Computational Algebraic Topology Topic B: Sheaf cohomology and applications to quantum non-locality and contextuality Lecture 4

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$$\rho_U^V : P(V) \longrightarrow P(U)$$

subject to the functoriality requirements: if $U \subseteq V \subseteq W$, then

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Functoriality is easily verified: in this notation

$$(f|_V)_U = f|_U.$$

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- The category of all presheaves on a space X has a very rich structure it is a **topos**. We shall not go into this aspect.
- However, there is an important conceptual aspect which should be understood. Presheaves allow us to formalise the concept of **variable set**. The variation is essentially over **contexts**. So presheaves provide the natural setting for talking about contextuality!

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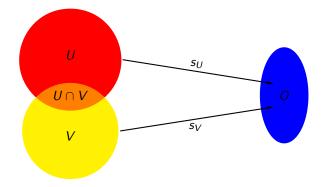
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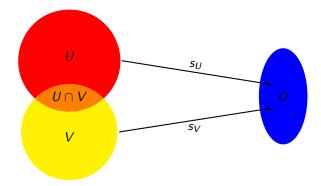
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The presheaf *P* is a **sheaf** if for every open cover \mathcal{U} , it satisfies the sheaf condition for \mathcal{U} .

Gluing functional sections



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If $s_U|_{U\cap V} = s_V|_{U\cap V}$, they can be glued to form

$$s: U \cup V \longrightarrow O$$

such that $s|_U = s_U$ and $s|_V = s_V$.

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In particular, this is one of the main intuitions behind **sheaf cohomology**.

		(0,0)	(1,0)	(0,1)	(1,1)	
а	Ь	0	1/2	1/2	0	
a'	Ь	3/8	1/8	1/8	3/8	
а	b'	3/8	1/8	1/8	3/8	
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Each row of the table specifies a **probability distribution** on events O^C for a given choice of measurements C.

Presheaves, Sheaves and Gluing

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The fact that the behaviour of these observable outcomes cannot be accounted for by some context-independent global description of reality corresponds to the geometric fact that these local sections cannot be glued together into a **global section**.

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The quantum phenomena of **non-locality** and **contextuality** correspond exactly to the existence of obstructions to global sections in this sense.

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Features in tropical geometry (the max-plus semiring).

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Contextual Probability Theory!

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Then we can define the presheaf

$$\mathcal{F}: \mathcal{P}(X)^{\mathrm{op}} \longrightarrow \mathbf{Set} :: U \mapsto \mathcal{D}_R(O^U)$$

A measurement scenario is a structure $\langle X, \mathcal{M}, O \rangle$ where:

- X is a set of "measurement labels" or "variables"
- \mathcal{M} is a family of subsets of X with $\bigcup \mathcal{M} = X$; the "measurement contexts"
- O is a set of "outcomes" or "values"

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A setting for contextual probability.

Empirical Models: Reconstructing Probability Tables

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An empirical model for μ is a family $\{e_C\}_{C \in \mathcal{M}}, e_C \in \mathcal{D}_R \mathcal{E}(C)$, which is compatible: for all $C, C' \in \mathcal{M}$,

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E.g. in the bipartite case, consider $C = \{m_a, m_b\}$, $C' = \{m_a, m'_b\}$. Fix $s_0 \in \mathcal{E}(\{m_a\})$. Compatibility implies

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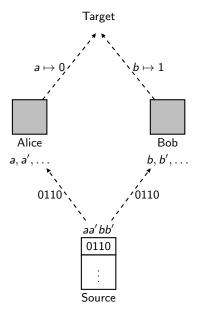
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This says that the probability for Alice to get the outcome $s_0(m_a)$ is the same, whether we marginalize over the possible outcomes for Bob with measurement m_b , or with m'_b .

In other words, Bob's choice of measurement cannot influence Alice's outcome.

Hidden Variables: The Mermin instruction set picture



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Question: does there exist a global section for this family?

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If *d* is a global section for the model $\{e_C\}$, we recover the predictions of the model by **averaging over the values of these hidden variables**:

$$e_{\mathcal{C}}(s) = d|\mathcal{C}(s) = \sum_{s'\in\mathcal{E}(X),s'|\mathcal{C}=s} d(s') = \sum_{s'\in\mathcal{E}(X)} \delta_{s'|\mathcal{C}}(s) \cdot d(s').$$

Note also that this is a **local** model:

$$\delta_s|C(s') = \prod_{x\in C} \delta_{s|x}(s'|x).$$

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Theorem

Any factorizable (i.e. local) hidden-variable model defines a global section.

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So:

existence of a local hidden-variable model for a given empirical model IFF empirical model has a global section

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Any factorizable (i.e. local) hidden-variable model defines a global section.

Hence:

No such h.v. model exists (the empirical model is **non-local/contextual**) IFF there is an **obstruction to the existence of a global section**

S. Abramsky and L. Hardy, Logical Bell Inequalities, *Phys. Rev. A* 85, 062114 (2012).

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Sheaf cohomology.

S. Abramsky, S. Mansfield and R. Soares Barbosa, The Cohomology of Non-Locality and Contextuality, in *Proc. QPL 2011*, EPTCS v. 95:1–15, 2012.

The Hardy Model and the PR Box

The Hardy Model and the PR Box

	(0,0)	(0, 1)	(1, 0)	(1, 1)
(a_1,b_1)	1	1	1	1
$egin{array}{llllllllllllllllllllllllllllllllllll$	0	1	1	1
(a_2, b_1)	0	1	1	1
(a_2, b_2)	1	1	1	0

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		(0,0)				
a_1	b_1	1 1 1 0	0	0	1	
a_1	<i>b</i> ₂	1	0	0	1	
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a ₂	<i>b</i> ₂	0	1	1	0	
		· -				

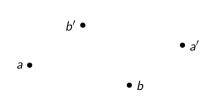
The PR Box

- Ignore precise probabilities
- Events are possible or not
- E.g. the Hardy model:

	00	01	10	11
ab	\checkmark	\checkmark	\checkmark	\checkmark
ab'	×	\checkmark	\checkmark	\checkmark
a' b	×	\checkmark	\checkmark	\checkmark
a' b'	\checkmark	\checkmark	\checkmark	×

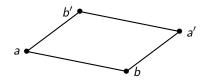
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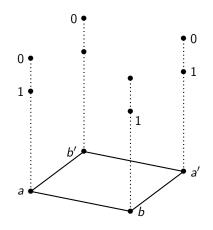
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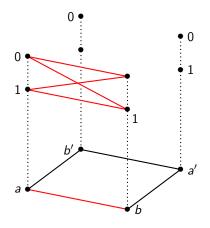
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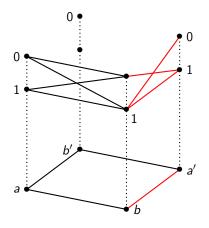
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a' b'	\checkmark	\checkmark	\checkmark	×



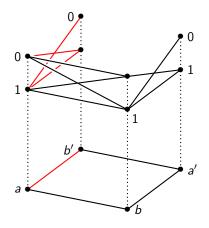
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a' b'	\checkmark	\checkmark	\checkmark	×



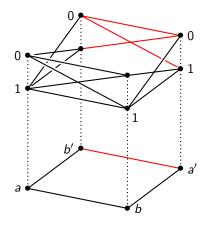
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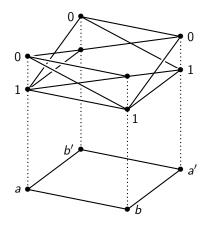
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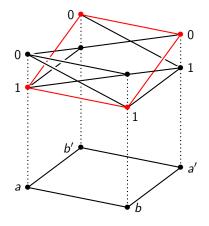
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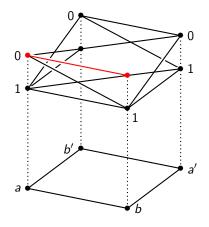
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a' b'	\checkmark	\checkmark	\checkmark	×



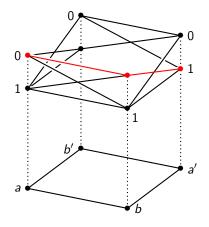
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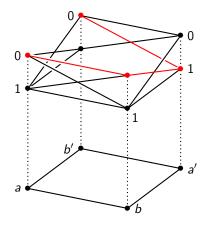
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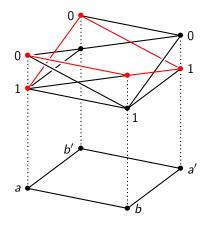
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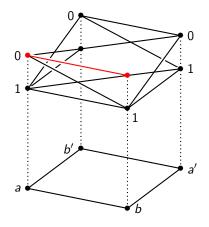
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a' b'	\checkmark	\checkmark	\checkmark	×



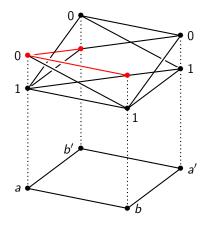
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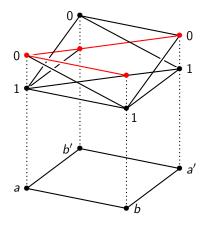
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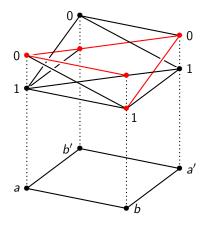
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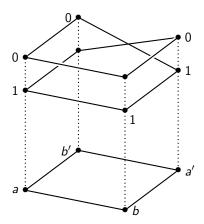
Strong Contextuality

			(1,0)	(0,1)	(1, 1)	
a_1	b_1	1 1 1 0	0	0	1	
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a ₂	b ₂	0	1	1	0	
The PR Box						

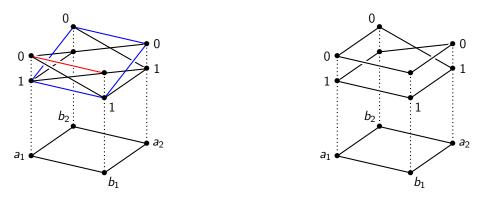
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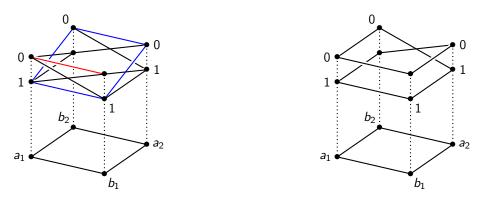


Visualizing Contextuality



The Hardy table and the PR box as bundles

Visualizing Contextuality

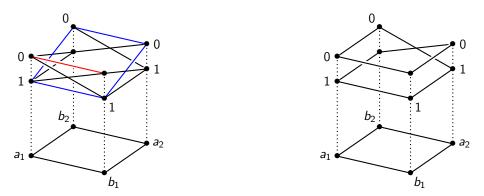


The Hardy table and the PR box as bundles

A hierarchy of degrees of contextuality:

 $\mathsf{Bell}\ <\ \mathsf{Hardy}\ <\ \mathsf{GHZ}$

Visualizing Contextuality



The Hardy table and the PR box as bundles

Formally, take

$$X := \prod_{i \in I} X_i := \{(i, x) : i \in I, x \in X_i\}$$

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Conversely, given $p: X \longrightarrow I$, we can form the indexed family $\{X_i\}_{i \in I}$, where $X_i := p^{-1}(\{i\})$.

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$$X := \prod_{i \in I} X_i := \{(i, x) : i \in I, x \in X_i\}$$

The family is $\phi: I \longrightarrow \mathcal{P}(X)$, $\phi(i) = \{(i, x) : x \in X_i\}$.

There is also a natural projection function

$$p: X \longrightarrow I \qquad p: (i, x) \mapsto i$$

Conversely, given $p: X \longrightarrow I$, we can form the indexed family $\{X_i\}_{i \in I}$, where $X_i := p^{-1}(\{i\})$.

These are equivalent ways of looking at the same idea.

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Sheaves on X are equivalently formulated as continuous maps $p: Y \longrightarrow X$ which are **local homeomorphisms** (*espaces étalé*).

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Thus in terms of well-known examples, we have

 $\mathsf{Bell} < \mathsf{Hardy} < \mathsf{GHZ}$