

## MODEL THEORY I — EXERCISE 6

Question 1 (repeat question from Exercise 4)

Let  $L = \{R\}$  where  $R$  is a binary relation symbol.

Let  $Gr$  be the theory of graphs (saying that  $R$  is symmetric, anti-reflexive).

Fix  $m, n$  and consider  $\alpha_{m,n}$  which states:

$$\forall x_1 \dots x_m \forall y_1 \dots y_m \left( \bigwedge_{i,j} x_i \neq y_j \rightarrow \exists z \bigwedge_i R(x_i, z) \wedge \bigwedge_j \neg R(y_j, z) \wedge y_j \neq z \right).$$

- (1) Let  $N \in \mathbb{N}$  and let  $p_N$  be the probability that  $\alpha_{m,n}$  holds in a graph  $G$  with vertex set  $V_N = \{1, \dots, N\}$  (there are  $2^{\binom{N}{2}}$  such graphs, and each is given the same probability). Show that  $\lim_{n \rightarrow \infty} p_N = 1$ .

Hint: fix  $a_1, \dots, a_m, b_1, \dots, b_n \in V_N$  such that  $a_{i'} \neq a_i \neq b_j \neq b_{j'}$  and compute the probability  $p_N(\bar{a}, \bar{b})$  that  $G \models \neg \exists z \bigwedge_i R(a_i, z) \wedge \bigwedge_j \neg R(b_j, z) \wedge b_j \neq z$ . Show that it is  $c^{-(N-m-n)}$  where  $c = 1 - 2^{-m-n} < 1$  (so exponentially goes to 0). Note that  $1 - p_N \leq \sum_{\text{choices of } \bar{a}, \bar{b}} p_N(\bar{a}, \bar{b})$ . (You should be a bit careful as  $\alpha_{m,n}$  didn't ask that  $x_i \neq x_{i'}$ .)

- (2) Let  $RG = \{\alpha_{m,n} \mid m, n \in \mathbb{N}\} \cup Gr$  ( $RG$  stands for Random Graph). Conclude that  $RG$  is consistent.
- (3) Show by induction on  $k < \omega$  that any finite graph of size  $k$  can be embedded in a model of  $RG$ . Conclude that  $RG_{\forall} = Gr$ .
- (4) Show that  $RG$  eliminates quantifiers.

Question 2\*

Let  $T = Th(\mathbb{N}, +, \cdot, 0, 1, \leq)$  (or, if you prefer, Peano arithmetic). A model  $N$  is said to be an *end extension* of  $M$  iff  $M \neq N$ ,  $M \prec N$  and, for all  $b \in N$  and  $a \in M$ , if  $b < a$  then  $b \in M$ . Prove that every countable model of  $T$  has an end extension.

Hint: the omitting types theorem.

Question 3

- (1) Let  $T$  be countable and consistent. Then any meager subset  $X \subseteq S_n(T)$  can be omitted (see Exercise 5, Question 4).

- (2) Find an example of an uncountable theory where the omitting types theorem fails.

Hint: look at the formulation of the omitting types theorem, there is a very simple example.

- (3) Show that any complete theory whose language consists of only unary predicates has elimination of quantifiers.

#### Question 4

- (1) Let  $\Sigma$  be a (consistent) theory in a countable language and let  $S(\Sigma)$  be the set of its completions. Show that  $S(\Sigma)$  is either countable or has size  $2^{\aleph_0}$ .

Hint: try to find a sentence  $\alpha$  such that both  $\Sigma \cup \{\alpha\}$  and  $\Sigma \cup \{\neg\alpha\}$  have uncountably many completions.

- (2) \* Let  $\Sigma$  be as above. Let  $p$  be a partial type and let  $S(\Sigma, p)$  be the set  $\{Th(M) \mid M \models \Sigma, M \text{ omits } p\}$ . Show that  $S(\Sigma, p)$  is either countable or has size  $2^{\aleph_0}$ .

#### Question 5

Suppose  $G$  is a group acting on a set  $X$  (recall the definition of group action, if needed). We say that the action is *oligomorphic* if for all  $n$  the number of orbits of  $G$  acting on  $X^n$  is finite.

When  $M$  is a structure, we can think of  $G = \text{Aut}(M)$  as acting on  $M$ .

Show that the following are equivalent for a complete countable theory  $T$ .

- (1)  $T$  is  $\omega$ -categorical.
- (2) Every countable model of  $T$  has an oligomorphic automorphism group.
- (3) Some countable model of  $T$  has an oligomorphic automorphism group.

#### Question 6

- (1) Show that if  $L \subseteq L'$  and  $T \subseteq T'$  are countable complete theories, and  $T'$  is  $\omega$ -categorical, then so is  $T$ .
- (2) Show that if  $T$  countable and is  $\omega$ -categorical,  $M \models T$  and  $A \subseteq |M|$  is finite then  $T_A = Th(M_A)$  is also  $\omega$ -categorical.
- (3) Does (2) hold with  $A$  infinite?

#### Question 7

Suppose  $T$  is a theory in  $L$ , not necessarily countable. We say that  $T$  is *small* if  $S_n(T)$  is countable for all  $n < \omega$ .

- (1) Suppose that  $T$  is small. Show that  $T$  is essentially countable: there is a countable sub-language  $L_0 \subseteq L$  such that any formula  $\varphi(\bar{x}) \in L$  is equivalent (modulo  $T$ ) to some formula  $\psi(\bar{x}) \in L_0$ .
- (2) Conclude that if  $M \models T \upharpoonright L_0$  (so  $M$  is an  $L_0$ -structure) then  $M$  can be enriched to be a model of  $T$ .
- (3) Show that if  $T$  is countable and not small then it has uncountably many countable models up-to isomorphism.

Question 8

Suppose  $T$  is a complete countable theory. Show that any countable model  $M \models T$  has a countable elementary extension which is  $\omega$ -homogeneous.

Question 9

- (1) Let  $M$  be countable with a countable language and  $(Th(M))$  is  $\omega$ -categorical. Show that if  $X \subseteq M^n$  is invariant under the action of  $\text{Aut}(M)$  then  $X$  is definable (without parameters).
- (2) \* Show that if  $M$  is an infinite field, then  $Th(M)$  is never  $\omega$ -categorical.
- (3) Suppose now that  $M$  is a structure (perhaps not  $\omega$ -categorical). Assume that for some  $n$  only finitely many  $n$ -types are realized in  $M$ . Then in any  $N \equiv M$ , the exact same  $n$ -types are realized.