edges $G_A'$ of index $A$ consist of all $|A|$-tuples $(\phi_i)_{i \in A}$ for which the associated $|A|$-partite map $\phi_A : E_A \to Y_A$ is a homomorphism, i.e. $\phi_A(J_A) \subseteq G_A$. Note that $G' = J \to G$ is formally a $t$-complex, but its sets of size bigger than $k$ have a trivial structure, in that they are present exactly when all their subsets of size (at most) $k$ are present.

**Theorem 5.4.** Suppose $0 \leq 1/n \ll \eta_2 \ll \epsilon_2 \ll \eta_2' \ll d_2 \ll \eta_3 \ll \epsilon_3 \ll \eta_3' \ll \cdots \ll \eta_k \ll \epsilon_k \ll \eta_k' \ll 1/t, d_k < 1$, that $J$ is a $t$-partite $k$-complex on $E = E_1 \cup \cdots \cup E_t$ and $G$ is an $\eta$-quasirandom $(\bar{d}, \bar{d})$-dense $t$-partite $k$-complex on $Y = Y_1 \cup \cdots \cup Y_t$, where $G_i = Y_i$, $1 \leq i \leq t$. Then $G' = J \to G$ is an $\eta'$-quasirandom, where $\eta' = (\eta_2', \ldots, \eta_t')$ with $\eta_i' = 0$ for $k < i \leq t$. Furthermore, the densities may be estimated as $d_A(G') = (1 \pm \epsilon) d_A(G)^{|A|}$ for $|A| \leq k$ and $d_A(G') = 1$ for $|A| > k$. 

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Proof: The argument is very similar to that used in the proof of Theorem 3.5. We write \( d_A = d_A(G) \) and \( d'_A = d_A(G') \) for \( A \subseteq [t], |A| \geq 2 \). Fix \( A \subseteq [t], |A| \geq 2 \). The counting lemma gives

\[
\frac{d'_A}{d_A} = \sum_{S \in J_A} \frac{\prod_{\phi \in \Phi(E,Y)} (\phi(J_A^S) \subseteq G)}{\prod_{\phi \in \Phi(E,Y)} (\phi(J_A^S) \subseteq G)} = (1 \pm \epsilon_A/3) \prod_{S \in J_A} \frac{d_S}{d_A} = (1 \pm \epsilon_A) d'_{A^*}.
\]

Next we need to estimate \( \text{Oct}(\overline{G}_A^t) \). Write \( \overline{G}_A^t = g_0' - g_1' \), where \( g_0' = G_{A^*}^t \) and \( g_1' = d_A G_{A^*}^t \). We can assume that \( |A| \leq k \); otherwise we have \( G_{A^*}^t = d_A G_{A^*}^t \), \( d_A = 1 \) and \( \overline{G}_A = 0 \). Then

\[
\text{Oct}(\overline{G}_A^t) = \sum_j (-1)^j \sum_{\phi \in \Phi(U,Y)} \prod_{B \in O_A} \overline{G}_A^{t}(\phi(B)) = \sum_j \prod_{B \in O_A} \overline{G}_A^{t}(\phi(B)) \prod_{j \in J_A} g_{jB}^t(\phi(B)),
\]

where we recall that \( O_A \) is the \( A \)-octahedron on \( U \), \( j \) ranges over all \([0,1]\)-sequences \((j_B)_{B \in O_A}\), and \( \sum_j = \sum_{B \in O_A} j_B \).

We interpret each summand as counting homomorphisms to functions from \( J_A \), where we define the clone \( J \) of \( J \) as the following \( t \)-partite \( k \)-complex on \( EE = EE_1 \cup \cdots \cup EE_j \). The parts \( EE_j = E_j^0 \cup E_j^1 \), \( 1 \leq i \leq t \). Each point \( x \) has two copies \( x_0, x_1 \) of each point \( x \). The edges of \( J \) consist of \( 2^{|A|} \) copies of \( J \), indexed by \( O_A \), thus \( J = \bigcup_{B \in O_A} J_B^t \), writing \( J_B^t = \{ (x_B^i)_{i \in A} : x_B^i \in J_A \} \), where \( x_B^i \) is \( x_0 \) if \( B \cap U_i = \{ u^0 \} \) or \( x_1 \) if \( B \cap U_i = \{ u^1 \} \). Given a \( t \)-partite map \( \omega \) from \( U \) to \( Y \), we obtain \( t \)-partite maps \( \phi^B = \omega(B), B \in O_A \) acting on various copies of \( E \) in \( EE \), which together give a \( t \)-partite map \( \phi \) from \( EE \) to \( Y \) defined by \( \phi(x_i^j) = \omega(u_i^j(x)) \). Now we can rewrite the summands above using

\[
\prod_{B \in O_A} \overline{G}_A^{t}(\phi(B)) = \prod_{S \in J_A} \sum_{T \in J_A^S} \prod_{j \in T} g_{jB}^t(\phi(T)) = \sum_{T \in J_A} d_T^t(G),
\]

where we write \( B(T) \) for that \( B \in O_A \) such that \( Y \in J_B^t \), \( g_0 = G_A \) (the characteristic function) and \( g_1 = d_A G_{A^*}^t \).

Now we can estimate the contribution from the summand corresponding to \( j = (j_B)_{B \in O_A} \) using the counting lemma for homomorphisms to \( G \) from the complex \( K_j = \bigcup_{B \in O_A} J_B^t \). We have

\[
\sum_j (-1)^j \sum_{\phi \in \Phi(U,Y)} \prod_{B \in O_A} g_{jB}^t(\phi(B)) = \sum_j (-1)^j \sum_{\phi \in \Phi(U,Y)} \prod_{T \in J_A} g_{jB}(\phi(T)) = \sum_j (-1)^j \sum_{T \in J_A} d_T^t(G).
\]

Therefore

\[
\text{Oct}(\overline{G}_A^t) = \sum_j (-1)^j \sum_{T \in J_A} d_T^t(G) = \sum_j (-1)^j \sum_{B \in O_A} d_T^t(G) = \sum_j (-1)^j \sum_{B \in O_A} d_T^t(G) < n_A \prod_{B \in O_A} d_T^t(G).
\]

Since \( \epsilon_A \ll n_A \). This proves that \( G_{A^*}^t \) is \( n_A \)-quasirandom, so we are done. 

Random Structures and Algorithms DOI 10.1002/rsa
5.4. An Application to Generalized Turán Problems

In this subsection we apply the quasirandom counting lemma to derive information about generalized Turán problems for configurations that have the following particular structure. Suppose $F$ is a $k$-graph and $s$ is an integer. The $s$-blowup $F(s)$ of $F$ is defined as follows. For each vertex $x$ of $F$ there are vertices $x_1, \ldots, x_s$ of $F(s)$. For each edge $x^1 \cdots x^t$ of $F$ we have all $s^t$ edges $x_1^{i_1} \cdots x_s^{i_s}$ with $1 \leq i_1, \ldots, i_s \leq s$ in $F(s)$. Note that $F(2) = FF$ is the clone of $F$ as defined in the previous section. For $1 \leq s \leq k - 1$ we define the $s$-augmentation $F^+$ of $F$ as the $k$-graph obtained from $F(s)$ by adding a set $V^+$ of $k - s$ new vertices and all edges $V^+ \cup \{x_1, \ldots, x_s\}$, with $x \in V(F)$. The following theorem is a generalization of Theorem 4.1 in two respects: (i) the relationship between a single edge and the Fano plane is replaced by the relationship between $F$ and $F^+$, (ii) the assumptions have been relaxed to “$\theta$-approximate” assumptions.

Theorem 5.5. Suppose $0 < 1/n \ll \theta \ll \epsilon \ll 1/k$, $1 \leq s \leq k - 1$, $0 < \delta < 1$, $F$ is a $k$-graph, $H$ is a $k$-graph on a set $X$ of $n$ vertices, $d_{F^+}(H) < \theta$ and the $s$-graph $G = \{S \in \binom{X}{s} : |N_H(S)| < \delta|\binom{X}{s}|\}$ has density $d(G) < \theta$. Then there is a subset $Z \subseteq X$ with $|Z| > (\delta - \epsilon)n$ so that $d_F(H[Z]) < \epsilon$.

Proof of Theorem 5.5. We operate with a constant hierarchy $0 < 1/n \ll \theta \ll \alpha \ll d_1 \ll \eta_2 \ll d_2 \ll \eta_3 \ll \cdots \ll d_k \ll \eta_k \ll \gamma, d_k \ll \omega \ll \epsilon \ll 1$. Set $r = [\omega^{-2}]$ and let $X = X_1 \cup \cdots \cup X_r$ be an arbitrary partition with $||X_i| - n/r| < 1$ for $1 \leq i \leq r$. Let $H_1$ be the edges of $H$ that respect this partition, i.e. have at most one point in each $X_i$. Consider the partition system $P$ naturally associated with $H^r_1$, in which $P_A$ is the partition $K_A(X) = (H^r_1)_A \cup K_A(X) \setminus (H^r_1)_A$. Applying Theorem 5.1, using $\gamma$ instead of $\epsilon$, we obtain a partition system $Q$ refining $P$ such that

- $Q_A = P_A$ when $|A| = k$,
- $|Q_A| \leq d_A^{-1/2}$ when $|A| < k$,
- $Q_i = Q_{(i)}$, $1 \leq i \leq r$ are of the form $X^{0,1}_i \cup \cdots \cup X^{0,mi}_i \cup X^1_i \cup \cdots \cup X^{m_i}_i$, where $X^{0,1}_i = X^{0,1}_i \cup \cdots \cup X^{0,mi}_i$ are exceptional classes of size $|X^{0}_i| < \gamma|X_i|$, and the $X^1_i$, $1 \leq i \leq r$, are all of equal size, say $m_i$, and
- if $x = (x_1, \ldots, x_r)$ is chosen uniformly at random from $K_{r}(X)$, then with probability at least $1 - \gamma$ the induced complex $Q(x)$ is $\overline{\eta}$-quasirandom and $\overline{d}$-dense, for some $0 < d_k < 1$.

Consider the set of all $s$-tuples contained in the vertex classes of $Q$: denote it $Y = \cup_{(i,t) \in I} Y_{i.t}$, where $I = \{(i,t) : 1 \leq i \leq r, 1 \leq t \leq n_i\}$ and $Y_{i.t} = \binom{X_i}{t}$. Let $W = \binom{X}{k-s}$ and $W' = \{R \in W : |N_H(R) \cap Y| > (\delta - \omega)|Y|\}$. Double counting gives

\[
(|Y| - \theta n'/s!)\delta\binom{n-s}{k-s} = \sum_{S \in Y \setminus G} d_H(S) = \sum_{R \in W} |N_H(R) \cap Y| < |W'||Y| + |W'|(\delta - \omega)|Y|,
\]

which gives $|W'| > \frac{1}{\omega}n'!$. Write $C_R = \{(i,t) \in I : |N_H(R) \cap Y_{i.t}| > \omega(n'_s!\rangle$. Since $\alpha \ll d_1$ there are at most $2^{2|I|} \leq \alpha^{-1/2}$ possibilities for $C_R$, so there is some $C$ that occurs as
$C_R$ for at least $\alpha(\binom{n}{k})$ of the $R$ in $W'$. For now we focus on one such $R$. Then, since $|X^i_j| = m$ for every $(i,t) \in I$,

$$(\delta - \omega)|I| \binom{m}{s} < |N_R(R) \cap Y| < \sum_{(i,t) \in C} \binom{m}{s} + \sum_{(i,t) \in F \setminus C} \omega \binom{m}{s},$$

so $|C| > (\delta - 2\omega)|I|$. Let $Z = \cup_{(i,t) \in C} X^i_t$. Then $|Z| = |C|m > (\delta - 3\omega)n$.

Suppose $V(F) = \{v^1, \ldots, v^t\}$ for some $t$. We think of $F^\pm$ as a $t$-partite $k$-complex with one vertex in each part. For any (ordered) $t$-tuple of vertices $x = (x_1, \ldots, x_t)$ with $x_a \in X^i_{i_a}$, $(i_a, t_a) \in C$, $i_a$ distinct, $1 \leq a \leq t$ we consider the induced $t$-partite $k$-complex $Q(x)$, i.e. the union of the induced complexes corresponding to each $k$-subset of $x$. For convenient notation we temporarily identify $(i_1, \ldots, i_t)$ with $|t| = \{1, \ldots, t\}$. We let $H(x) \subseteq Q(x)$ be obtained by keeping those $k$-tuples corresponding to edges of $H$, i.e. $H(x)_A = Q(x)_A$ if $Q(x)_A \subseteq H_A$ and $H(x)_A = \emptyset$ if $Q(x)_A \subseteq X_A \setminus H_A$, for $A = \{i_a\}$ of the $\overline{\eta}$-quasirandom and $(F^\pm, \bar{\eta})$-dense. Otherwise, we can calculate as follows that there is too high a probability that a random map $\phi : V(F^+) \to X$ is a homomorphism from $F^+$ to $H$. Using the definition of $F^+$, we write the event $\{\phi(F^+) \subseteq H\}$ as $E_1 \cap E_2 \cap E_3 \cap E_4$, where $E_1$ is the event that the additional vertex set $V^+$ is mapped to some $R$ in $W'$ with $C_R = C$, $E_2$ is the event that $V(F(s))$ is mapped in $r$-partite fashion to the classes $X^i_t$, $1 \leq a \leq t$, $E_3$ is the event that each $s$-tuple $v^i_1 \cdots v^i_s$ is mapped to $N_H(R) \cap Y_{i_a,i_t}$ for $1 \leq a \leq t$, and $E_4$ is the event $\{\phi(F(s)) \subseteq H\}$. Now $\Pr(E_1) > \alpha/2$ by definition of $C$, $\Pr(E_2|E_1) = (m/n)^{\alpha} > \alpha$ and $\Pr(E_3|E_1 \cap E_2) > (\omega/2)^t$ by definition of $C_R$. Also, Theorem 5.4 and Corollary 5.3 applied to $J = F(s) \leq$ give

$$\Pr(E_4|E_1 \cap E_2 \cap E_3) = (1 + \epsilon) \prod_{A \in F(s)} d_A(H(x)) > \frac{1}{2} \prod_{A \in F} d^{ZA}_{|A|} > \alpha.$$ 

We deduce that $\Pr[\phi(F^+) \subseteq H] > \alpha/2 \cdot \alpha \cdot (\omega/2)^t \cdot \alpha > \delta$, which contradicts our assumptions. Therefore $H(x)$ is not both $\overline{\eta}$-quasirandom and $(F^\pm, \bar{\eta})$-dense.

Now we want to estimate $d_F(H[Z]) = \Pr_{f : E \to Z} (f(F) \subseteq H[Z])$. The event $f(F) \subseteq H[Z]$ is contained in the event $B_1 \cup B_2 \cup B_3$ where $B_1$ is the event that $f$ is not a partite map (i.e. two vertices of $F$ are mapped to the same $X_i$), $B_2$ is the event that $f$ is a partite map but $F$ is mapped to $H(x)$ for some $t$-tuple $x$ that is not $\overline{\eta}$-quasirandom, and $B_3$ is the event that neither $B_1$ or $B_2$ holds and $f(F) \subseteq H$. Then $\Pr(B_1) < (\frac{\omega}{r})/r < \omega$. Also, $Q$ was chosen so that at most $\gamma n^k$ edges of $H_1$ belong to induced complexes which fail a quasirandomness condition, so $\Pr(B_2) < |E(F)|/\gamma < \omega$. To estimate $\Pr(B_3)$ we note that on this event $f$ maps $V(F)$ to some $H(x)$ that is $\overline{\eta}$-quasirandom, so by the previous paragraph it is not $(F^\pm, \bar{\eta})$-dense. Since $|Q_A| \leq d^{-1/2}_{|A|}$ when $|A| < k$, for any $a < k$, at most $(\frac{n}{r})^d d^{1/2}_{a} n^k$ edges of $H_1$ belong to an induced complex in which some $a$-graph has relative density at most $d_{a}$. Also, at most $d_{k} n^k$ edges of $H_1$ belong to an induced complex in which the $k$-graph has relative density at most $d_{k}$. Summing these contributions, we deduce that $\Pr(B_3) < \omega$. Finally we have $d_F(H[Z]) < \Pr(B_1) + \Pr(B_2) + \Pr(B_3) < 3\omega < \epsilon$, as required.

**Remark.** As a further extension, note that we could have taken any $I' \subseteq I$ and applied the argument with $Y' = \cup_{(i,t) \in I'} Y_{i_t}$ instead of $Y$. Thus we find some $C' \subseteq I'$ with $|C'| > (\delta - \epsilon) |I'|$ such that $Z' = \cup_{(i,t) \in C'} X^i_t$ satisfies $|C'| m > |I'|(\delta - \epsilon) n$ and $d_F(H[Z']) < \epsilon$.
6. CODEGREE PROBLEMS FOR PROJECTIVE PLANES

In this section we apply Theorem 5.5 to codegree problems for projective planes. The first subsection contains definitions and a summary of previous results on codegree problems for projective planes obtained in [15]. The second subsection generalises the approach used for the Fano plane to obtain strong structural information for the general problem, which already determines the codegree threshold for planes over a field of odd size up to an additive constant (the first part of Theorem 1.2). In the third subsection we complete the proof of Theorem 1.2 with a more detailed analysis for the plane over \( \mathbb{F}_3 \) in terms of its blocking sets: we see that the true codegree threshold in this case lies strictly between the natural upper and lower bounds found in the first subsection. The last subsection deals with the plane over \( \mathbb{F}_4 \), where we demonstrate a surprisingly different behavior from \( \mathbb{F}_2 \) and fields of odd size.

6.1. Definitions and Previous Results

First we give some definitions. Let \( \mathbb{F}_q \) be the field with \( q \) elements, for any prime power \( q \). The projective geometry \( PG_m(q) \) of dimension \( m \) over \( \mathbb{F}_q \) is a \((q^m + 1)\)-graph with vertex set equal to the one-dimensional subspaces of \( \mathbb{F}_q^{m+1} \) and edges corresponding to the two-dimensional subspaces of \( \mathbb{F}_q^{m+1} \), in that for each two-dimensional subspace, the set of one-dimensional subspaces that it contains is an edge of the hypergraph \( PG_m(q) \). A vertex of \( PG_m(q) \) can be described by projective co-ordinates as \((x_1 : \cdots : x_{m+1})\), where \((x_1, \ldots, x_{m+1})\) is any non-zero vector of \( \mathbb{F}_q^{m+1} \) and \((x_1 : \cdots : x_{m+1})\) denotes the one-dimensional subspace that it generates.

The main result of [15] is the following upper bound on the codegree density for general projective geometries, which is tight in many cases.

\[
\pi_q(PG_m(q)) \leq 1 - \frac{1}{m}. \quad \text{Equality holds whenever } m = 2 \text{ and } q \text{ is } 2 \text{ or odd, and whenever } m = 3 \text{ and } q \text{ is } 2 \text{ or } 3.
\]

The results in [15] can be summarized by the following table, in which the entry in the cell indexed by row \( m \) and column \( q \) is either a number indicating the exact value of \( \pi_q(PG_m(q)) \) or an interval in which \( \pi_q(PG_m(q)) \) lies.

<table>
<thead>
<tr>
<th>( m \backslash q )</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>( 2^t, t \geq 3 )</th>
<th>( p^t, p \text{ odd} )</th>
</tr>
</thead>
<tbody>
<tr>
<td>2</td>
<td>1/2</td>
<td>1/2</td>
<td>[1/3, 1/2]</td>
<td>[0, 1/2]</td>
<td>1/2</td>
</tr>
<tr>
<td>3</td>
<td>2/3</td>
<td>2/3</td>
<td>[1/2, 2/3]</td>
<td>[0, 2/3]</td>
<td>[1/2, 2/3]</td>
</tr>
<tr>
<td>( m \geq 5 )</td>
<td>[3/4, 1 - 1/m]</td>
<td>[2/3, 1 - 1/m]</td>
<td>[1/2, 1 - 1/m]</td>
<td>[0, 1 - 1/m]</td>
<td>[1/2, 1 - 1/m]</td>
</tr>
</tbody>
</table>

6.2. General Structure

A useful property of projective geometries proved in [15] is that \( PG_m(q) \subseteq PG_{m-1}(q)^+e \); in particular \( PG_2(q) \subseteq e^{+q} \), where \( e \) is a single \((q+1)\)-edge. Thus we may specialise Theorem 5.5 as follows.

Random Structures and Algorithms DOI 10.1002/rsa
Theorem 6.2. Suppose $q$ is a prime power, $0 < 1/n \ll \theta \ll \epsilon \ll 1/q$, $0 < \delta < 1$, $H$ is a $(q+1)$-graph on a set $X$ of $n$ vertices, $d_{PG_2(q)}(H) < \theta$ and the $q$-graph $G = \{Q \in \binom{X}{q} : |N_H(Q)| < \delta n\}$ has density $d(G) < \theta$. Then there is a subset $Z \subseteq X$ with $|Z| > (\delta - \epsilon)n$ so that $d(H[Z]) < \epsilon$.

Similarly to Theorem 4.2, we can use this to deduce a stability result, describing the approximate structure of a $(q+1)$-graph that does not contain $PG_2(q)$ and has most of its $q$-degrees at least $n/2 - o(n)$.$^{16}$

Theorem 6.3. Suppose $q$ is a prime power, $0 < 1/n \ll \theta \ll \epsilon \ll 1$, $H$ is a $(q+1)$-graph on a set $X$ of $n$ vertices, $d_{PG_2(q)}(H) < \theta$ and the $q$-graph $G = \{Q \in \binom{X}{q} : |N_H(Q)| < (1/2 - \theta)n\}$ has density $d(G) < \theta$. Then there is a partition $X = A \cup B$ such that at most $\epsilon n^{q+1}$ edges are contained entirely within $A$ or within $B$.

Before giving the proof we remark that any $k$-graph $H$ with density $d(H) > 1 - \binom{n}{k}^{-1}$ contains $K_m^k$, the complete $k$-graph on $m$ vertices. Indeed, by averaging $H$ contains a set $M$ of $m$ vertices with $d(H[M]) \geq d(H)$, and then every $k$-tuple in $M$ must be an edge. (In other words we are using the easy bound $\pi(K_m^k) \leq 1 - \binom{n}{k}^{-1}$ on the Turán density, which is far from being best possible, but suffices for our purposes here.)

Proof: Introduce a hierarchy of constants $\epsilon = \theta_1 \gg \theta_2 \gg \theta_3 \gg \theta_4 = \theta$. By Theorem 6.2 we can find $A \subseteq X$ with $|A| > (1/2 - \theta_2)n$ and $d(H[A]) < \theta_2$, i.e. $|H[A]| < \theta_2 |A|^{q+1}/(q+1)!$. Let $B = X \setminus A$. Now

\[(q+1)|H[A]| + \sum_{Q \in \binom{A}{q}} |N_H(Q) \cap B| = \sum_{Q \in \binom{A}{q}} |N_H(Q)|
\]
\[> \left|\binom{A}{q}\right| \cdot G(1/2 - \theta)n
\]
\[> (1/2 - \theta_2)n \binom{|A|}{q}
\]

so

\[\sum_{b \in B} \left|N_H(b) \cap \binom{A}{q}\right| = \sum_{Q \in \binom{A}{q}} |N_H(Q) \cap B|
\]
\[> (1/2 - \theta_2)n \binom{|A|}{q} - \theta_2 |A|^{q+1}/q!
\]
\[> (1/2 - \theta_3)n \binom{|A|}{q}.
\]

This implies that $|B| > (1/2 - \theta_3)n$, so $|A|$ and $|B|$ are $(1/2 \pm \theta_3)n$. Let

\[B_0 = \left\{ b \in B : \left|N_H(b) \cap \binom{A}{q}\right| < \left(1 - \theta_3^{1/2}\right) \binom{|A|}{q}\right\}.
\]

$^{16}$This weakened assumption may strike the reader as strange at first sight, as in a straightforward application of the stability method one has information about all $q$-degrees, but its importance will become clearer in the last subsection.

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Then

\[(1/2 - \theta_3)n\left(\frac{|A|}{q}\right) < \sum_{b \in B} \left| N_H(b) \cap \binom{A}{q} \right| < |B_0|(1 - \theta_3^{1/2})\left(\frac{|A|}{q}\right) + |B \setminus B_0|\left(\frac{|A|}{q}\right)\]

\[= |B|\left(\frac{|A|}{q}\right) - |B_0| \cdot \theta_3^{1/2}\left(\frac{|A|}{q}\right),\]

so

\[|B_0| < \theta_3^{-1/2}(|B| - (1/2 - \theta_3)n) < 2\theta_3^{1/2}n.\]

We claim that there is no edge \(e \in H[B \setminus B_0]\). If there were we would have \(|\bigcap_{b \in e} N_H(b) \cap \binom{A}{q}| > (1 - (q + 1)\theta_3^{1/2})|A|\). But then, by the remark before the proof \(|\bigcap_{b \in e} N_H(b) \cap \binom{A}{q}| contains a \(K^2_q\), and together with \(e\) we have a copy of \(PG_2(q)\). Therefore every edge of \(B\) contains a point of \(B_0\), so \(|H[B]| \leq |B_0|(|A|/q^2) < 2\theta_3^{1/2}n((1/2 + \theta_3)n)^q/q! < \frac{1}{2}e^{nq+1}\). Since \(|H[A]| < |e|A|^{q+1}/(q + 1)\) we are done.

Finally, just as in the proof of Theorem 1.1, we use the previous stability result to obtain very precise information about a \((q + 1)\)-graph on a set \(X\) of \(n > n_0\) vertices with minimum \(q\)-degree \(\delta_q(H) \geq n/2\) that does not contain \(PG_2(q)\).

**Theorem 6.4.** For any prime power \(q\) there is a number \(n_0\) so that if \(H\) is a \((q + 1)\)-graph on a set \(X\) of \(n > n_0\) vertices with minimum \(q\)-degree \(\delta_q(H) \geq n/2\) that does not contain \(PG_2(q)\) then \(n\) is even, there is a partition \(X = X_0 \cup X_1\) where \(|X_0| = |X_1| = n/2\) and \(H[X_0] = H[X_1] = \emptyset\).

**Proof:** Let \(\epsilon\) be sufficiently small and \(n > n_0(\epsilon, q)\) sufficiently large. By Theorem 6.3 we have a partition \(X = X_0 \cup X_1\) so that \(|H[X_0]| + |H[X_1]| < \epsilon n^{q+1}\). Choose the partition that minimises \(|H[X_0]| + |H[X_1]|\). We will show that this partition satisfies the conclusion of the theorem. Note first that the same argument used in the proof of Theorem 6.3 shows that \(|X_0|\) and \(|X_1|\) are \((1/2 \pm 2\epsilon)n\).

First we show that there is no vertex \(x \in X_0\) with degree at least \(\epsilon^{1/2}n^{q}\) in \(H[X_0]\). For suppose there is such a vertex \(x\). By choice of partition we have \(|N_H(x) \cap \binom{X_0}{q}| \geq |N_H(x) \cap \binom{X_1}{q}| \geq \epsilon^{1/2}n^{q}\), or we could reduce \(|H[X_0]| + |H[X_1]|\) by moving \(x\) to \(X_1\). We can choose matchings \(M = \{x_1^1 \cdots x_1^q, \ldots, x_m^1 \cdots x_m^q\}\) in \(N_H(x) \cap \binom{X_0}{q}\) and \(M' = \{y_1^1 \cdots y_1^q, \ldots, y_m^1 \cdots y_m^q\}\) in \(N_H(x) \cap \binom{X_1}{q}\), with \(m = \epsilon^{1/2}n^q/q\) (as in the case \(q = 2\) we are using the well-known observation that any maximal matchings will be at least this large). Now

\[\sum_{I = \{i_1, \ldots, i_q\} \in \binom{[m]}{q}} \sum_{J = \{j_1, \ldots, j_q\} \in \binom{[q]}{q}} \left| N_{H[X_0]}(x_{i_1}^1 \cdots x_{i_q}^q) \right| < (q + 1)|H[X_0]| < (q + 1)e^{n^{q+1}},\]

so we can choose \(I \in \binom{[m]}{q}\) such that

\[\sum_{J = \{j_1, \ldots, j_q\} \in \binom{[q]}{q}} \left| N_{H[X_0]}(x_{i_1}^1 \cdots x_{i_q}^q) \right| < \left(\frac{m}{q}\right)^{-1} (q + 1)e^{n^{q+1}} < 2q^q(q + 1)e^{1/2}n.\]

Also, for every \(1 \leq k \leq m\), \(H\) cannot have all \(q^{q+1}\) edges \(x_{i_1}^1 \cdots x_{i_q}^j \cdot x_{i_q}^{j+1}\), with \(j_1, \ldots, j_{q+1}\) in \([q]\), as then together with \(x\) we have a \(q\)-augmented edge, which contains \(PG_2(q)\). Therefore
there is some $q$-tuple $Q = x_1^1 \cdots x_q^q$ such that there are at least $q^q m$ vertices $y_k^{i+1}$ that do not belong to $N_H(Q)$. This gives $|N_H(Q)| < 2q^q(q+1)^{1/2}n + |X_1|^q - q^q - 1 q^{1/2} n < (1/2 - \varepsilon)n$, which contradicts our assumptions. We deduce that there is no vertex in $X_0$ with degree at least $e^{1/2} n^q$ in $H[X_0]$. Similarly there is no vertex in $X_1$ with degree at least $e^{1/2} n^q$ in $H[X_1]$.

Write $|X_0| = n/2 + t$ and $|X_1| = n/2 - t$, where without loss of generality $0 \leq t \leq 2\varepsilon n$. Suppose for a contradiction that either $t > 0$ or $t = 0$ and there is an edge in $H[X_0]$ or $H[X_1]$. Note that for every $q$-tuple $Q \in (X_q)$ we have $|N_H(Q) \cap X_0| \geq |N_H(Q) - |X_0| > t$. Thus we can assume there is at least one edge in $X_0$ (since the case $t = 0$ is symmetrical). Let $Y_i \subseteq X_i$ be minimum size transversals of $H[X_i]$. Then $Y_0 \neq \emptyset$. Also, by the previous paragraph

$$|Y_0| e^{1/2} n^q > \sum_{x \in Y_0} |N_H(X_0)(x)| \geq |H[X_0]| = \frac{1}{q+1} \sum_{Q \in (X_q)} |N_H(X_0)(Q)| \geq \frac{t}{q+1} |X_0|$$

so $|Y_0| > e^{-1/2} e^{1/2} n^q n^{-q} (n/2 + t) > 2q^t$, say.

For each edge $e \in X_0$ consider all possible ways to extend it to a copy of $PG_2(q)$ using some $F \in (X_q)$. Since $H$ does not contain $PG_2(q)$ there is some $(q+1)$-tuple with $q$ points in $F$ and one point in $e$ that is not an edge. We count each such $(q+1)$-tuple at most $\binom{|X_0| - 1}{q-1}$ distinct $(q+1)$-tuples. Thus there is some point $x$ in $e$ for which we have a set $M_x$ of at least $\frac{1}{q+1} \binom{|X_1| - 1}{q-1} > q^3 q \frac{|X_1|}{q}$ “missing” $(q+1)$-tuples involving $x$ and a $q$-tuple in $(X_q)$ not belonging to $N_H(x)$. Now varying $e$ over all edges in $X_0$ we get a set $M = \cup_x M_x$ of missing $(q+1)$-tuples with $|M| \geq |Y_0| \cdot q^3 q \frac{|X_1|}{q}$ (since $Y_0$ is a minimum size transversal). For each $q$-tuple $Q \in (X_q)$ we have

$$n/2 \leq |N_H(Q)| = |N_H(Q) \cap X_0| + |N_H(Q) \cap X_1| \leq |X_0| - |M|/q + |N_H(Q) \cap X_1|$$

so

$$(q+1)|H[X_1]| = \sum_{Q \in (X_q)} |N_H(Q) \cap X_1| \geq \sum_{Q \in (X_q)} (|N_M(Q)| - t) = |M| - t \binom{|X_1|}{q}.$$  

Since $q^3 |Y_0| - t > \frac{1}{2} q^3 |Y_0|$ we have $|H[X_1]| > \frac{1}{2(q+1)} q^3 |Y_0| \binom{|X_1|}{q}$. Also $|Y_1| e^{1/2} n^q > \sum_{x \in Y_1} |N_H(Y_1)(x)| > |H[X_1]|$, so $|Y_1| > e^{-1/2} e^{1/2} n^q |Y_0| q^{-q} (n/2 - t) > 2|Y_0|$. Finally we can apply the argument of the previous paragraph interchanging $X_0$ and $X_1$. We get a set $M'$ of at least $|Y_1| \cdot q^3 q \frac{|X_0|}{q}$ distinct $(q+1)$-tuples that are not edges, each having $q$ points in $X_0$ and 1 point in $X_1$. For each $q$-tuple $Q \subseteq X_0$ we now have $|N_H(Q) \cap X_0| \geq |N_M(Q) + t| > |H[X_0]| > \frac{1}{2(q+1)} q^3 |Y_0| + t \binom{|X_0|}{q}$ and $|Y_0| > e^{-1/2} e^{1/2} n^q |H[X_0]| > 2|Y_1|$. This contradiction completes the proof.

**Remark.** The argument applies more generally to any $F$ with $PG_2(q) \subseteq F \subseteq e^{+q}$ (where $e = PG_1(q)$ is a single edge.)

For general odd $q$ Theorem 6.4 determines the $q$-degree threshold to find $PG_2(q)$ up to a constant, thus proving the first part of Theorem 1.2. It provides an upper bound $ex_q(n, PG_2(q)) \leq n/2$. On the other hand, it was proved in [15] that there is no copy

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of $PG_2(q)$ in the complete oddly bipartite $(q + 1)$-graph, by which we mean the construction obtained by forming a balanced partition $X = X_0 \cup X_1$ and taking as edges all $(q + 1)$-tuples with an odd number of points in each $X_i$. This gives a lower bound $\text{ex}_q(n, PG_2(q)) \geq \lfloor n/2 \rfloor - q + 1$.

6.3. The Projective Plane Over $\mathbb{F}_3$

To nail down the constant uncertainty in the bound $\text{ex}_q(n, PG_2(q)) \geq \lfloor n/2 \rfloor - q + 1$ for odd $q$ requires analysis of the degrees of $q$-tuples not contained in $X_0$ or $X_1$, which is closely connected to the theory of blocking sets in projective planes (see [28]). This theory is far from complete, but the case $q = 3$ is sufficiently simple to analyse.

Say that $S \subset PG_m(q)$ is a blocking set if $0 < |S \cap L| < |L|$ for every line $L$ of $PG_m(q)$. Note that the complement of a blocking set is also a blocking set, so the existence of a blocking set is equivalent to the existence of a bipartition of $PG_m(q)$. The blocking sets of $PG_2(3)$ may be classified as follows (see [28] or [10]): they all have size 6 or 7, and those of size 6 are of the form $L(x, y) \cup L(y, z) \cup L(x, z) \setminus \{x, y, z\}$, where $x, y, z$ are non-collinear points and we use the notation $L(a, b)$ for the line containing $a$ and $b$. (Those of size 7 are their complements.)

Consider a partition $V(PG_2(3)) = A_0 \cup A_1$ of $PG_2(3)$ into blocking sets, where $A_1 = L(x, y) \cup L(y, z) \cup L(x, z) \setminus \{x, y, z\}$. We refer to the type of a line $L$ as $A_0^0A_1^i$ if $|L \cap A_i| = t_i$, $i = 0, 1$. Then the lines $L(x, y)$, $L(y, z)$ and $L(x, z)$ each have type $A_0^2A_1^1$, and their restrictions to each part form a triangle in $A_0$ and a matching in $A_1$. There are 6 lines of type $A_0^3A_1$, and we note that the 3 points of $A_0$ consist of 1 point from $\{x, y, z\}$ and 2 of the other 4 points.

This leaves 4 lines of type $A_0A_1^4$, for which we note that the point in $A_0$ is not one of $\{x, y, z\}$. In the picture, we have $x = 100$, $y = 010$, $z = 001$, $A_0$ = white discs, $A_1$ = black discs.

Construction. Now we can describe an improvement to the complete oddly bipartite construction. Take a set $X$ of $n$ points, $n$ even, and partition it as $X = X_0 \cup X_1$ with $|X_0| = n/2 + 1$. Choose 2 special points $a, b$ in $X_0$. Form a 4-graph $H$ by taking as edges all 4-tuples which either have 3 points in one $X_i$ and 1 point in $X_{1-i}$, or have 2 points in $X_1$ and 2 points in $X_0$, exactly one of which is $a$ or $b$. Then $H$ does not contain $PG_2(3)$, as the

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edges with 2 points in each part do not contain the triangle with matching neighborhood configuration described in the previous paragraph. Also, the minimum 3-degree of $H$ is $n/2 - 1$.

We will show that this construction is optimal. First we need the following lemma, which is similar to (but does not follow from) a case of a result of Diwan and Sobhan Babu [5]. The proof is a simple but slightly tedious case analysis which we outline here for the reader’s convenience.

**Lemma 6.5.** Suppose $G_1$, $G_2$ and $G_3$ are graphs on the same set $Y$ of at least 8 vertices, so that each has minimum degree at least 2 and there is no “multicolored” matching $e_1$, $e_2$, $e_3$ with $e_i \in G_i$, for each $i$. Then there are two points $a, b$ that meet each edge of each $G_i$.

**Proof.** Suppose that no two points meet each edge of each $G_i$. We claim that there is a matching $M$ of size 3 in $G = G_1 \cup G_2 \cup G_3$. To see this, consider a possible counterexample $G$. Divide into cases according to the connectivity $\kappa(G)$ of $G$. The minimum degree condition implies that each component of $G$ contains a matching of size 1 (i.e., an edge), and even a matching of size 2 unless it is a triangle, so $G$ must be connected. If $\kappa(G) = 1$ and $v$ is a cutvertex then every component of $G \setminus v$ contains an edge, so $G \setminus v$ has exactly two components $C, C'$. Since there is no matching of size edge in $C'$. Next suppose that $\kappa(G) = 2$ and $\{u, v\}$ say $C$, is a star with at least 4 vertices, i.e. its edges all contain some vertex $x$ in $C$. By the minimum degree condition any other vertex $y$ in $C$ is joined to $v$. But now we can find a matching of size 3: take $ya, xz$ for some $z \in C \setminus \{x, y\}$ and an edge in $C'$. Next suppose that $\kappa(G) = 2$ and $\{u, v\}$ is a cutset. By hypothesis $\{u, v\}$ does not meet every edge, so some component $C$ of $G \setminus \{u, v\}$ contains an edge. If there are at least two other components $C'$, $C''$ of $G \setminus \{u, v\}$ then we can extend this to a matching of size 3 with an edge from $C$ to $u$ and an edge from $C''$ to $v$. If there is one other component $C'$ of $G \setminus \{u, v\}$ then we can find a matching of size 3 in which there is an edge from $C$ to $u$ or $v$, from $C'$ to the other of $u$ or $v$, and a third edge in $C$. Finally, suppose $\kappa(G) \geq 3$ and $S$ is a minimum cutset. Since every $x \in S$ has a neighbor in every component of $G \setminus S$ there are exactly 2 components of $G \setminus S$. Then we can find a matching of size 3 using two edges from $S$ to components of $G \setminus S$ and one edge inside a component of $G \setminus S$. In all cases we see that there is a matching $M$ of size 3 in $G = G_1 \cup G_2 \cup G_3$.

If not all three colors $G_i$ are used in $M$ we will show how to increase the number of colors. Suppose first that all three edges are in $G_1$. Pick any new point $x$ and an edge $e$ of $G_2$ that contains it. We can include $e$ and discard whichever edge of our original matching meets it (or any if none does) to obtain a new matching $M'$ of size 3 on which the color $G_2$ also appears (on edge $e$). Now any new point $y$ must be incident to exactly two edges of $G_3$, joining it to the endpoints of $e$ (otherwise we would have a multicolored matching). However, we could take another new point $z$, an edge $e'$ in $G_2$ containing $z$, and include $e'$ in a new matching (discarding an appropriate edge). Now we have either 2 edges of $G_1$ and 1 of $G_2$ (if $e$ was discarded) or 2 edges of $G_2$ and 1 of $G_1$ (otherwise). Either way, the same reasoning as before tells us that every new point $y'$ is incident in $G_1$ to exactly to the endpoints of $f$, where $f \neq e$ is some edge of the new matching, contradicting the fact that it is incident exactly to the endpoints of $e$.

**Theorem 6.6.** If $n$ is sufficiently large and $H$ is a 4-graph on a set $X$ of $n$ vertices that does not contain $PG_2(3)$ then $\delta_3(H) < n/2$.

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Proof of Theorem 6.6. Suppose for a contradiction that \( \delta_3(H) \geq n/2 \). By Theorem 6.4 we have \( n \) even, \( X = X_0 \cup X_1 \) with \( |X_0| = |X_1| = n/2 \) and \( H[X_0] = H[X_1] = \emptyset \). Also, for every 3-tuple contained in one of the \( X_i \), to get minimum 3-degree \( n/2 \) every 4-tuple obtained by adding a point from \( X_{1-i} \) must be an edge. For every pair \( a, b \) in \( X_0 \) we have a graph \( G_{a,b} = N_H(a,b) \cap \binom{X}{2} \) on \( X_1 \) with minimum degree at least 2. These graphs do not contain a “triangle-colored matching,” i.e., a triple \( a, b, c \) in \( X_0 \) and a matching \( e_1, e_2, e_3 \) in \( X_1 \) with \( e_1 \in G_{b,c}, e_2 \in G_{a,c}, e_3 \in G_{a,b} \). For every 4-tuple with 3 points in one \( X_i \) and one point in \( X_{1-i} \) is an edge, so using the description of the blocking sets in \( PG_2(3) \), any triangle-colored matching could be completed to a copy of \( PG_2(3) \), contrary to assumption. By the lemma there must be two points \( a_1, b_1 \) in \( X_1 \) that meet every edge of each of \( G_{b,c} \), \( G_{a,c} \) and \( G_{a,b} \). In fact \( a_1, b_1 \) must meet every edge of \( G_{d,a'} \) for any pair \( a', b' \) in \( X_0 \), as may be seen by applying the previous reasoning in the triangles \( a'b'a' \) and \( b'ab' \) (without loss of generality). Similarly there are two special points \( a_0, b_0 \) in \( X_0 \) that meet every edge with 2 points in each \( X_i \). But now any triple \( a, b, c \) with say \( a, b \) in \( X_0 \setminus \{a_0, b_0\} \) and \( c \) in \( X_1 \setminus \{a_1, b_1\} \) is not contained in any edge with 2 points in each \( X_i \), so has 3-degree equal to \( n/2 - 2 \). This contradiction completes the proof.

6.4. The Projective Plane Over \( \mathbb{F}_4 \)

We conclude by demonstrating a somewhat surprising phenomenon for the projective plane over \( \mathbb{F}_4 \): its codegree density is less than \( 1/2 - c \) for some absolute \( c > 0 \), unlike the cases of \( \mathbb{F}_2 \) and \( \mathbb{F}_q \), \( q \) odd where the codegree density is \( 1/2 \).

Before proving our bound we need some information about the blocking sets of \( PG_2(4) \). A classification was given in [1], but we will just need two specific examples. For the first, suppose more generally that \( q \) is a prime power. Then one blocking set in \( PG_2(q^2) \) is a Baer subplane \( B \), which may be constructed by restricting to those points \( (x_0 : x_1 : x_2) \) that have some representative \( (a_0, a_1, a_2) \) with each \( a_i \) in the base field \( \mathbb{F}_q \). Each line of \( PG_2(q^2) \) contains either 1 or \( q + 1 \) points of \( B \). The intersections of size \( q + 1 \) in \( B \) form the lines of a copy of \( PG_2(q) \). Also, since every pair of lines in \( PG_2(q) \) intersect, the lines of \( PG_2(q^2) \) containing them do not intersect outside of \( B \). For the sake of being more concrete, we remark that this can be described more explicitly in \( PG_2(4) \) using the representation \( PG_2(4) = \{A + x : x \in \mathbb{Z}_{21}\} \), where \( A = \{3, 6, 7, 12, 14\} \). An example of a Baer subplane is \( B = \{ \ x : x \equiv 0 \mod 3 \} \). Dividing by 3 we can represent the lines of \( B \) as \( \{A' + x : x \in \mathbb{Z}_7\} \), where \( A' = \{1, 2, 4\} \): a well-known description of the Fano plane.

We also need to consider the blocking set \( \Lambda = L(x, y) \cup L(x, z) \cup \{w\} \setminus \{y, z\} \), where \( x, y, z \) are 3 non-collinear points, and \( w \in L(y, z) \). Consider the associated 2-coloring \( C_0 = \Lambda, C_1 = V(PG_2(4)) \setminus C_0 \). Recall that the type of an edge \( L \) is \( L[\Lambda \setminus C_0, C_1[\Lambda \setminus C_1] \). There are 2 edges of type \( C_0^2 C_1^2 \), 9 of type \( C_0^1 C_1^2 \), 3 of type \( C_0^2 C_1^3 \) and 7 of type \( C_0^3 C_1^2 \). For the \( C_0^2 C_1^2 \) edges the triples in \( C_0 \) can be described as \( w d_i' \) with \( i = 1, 2, 3 \), where \( L(x,y) = \{ x, y, d_1, d_2, d_3 \} \), and \( L(x,z) = \{ x, z, d_1', d_2', d_3' \} \). For the \( C_0^3 C_1^1 \) edges the triples in \( C_1 \) can be described as \( e_1 e_2 e_3, a e_1', b e_1' \) with \( i = 1, 2, 3 \), where we write \( L(y, z) = \{ y, z, w, a, b, L(x, w) = \{ x, w, e_1, e_2, e_3 \} \), \( L(a, e_1) \cup L(x, y) \cup L(x, z) = \{ e_1 \} \), and \( L(b, e_1) \cup L(x, y) \cup L(x, z) \cup \{ b, e_1 \} = \{ e_1' \} \). Figure 3 shows an incomplete representation of \( PG_2(4) \), with \( C_0 \) represented by black discs and \( C_1 \) by white discs.

The following statement generalises Theorem 1.3, in that it makes an allowance for a small exceptional set of small codegrees.

Random Structures and Algorithms DOI 10.1002/rsa
Theorem 6.7. For any $\epsilon > 0$ sufficiently small there is $\theta > 0$ and $n_0$ so that if $H$ is a 5-graph on a set $X$ of $n > n_0$ vertices and $G = \{Q \in \binom{X}{4} : |N_H(Q)| < (1/2 - \theta)n\}$ satisfies $d(G) < \theta$ then $d_{PG_2(4)}(H) > \epsilon$.

Proof. Let $\epsilon_5$ be a small constant (say $10^{-6}$), introduce a hierarchy of constants $\epsilon_5 \gg \cdots \gg \epsilon_2 \gg \epsilon_1 \gg \epsilon_0 \gg \epsilon \gg \theta$ and suppose $n$ is sufficiently large. Suppose for a contradiction that $d(G) < \theta$ but $d_{PG_2(4)}(H) < \epsilon$. By Theorem 6.3 we can find a partition $X = X_0 \cup X_1$ so that $|X_0|$ and $|X_1|$ are $(1/2 \pm \epsilon_0)n$ and $|H[X_0]| + |H[X_1]| < \epsilon_0 n^5$.

We introduce two 3-graphs for $i = 0, 1$:

$$ J_i = \left\{ e \in \binom{X_i}{3} : \left| N_H(e) \cap \binom{X_{1-i}}{2} \right| \geq \epsilon_1 \left| \binom{X_{1-i}}{2} \right| \right\}. $$

We claim that $d_{PG_2(4)}(J_i) < \epsilon_2$. For suppose otherwise, say that $d_{PG_2(4)}(J_0) > \epsilon_2$, and let $\phi$ be a random map from $V(PG_2(4))$ to $X$, conditioned on the event $E(\phi)$ that a Baer subplane $B$ is mapped to $J_0$, and the other points are mapped to $X_1$. We estimate the probability that $\phi$ embeds $PG_2(4)$ in $H$. There are 2 types of edges to consider: 14 of type $X_0^1 X_1^1$ and 7 of type $X_0^3 X_1^2$. Suppose $L$ is an edge of $PG_2(4)$ that we are attempting to embed with type $X_0^3 X_1^2$. For each point $a$ in $X_0$ let $m(a)$ be the number of 4-tuples $Q \in \binom{X}{4}$ such that $a \cup \{Q\}$ is not an edge of $H$. We have

$$ (1/2 - \theta)n \left( \binom{|X_1|}{4} - \theta \binom{n}{4} \right) < \sum_{Q \in \binom{X}{4}} |N_H(Q)| = \sum_{a \in X_0} |N_H(a) \cap \binom{X_1}{4}| + 5|H[X_1]| $$

$$ < |X_0| \binom{|X_1|}{4} - \sum_{a \in X_0} m(a) + 5\epsilon_0 n^5, $$

so $\sum_{a \in X_0} m(a) < 10\epsilon_0 n^5$. Now

$$ \mathbb{P}(\phi(L) \not\in H) = \sum_{a \in X_0} \mathbb{P}(\phi(L) \cap X_0 = \{a\}) m(a) |X_1|^{-4}, $$

and since there are at most $\binom{|X_0| - 1}{6}/\binom{|X_0|}{6}$ copies of $PG_2(2)$ in $J_0$ that use $a$ we have $\mathbb{P}(\phi(L) \cap X_0 = \{a\}) < \binom{|X_0| - 1}{6}/\binom{|X_0|}{6} = \frac{7}{2^{2|X_0|}}$, so $\mathbb{P}(\phi(L) \not\in H) < \frac{7}{2^{2|X_0|}} \cdot 10\epsilon_0 n^5 |X_1|^{-4} < 5000\epsilon_2^{-1} \epsilon_0$. We
are attempting to embed 14 edges of type $X^1_0X^4_0$, so the probability that any fails is less than $10^5\epsilon^{-1}_2\epsilon_0$. On the other hand, we have conditioned on the event that the Baer subplane $B$ is mapped to $J_0$, and so by definition each of the 7 attempted embeddings of type $X^1_0X^2_0$ is successful with probability at least $\epsilon_1 - O(1/n)$. Furthermore, if $T_1, \ldots, T_7$ are the triples of the Baer subplane then $N_{PG_2(4)}(T_1), \ldots, N_{PG_2(4)}(T_7)$ are disjoint sets, so the events are independent. With probability $1 - O(1/n)$ their images under $\psi$ are disjoint, so we estimate $\mathbb{P}(\phi(PG_2(4)) \subseteq H) > \epsilon^{-1}_1 - O(1/n) - 10^5\epsilon^{-1}_2\epsilon_0 > \epsilon^{1/2}$, say. Then, letting $\psi$ be a random map from $V(PG_2(4))$ to $X$, we have $d_{PG_2(4)}(H) = \mathbb{P}(\psi(PG_2(4)) \subseteq H) \geq \mathbb{P}[E(\psi)]\epsilon^{1/2} > \epsilon_2(1/4)^{14}\epsilon^{1/2} > \epsilon$. This contradiction shows that $d_{PG_2(4)}(J_i) < \epsilon_2$.

Next we introduce 2-graphs for $i = 0, 1$:

$$P_i = \left\{ \{a_i, a'_i\} \in \left( \frac{X_i}{2} \right) \mid |N_J(a_i, a'_i)| < (1/2 - \epsilon_2)|X_i| \right\}.$$  

We claim that $|P_i| < \epsilon_2(\frac{|X_i|}{2})$ for at least one of $i = 0, 1$. For suppose otherwise, and consider a 4-tuple $Q = \{a_0, a'_0, a_1, a'_1\}$, where $\{a_i, a'_i\}$ are chosen uniformly at random from $P_i$, $i = 0, 1$.

We have $\mathbb{E}[\epsilon_{N_H(Q)}] = \mathbb{E}[\epsilon_{N_H(Q) \cap X_0}] + \mathbb{E}[\epsilon_{N_H(Q) \cap X_1}]$.

Now

$$\mathbb{E}[\epsilon_{N_H(Q) \cap X_0}] = \mathbb{E}_{a_0, a'_0} \sum_{a \in X_0} \mathbb{P}_{a, a'_0} \left( a \in N_H(Q) \mid a_0, a'_0 \right)$$

and $\mathbb{P}_{a, a'_0}(a \in N_H(Q) \mid a_0, a'_0) = \mathbb{P}(\{a, a'_0\} \in N_H(a_0, a'_0, a))$. For $a \in N_{J_0}(a_0, a'_0)$ we estimate this probability trivially by 1, but for $a \in X_0 \setminus N_{J_0}(a_0, a'_0)$ we use the definition of $J_0$ and the lower bound on $P_1$ to estimate $\mathbb{P}(\{a, a'_0\} \in N_H(a_0, a'_0, a)) < \epsilon_1\epsilon_2^{-1}$. This gives $\mathbb{E}[\epsilon_{N_H(Q) \cap X_0}] \leq |N_H(a_0, a'_0)| + \epsilon_1\epsilon_2^{-1}|X_0| < (1/2 - \epsilon_2/2)|X_0|$. Similarly we estimate $\mathbb{E}[\epsilon_{N_H(Q) \cap X_1}] < (1/2 - \epsilon_2/2)|X_1|$, so $\mathbb{E}[\epsilon_{N_H(Q)}] < (1/2 - \epsilon_2/2)n$. However, we also have $\mathbb{E}[\epsilon_{N_H(Q)}] > \mathbb{P}(Q \notin G) \cdot (1/2 - \theta)n$, and $\mathbb{P}(Q \notin G) \leq \frac{|G|}{|P_0|^2} < \epsilon_2(\frac{\alpha_0^2}{24})^2 < \epsilon$, so $\mathbb{E}[\epsilon_{N_H(Q)}] > (1 - \epsilon)(1/2 - \theta)n$. This contradiction shows that at least one $P_i$ is small, say $|P_0| < \epsilon_2(\frac{|X_0|}{2})$.

Now we can apply Theorem 6.3 to find a partition $X_0 = A_0 \cup B_0$ where $|A_0|$ and $|B_0|$ are $(1/2 \pm \epsilon_3)|X_0|$ such that at most $\epsilon_3|X_0|^3$ edges of $J_0$ are contained entirely within $A_0$ or within $B_0$. Next we repeat the argument to deduce similar structural information on $J_1$. Let

$$P'_0 = \left\{ \{a_0, a'_0\} \in \left( \frac{A_0}{2} \right) \mid |N_{J_0}(a_0, a'_0) \cap A_0| < 400\epsilon_3|X_0| \right\};$$

$$P'_1 = \left\{ \{a_1, a'_1\} \in \left( \frac{X_1}{2} \right) \mid |N_{J_1}(a_1, a'_1)| < (1/2 - \epsilon_4)|X_1| \right\}.$$

We must have $|P'_0| > \frac{1}{2}(\frac{|A_0|}{2})$, otherwise

$$\epsilon_3|X_0|^3 > |(J_0)_{A_0}| = \frac{1}{2} \sum_{a_0, a'_0 \in A_0} |N_{J_0}(a_0, a'_0) \cap A_0| > \frac{1}{6} \left( \frac{|A_0|}{2} \right) \cdot 400\epsilon_3|X_0|,$$

which is a contradiction. Now we cannot have $|P'_1| > \epsilon_4(\frac{|X_1|}{2})$, as then considering $Q = \{a_0, a'_0, a_1, a'_1\}$, where $\{a_i, a'_i\}$ are chosen uniformly at random from $P'_i$, $i = 0, 1$, we estimate (similarly to before) $\mathbb{E}[\epsilon_{N_H(Q) \cap X_0}] < (1/2 + 500\epsilon_3)|X_0| + \epsilon_1\epsilon_2^{-1}|X_0|$ and $\mathbb{E}[\epsilon_{N_H(Q) \cap X_1}] < (1/2 - \epsilon_4)|X_1| + 10\epsilon_1|X_1|$, so $\mathbb{E}[\epsilon_{N_H(Q)}] < (1/2 - \epsilon_4/2)n$, contradiction. Again, by Theorem

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6.3 we find a partition $X_1 = A_1 \cup B_1$ where $|A_1|$ and $|B_1|$ are $(1/2 \pm \epsilon_5)|X_1|$ such that at most $\epsilon_5|X_1|^3$ edges of $J_1$ are contained entirely within $A_1$ or within $B_1$.

Let $T_i$ count edges of $H$ of type $A_i^2B_i^1$, $i = 0, 1$. We can bound $T_i$ by summing degrees of quadruples $Q = \{a_0, a'_0, a_1, a'_1\}$ with $\{a_i, a'_i\} \in \binom{\Sigma_i}{2}$:

$$(1/2 - \theta)n \left(\binom{|A_0|}{2} \binom{|A_1|}{2} - \theta n^2/4\right) < T_0 + T_1 + 3 \sum_{e \in \binom{\Sigma_0}{3}} \left| N_H(e) \cap \binom{A_1}{2} \right|
+ 3 \sum_{e \in \binom{\Sigma_1}{3}} \left| N_H(e) \cap \binom{A_0}{2} \right|.$$

Now

$$\sum_{e \in \binom{\Sigma_0}{3}} \left| N_H(e) \cap \binom{A_1}{2} \right| < |(J_0)_{A_0}| \binom{|A_1|}{2} + \binom{|A_0|}{3} \cdot \epsilon_1 \binom{|X_1|}{2} < 5\epsilon_3 n \binom{|A_0|}{2} \binom{|A_1|}{2}$$

and

$$\sum_{e \in \binom{\Sigma_1}{3}} \left| N_H(e) \cap \binom{A_0}{2} \right| < |(J_1)_{A_1}| \binom{|A_0|}{2} + \binom{|A_1|}{3} \cdot \epsilon_1 \binom{|X_1|}{2} < 5\epsilon_3 n \binom{|A_0|}{2} \binom{|A_1|}{2},$$

so $T_0 + T_1 > (1/2 - 40\epsilon_4)n\binom{|A_0|}{2} \binom{|A_1|}{2}$. Therefore

$$T_0 > |B_0| \left(\binom{|A_0|}{2} \binom{|A_1|}{2} - \left(|B_0| + |B_1|\right) \binom{|A_1|}{2} \binom{|A_0|}{2} \right) - (1/2 - 40\epsilon_5)n \binom{|A_0|}{2} \binom{|A_1|}{2} \right)$$

By symmetry, similar bounds hold for the number of edges in each case when we specify a triple in one $X_i$ respecting the partition $(A_i, B_i)$ and a pair in $A_{1-i}$ or $B_{1-i}$, i.e. the types

$A_0^2B_0^1A_1^2, A_0^2B_0^1B_1^2, A_0^3B_0^0B_1^2, A_0^1B_0^3A_1^1B_1^0, A_0^1B_1^3A_0^1B_0^0, A_0^1B_0^1A_1^1B_1^0$.

Now we find $PG_2(4)$ using the $\land$ coloring. Let $\phi$ be a random map from $V(PG_2(4))$ to $X$, conditioned on the event $E'(\phi)$ that $w, x$ are in $A_0, Y, Z, e_1, e_2, e_1', e_2''$ are in $A_1$, $d_i, d_i'$ are in $B_0$ and $e_3, a, b$ are in $B_1$ ($i$ ranges from 1 to 3). Note that $M = \min(|A_0|, |A_1|, |B_0|, |B_1|) > (1/4 - 2\epsilon_5)n$. There are 11 attempted embeddings of type $X_i^4X_{1-i}^4$. Recalling that $m(a)$ is the number of 4-tuples $Q \in \binom{\Sigma_i}{4}$ such that $a \cup \{Q\}$ is not an edge of $H$, we estimate that the probability that we fail to embed some such $L$ is at most $11 \sum_{a \in X_{1-i}} m(a) M^{-5} < 10^6 \epsilon_0$. There are 10 attempted embeddings that have the type discussed in the previous paragraph, i.e. one of the types equivalent to $A_1^2B_1^1A_1^3$. Each fails with probability at most 400\epsilon_5, so the probability that any fails is at most 4000\epsilon_5. Now a random map $\psi$ from $V(PG_2(4))$ to $X$ satisfies $E'(\psi)$ with probability at least $(1/5)^{21}$ and so succeeds in embedding $PG_2(4)$ in $H$ with probability at least $(1/5)^{21}/2 > \epsilon$. This contradiction completes the proof.

\textbf{Remark.} The question of what happens for $PG_2(2^s)$ in general is intriguing. It seems plausible that the above approach of going from $PG_2(2)$ to $PG_2(4)$ could be adapted to an inductive argument when $s = 2^r$ is a power of 2. Much of the argument would go through as

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If we stick to PG exact results are desired we also have the problem of recovering the exact structure from assuming this it is still not clear how to recover the global approximate structure of H above: our hypergraph has an approximate bipartition X = X₀ ∪ X₁ and Jᵢ = {e ∈ (xᵢ \ 2²ⁱ⁻¹ + 1) : |Nᵢ(e)| > ε₁(2²ⁱ⁻¹)} |Nᵢ(e)| satisfy d PG₂(2²ⁱ⁻¹, Jᵢ) < ε₂. The step that may fail is finding Q of low degree: our approach used the convenient coincidence of 2 · 2²ⁱ⁻¹ = 2²ⁱ which only occurs for r = 1.

7. FUTURE DIRECTIONS

The basic form of our hypergraph regularity method has been well illustrated by its application to projective planes, which are relatively easy to deal with (for reasons yet to be understood), although even here we cannot give exact answers in all cases, and fields of even size seem particularly strange. However, the quasirandom counting lemma has potential to be a powerful tool in the study of any Turán problem, whether generalized or standard. For example, in the Turán problem for the tetrahedron, if we consider any K₃⁻¹-free 3-graph H and a vertex x then the edges of H cannot be quasirandomly distributed with positive density within the triangles of the neighborhood graph NH(x).

If we restrict attention only to excluding simple k-graphs F (meaning that each pair of edges in F have at most one common point) then the projective geometries in higher dimensions point to one stumbling block that should be overcome in future developments of this method. For example, if we consider a 3-graph H on n vertices with no PG₃(2) and all but o(n²) codegrees at least (2/3 − o(1))n then our results will give a set Z of (2/3 − o(1))n vertices that induce a 3-graph with no Fano plane and all but o(n²) codegrees at least (1/2 − o(1))|Z|, so Z is approximately bipartite by our structure result. This suggests that PG₃(2) should be approximately tripartite, and hence an inductive approach for general m and q showing that a (q + 1)-graph H with no PGₘ(q) and all but o(n²) codegrees at least (1 − 1/m − o(1))n should be approximately m-partite.

A potential approach to filling in the gap is suggested by the remark after Theorem 5.5. If we stick to PG₃(2) for the sake of simplicity, then not only do we have an approximately bipartite subhypergraph of size about (2/3)n, but any set of vertices V obtained by taking some classes of the regularity partition contains some approximately bipartite Z’ of size about (2/3)|V’|. Thus we are faced with the problem of recovering structural information about H from various restrictions, which is perhaps best understood in the context of property testing (see [23] for a hypergraph property testing result and many references to the literature). Although a full investigation of this idea is beyond the scope we have set for this article introducing our basic method, we remark that it should be possible to carry our arguments over to this context via a random reducibility property of quasirandom complexes, i.e. that a random restriction of a quasirandom complex to sets of large constant size should be quasirandom with high probability: a high-level sketch is that the martingale used by Lovász and Szegedy [18] to show concentration of the probability that a random map from a fixed graph F to a random graph G is a homomorphism may be extended to show concentration of the octahedral counting function that appears in the definition of quasirandomness. This will allow us to conclude that if m is a large constant and M is a random m-set of vertices then M contains an approximately bipartite subhypergraph of size about (2/3)m with high probability (say 1 − exp −m²/² for some c > 0). However, even assuming this it is still not clear how to recover the global approximate structure of H. If exact results are desired we also have the problem of recovering the exact structure from the approximate structure. This seems to be quite a different type of question, and so far all
instances of its solution have been of a rather ad hoc character, so it would be interesting to develop some general principles here as well.

APPENDIX: VARIANT FORMS OF THE GOWERS QUASIRANDOMNESS FRAMEWORK

In this appendix we justify the variant forms of the decomposition theorem and counting lemma that we used in the article. We first note that the proof given for Theorem 5.1 in [8] also proves Theorem 5.2 in this article: it only helps to assume stronger inequalities for the $\eta_i$, and our parameter hierarchy is such that we can replace $|\mathcal{J}_0 \setminus \mathcal{J}_1| \epsilon$ in [8] by $\epsilon$ for simpler notation. Then the arguments in Section 6 of [8] go through as written. The proof of Theorem 5.1 is quite similar to the proof of Theorem 7.3 in [8]. We will not reproduce that proof in full, but will outline it to sufficient extent to explain what modifications are needed. First we recall a definition ([8], p. 36):

If $S = \{S_1, \ldots, S_s\}$ and $T = \{T_1, \ldots, T_t\}$ are partitions of the same set $U$, the mean-square density of $S$ with respect to $T$ is

$$\text{msd}_T(S) = \sum_{i=1}^{s} \sum_{j=1}^{t} \frac{|T_j|}{|U|} \left( \frac{|S_i \cap T_j|}{|T_j|} \right)^2.$$

We also recall Lemma 8.1 of [8], which states that if $T'$ is a refinement of $T$ then $\text{msd}_{T'}(S) \geq \text{msd}_T(S)$.

Proof of Theorem 5.1. The proof of Theorem 7.3 in [8] is by means of the following iterative procedure. Suppose we have a partition $k$-system $P$ (which may be the initial iterative or one produced by a some number of iterations: we will not complicate notation with a sequence $P_1, P_2, \ldots$). Consider a partition $k$-system $P^*$ defined by what is called weak equivalence in [8]: $S, S' \in K_k(X)$ are in the same class of $P^*_k$ exactly when $S_B$ and $S'_B$ are in the same class of $P^*_B$ for every proper subset $B \subset A$. It is shown (see [8], pp. 38–39) that if $\mathbb{P}_\epsilon(P(x) \in (\epsilon, j, k)\text{-quasirandom}) < 1 - \epsilon$, then there is $A \in \binom{\mathcal{I}}{\mathcal{I}_{\leq k}}$ and a refining partition $k$-system $Q$ of $P$, so that (a) $Q_B = P_B$ unless $B \subset A$, $|B| = |A| - 1$, (b) if $B \subset A$, $|B| = |A| - 1$ then $Q_B$ is a refinement of $P_B$ where each class of $P_B$ is partitioned into at most $c_A(\epsilon, P)$ further classes, and (c) $\text{msd}_{Q^*_A}(Q_A) \geq \text{msd}_{P^*_A}(P_A) + f_A(\epsilon, P)$. Here $c_A$ and $f_A$ are explicitly defined functions that depend only on $\epsilon$ and $P$, and furthermore the dependence of $c_A$ on $P$ depends only on $|P_B| : B \subset A$ and that of $f_A$ depends only on $|P_B| : |B| \geq |A|$. This property of $f_A$ implies that $f_A(\epsilon, Q) = f_A(\epsilon, P)$. Furthermore, this argument still applies in the context of our proof, i.e., we have the same conclusion if $P(x)$ is $\eta$-quasirandom and $\mathcal{I}$-dense with probability less than $1 - \epsilon$. (Note that the functions $c_A$ and $f_A$ now depend on the functions implicit in the $\ll$-notation for the parameter hierarchy.)

To see that the procedure terminates (with some system of partitions with the required property) introduce a function $\zeta_P$ for the system of partitions $P$, which is defined on $\binom{\mathcal{I}}{\mathcal{I}_{\leq k}}$ by $\zeta_P(A) = \left[ \frac{1 - \text{msd}_{P^*_A}(P_A)}{f_A(\epsilon, P)} \right]$. Choose an ordering $<_P$ of $\binom{\mathcal{I}}{\mathcal{I}_{\leq k}}$ in which $|B| \geq |B'|$ implies that $B < B'$. Order functions $\zeta$ on $\binom{\mathcal{I}}{\mathcal{I}_{\leq k}}$ by $\zeta < \zeta'$ if there is $B \in \binom{\mathcal{I}}{\mathcal{I}_{\leq k}}$ such that $\zeta(B') = \zeta(B)$ for all $B' < B$ and $\zeta(B) < \zeta'(B)$. This is a well-ordering, and the iteration takes the system $P$ to a system $Q$ with $\zeta_Q < \zeta_P$, so the procedure terminates.

To prove our version we introduce further refinements in each step of the procedure. First of all we make the general observation that given any partition $E = E_1 \cup \cdots \cup E_i$ there is an
“equalizing method” to find a partition \( E = E_0' \cup E_1' \cup \cdots \cup E_t' \), for which \( |E_i'| = \lfloor |E|/t^2 \rfloor \), \(|E_0'| < |E|/t\) and every \( E_i' \) with \( i \neq 0 \) is contained in some \( E_j \). The method is to repeatedly and arbitrarily select classes \( E_i' \) within some \( E_j \) that still has size at least \( \lfloor |E|/t^2 \rfloor \), and then remove its elements from consideration in later stages. Thus we are unable to use at most \( |E|/t^2 \) elements from each of the \( t \) original classes \( E_i \), and we put the unused elements together in an exceptional class \( E_0' = E_0'_{1} \cup \cdots \cup E_0'_{t} \) of size at most \( |E|/t \).

We start by using the equalizing method in an initial refinement to transform \( P \) into some \( Q \) with an equitable partition of the vertex set. By first arbitrarily refining \( P \) we can assume that \( |X_i'| < \epsilon |X_i|/2, 1 \leq i \leq r \). Then we repeatedly apply the same refinement procedure as above followed by the equalizing method: at each stage we obtain a new system \( Q \) from the original procedure and then refine it to some equitable \( Q' \). It is clear that the number of classes remains bounded by a function only of \( m, r, k \) and \( \epsilon \). By Lemma 8.1 of [8] we have 
\[
\text{msd}_{Q_A}(Q'_A) \geq \text{msd}_{Q_A}(P_A) + f_A(\epsilon, P).
\]
Also, we still have \( Q_B = P_B \) for all \( B \) with \( |B| \geq |A| \), so \( f_\ast (\epsilon, Q) = f_\ast (\epsilon, P) \) and the iterations terminate as before. The amounts added to the exceptional classes decrease rapidly with each iteration (certainly each is at most half of that at the previous iteration), and as we initially added at most \( \frac{1}{2} \epsilon |X_i| \) exceptional elements of \( X_i \) we end up with at most \( \epsilon |X_i| \).

ACKNOWLEDGMENTS

The author thanks Oleg Pikhurko and Yi Zhao for their discussions about Mubayi’s conjecture and the anonymous referees for their helpful comments.

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