Weirerstrass criterian for existence of minimizers (direct method in code of variations).

Pothorous the & Simber prelimening The A furtion f: X -> Rulted is soil to be lover semi-cationary (loss) if $\forall x_n \rightarrow x$ $f(x) \in Mind f(x_n)$. Ihm(V).) If f: X -> Rulton is lec & X is compact then int f(x) is obtained by some x & X. thoot If J= too re are done. Otherwise let l= int flor (Rul-ol). Let x n be a minimizing cognera Piech e conveying subsequence $x_n \to \infty$. Then lef(\bar{x}) $\leq b$ —if $f(x_n) = l$ so we have = l ($e l \in \mathbb{R}$). Let We say that a sequence of meaning pence P(x) converges weally (or narrowly) to pre P(x) it $\int \varphi(x) \, \varphi_n \left(\frac{1}{2} x - \frac{1}{2} \right) \varphi(x) \, \varphi(x) \, \varphi(x) \qquad \forall \varphi \in \mathcal{Z}_b(x) \quad \text{We write } p_n \Longrightarrow p_n \text{ or } p_n \Longrightarrow p_n \Longrightarrow p_n \Longrightarrow p_n \Longrightarrow p_n \text{ or } p_n \Longrightarrow p_n \Longrightarrow$ X=Q this is equivalent to Fr (a) -> Fr (a) * pult of continuity of Fr. Il Adenily of measures (price Il in P(X) is suit to be tright if to I KEX compact s.t. M: (X/K) < & 4: ex. I'm (Probehorn) Let & be Polish & (I'm) = P(X). (Mm) is relatively compet (i.e., Gran) = (Mm) = (Mm) is relatively compet (i.e., Gran) = (Mm) Many professione pre 3000) (iff) (fr) is tight. $\underbrace{\mathbb{R}}_{\Gamma} \left(\chi \setminus \mathsf{K}_{\Gamma} \right) \leq k_{-} \text{if } \mu_{\Gamma} \left(\chi \setminus \mathsf{K}_{\Gamma} \right) \leq \varepsilon$ by ostnucteur the so (fr.) is do tight. For ecompet $K \subseteq X$ we have $C_0(K) = C_0(K) = C(K)$ so the shad is the space of measures ℓ (m) is a bounded squence so (by Bened-Alaopha since $C_0(K)$ is separable) has a really conveyly subsequence: Man IX -> Mx. Take K. E. & through a ling and argument build one subsequence pum which converges weekly to some vi on Ky. Let p(A):= sor V: (Anky). For e = G(X), fy d(pm-p) = 2 hpli: \frac{1}{4} + Juic.)

Ly (X) = 1. 1 p (x) = 1. -- 211412/ K. "= >" + (> > . Le con cover X by spec bell B, B2, -. of redicin r. Let $G_u = B_1 \cup ... \cup B_k$. Then $\lim_{k \to \infty} \inf_{n} p_n(G_u) = 1$. (4)

In lead, otherwise and prove (Green) = c < 1. Taking subsquare, proven produced of the form (Green) = c < 1. I make gives I=p/x)-1. Salu r= - twike Gu. 420, by (8), 3 kg, Kg, ... int pro (6 mm) >1- 1- 2-m Let $A := \bigcap_{n} G_{nn}^{m}$ then int $p_{n}(A) > 1-\epsilon$ & \overline{A} is complete ϵ totally bounded $(\overline{A}^{c}) \leq p_{n}(A^{c}) = \mu_{n}(U(G, S^{c})) \leq \sum_{i} p_{n}(G_{nn}^{n}) \leq \epsilon \sum_{i=1}^{n} 2^{-n} = \epsilon$ $\left(\mu_{n}(\overline{A}^{c}) \leq \mu_{n}(A^{c}) = \mu_{n}(U(G_{c})^{c}) \leq \sum_{i} \mu_{n}(G_{i}^{n}) \leq \sum_{i \in I} 2^{-n} = C_{i}\right)$

(can be covered by a finite union of blue of any given resting).

Existence of Solitions to OT do

Lennelllet PED(X) 4 QED(Y) be tight. Then M(P,Q) = {T+D(Xx Y): Proje # T & P & proje # T & Q \ is kight. Your Fix 500 & Kx, Ky competer with ... $\forall \pi \in \Pi(P,Q), \pi(\chi \times \chi \times \chi \times \chi) \in \Pi(\chi \times \chi) \in \Pi(\chi \times \chi) + \Pi(\chi \times \chi)$ = p(xxx) + v(yxx) < 2 €.

Lennal M(M, N) is unpot.

Knot It is relatively compact by Prolehorov & the above lenne so so fix have to exhibit observer. Let II be e but of In. The to Jew + 4(5) ATT = 6- Jew + 4(5) ATT = Jew + 14(1)

It he used that $C_{\mathcal{C}}(\mathcal{X})$ determine elements in P(X). In fact one an constructive countrable family of functions that loss that.

Lenna 2.3 Suppose C: XXY wo Rultwolis Iscand bounded from below. Then TI -> JC-lt is lsc on B(XXY) with trology of week av. Lemma 2.4. For C: Z -> Ruspad bounded for blu c is loc (2) = surfu (2) for a family (fully 2) of Lipschitz further on Z. Prof (2.4). I fy (x) = but fx (xn) = but c(xn) since c > fx. letoling are c(x) & lengt c(xn). (Ph) Hore generally a sup of loc functions is loc. (2) (2" prot) c lsc => the equiphraph } (z, u): u > c(z) is closed in Z x R

but -, - of sup = () equiphraphs. Dlog c ≥0. Let f(z)= inf (c(u) + kd(z,u)) · | fu (2,) - fu (2,) | = | inf (clu) + k d(2,u1)) - inf (c(u2) + k d(2,u2)) | (wlop synum v2 eX < inf c(n,)+4 d(2,14,)- c(u,) -4 d(2,14,) u, eZ ∀ pich u_u + Z s.t. c(u_u) + k ol(u_u, 2) < f_u(≥) + l_u ≤ l + l_u $\left(\int d(u_{u}z) \leq \frac{l+l_{u}-c(u_{u})}{k} \leq \frac{l+l_{u}}{u} \longrightarrow 0$ toling but in) re get c (2) < built c(uu) = l e introdu. By Johny gu = In 1 k we may serious the sequence or of the furthers Proof (2.3)

We have the can take a seprence on the of Lip & bounded functions.

Then The Jobs John Constitution is content to Jost Hot Under Joseph State of the Ph above.

(by the Ph above).

Ihm 25 Let X, Y be Polish & pt D(X), v & D(Y) & c: Xx Y -> R Ufter loanshed below. Then the Kuntonovich of problem is extended P(pro)= inf Jcdv = Jcdin for some we M(pro) those We know that M/p, v) is compact, IT - fait is like so are conclude by bloompass. We dreedy noted that while M(PRV) is non-early, the cet of tourpits M(Prv)= for (Mpv): iney well be empty (e.g. p-Jos, r= N(a)). Lenne II pive P(Rd) & pir domber the T/piv) + \$. If piv are suppled on a compact then There is dense in M(m, v) I for a continuous: in f Scott = min Scott.

Exemple:

Extensions + HOT We can give here Y = X + ONE step martingules. Recoll the U(p,v) = { TI + T(p,v): ET [Y | T(x)] = X p-en { = { TI + T(p,v): TI = M&O } } fy Q (x, ly) = x p (4x) - e.e. } This set may be empty. In fact: y(.b) < 00 } Then (Schorsen) Let MINE B, (Rd) = { Je B(X): J |X| M(r,2) = 0 (f) p<0x , i.e., If the eff to F:R R conver Cenne M(riv) + & the M(riv) is compail. Front (for R i.e. dol)

Indeed, it is a subset of a compact
set so we just need to prove it is chosed. Recall that TEM (min) blue to M(min) of Jelx) (y-x) -III = 0 Fee C. (x) Jy (x- Jy O(x dy)) p (blx) Let $\overline{u}_n \in \mathcal{H}(p, \omega)$ can welly to $\overline{u} \in \mathcal{H}(p, \omega)$. Fix K > 0 & $f_n = \begin{cases} 1 & a & (-K, K)^2 \\ 0 & a & (-K, K)^2 \end{cases}$ $g_u = f(x)(y-x)f_u(x)y$ is $G_0 = \begin{cases} g_u \neq \overline{u}_n \longrightarrow \int g_u d\overline{u} \end{cases}$ cut. We by $f_n = f(x)(y-x)f_u(x)y$

= 1 13-9 1 10 - 9 1 10 = 1 6 (14+41) du = 6 (1 1x10++ 1 131/4) = 6 (1 1x10++ 1 131/4) = 6 (1 1x10++ 1 131/4) = 6 => | [8 str-olt/ = | 18.90 17, -11 | + | 34 15- 1- | = 35 00 TE M/MIN) The above extends to mg with faite discrete time M/p1,1-1,pm) but fails , e.g., in cations as time with M/p0, p1). Ils Going book to OT, a restriction of an optimal tr. plan is still optimal: Prof 2.6. In the setting of Th- 25, if TI is a minimiser of T' = T' is a non-negative measure with $T'(\chi \times y) > 0$ that $T' := \frac{T'}{T'(\chi \times y)}$ is an optimal to place for marginals $f' = f''(\chi \times y)$ If if not extend then take an minister \overline{u} , $\int c d\overline{r} = \int c d\overline{r}$ 下,命 (円/戸,3)

Lot T:= (T-T) + T'(XxY). T = T+ T'(XxY). (T-T) & (T/p10) L ∫ c d\(\tau < ∫ c d\(\tau \) = contradiction.

Some gropeties of the optimal solutions

A natural way to dry to inprove a given from port plex is is to consider if we can lower the cost who a cyclical relabelling of points.

Del (c-cyclical monotonicity) For c: 1x y -> Rulting a subset if a xx y

(s said to be c-cyclically monotone if the N (xxy),..., (x, yn) & M

I c(x, y;) \leq \int \frac{1}{121} c(x, y; h) \, there \frac{1}{121} \, \frac

correct cet I cost after cyclical re-rathly

A transport plus To = 3 (18x4) is said to be c-cyclically mentione if it is concontrated

on a -3 -5 -5 set.

Intrible: The M (p, D) optimal = IT is c-cyclically monotine

Tright from about: It also holds.