1. (a) False in N.  
(b) [  

$$\forall x : (x + x) = y$$
  
 $\exists y : (x + x) = y$   
 $\exists y : (x + x) = y$   
 $\forall x : \exists y : (x + x) = y$   
]  
 $\langle \forall x : (x + x) = y \supset \forall x : \exists y : (x + x) = y \rangle$   
[  
 $\sim \forall x : \exists y : (x + x) = y$   
 $\langle \neg \forall x : \exists y : (x + x) = y$   
 $\langle \neg \forall x : \exists y : (x + x) \supset \neg \forall x : (x + x) = y \rangle$   
 $\sim \forall x : (x + x) = y$   
 $\forall y : \neg \forall x : (x + x) = y$   
 $\forall y : \neg \forall x : (x + x) = y$   
]  
 $\langle \neg \forall x : \exists y : (x + x) = y \supset \forall x : \exists y : (x + x) = y \rangle$   
 $\langle \exists y : \forall x : (x + x) = y \supset \forall x : \exists y : (x + x) = y \rangle$ 

- (c) This is false in the structure of arithmetic modulo 2.
- (d) This is false in the structure obtained from ℕ by setting 0 to be interpreted as 37.
- (e) This is false in the structure obtained from  $\mathbb{N}$  by adjusting addition as follows: n + m := n + m if  $m \neq 2$ , but n + 2 = 73 for all n.
- 2. (a) TNT' can prove  $\forall x : (x \cdot S0) = (0 + x)$ , so if it could prove  $\forall x : (x \cdot S0) = x$  then it would also prove  $\forall x : (0 + x) = x$ . But as was mentioned in lectures, it can't prove that! So there must be a structure which satisfies the axioms of TNT' but in which  $\forall x : (x \cdot S0) = x$  is false.

How can we find such a structure? There are a few possibilities, here's a cute one: we add a second copy of  $\mathbb{N}$  to  $\mathbb{N}$ , and have addition and multiplication fold everything down to one copy or the other - folding to the left with addition, and to the right with multiplication.

More explicitly: we add new elements  $0', 1', 2', \ldots$  to  $\mathbb{N}$ , extend successor to the new elements by defining S(n') := (n + 1)', and extend addition by taking the primedness of the result to be the primedness of the first summand, and for multiplication take it from the second multiplicand. In other words:

$$n + m' = n + m; n' + m' = n' + m = (n + m)',$$

and

$$n' \cdot m = n \cdot m; \ n' \cdot m' = n \cdot m' = (n \cdot m)'$$

One can see that the axioms of TNT' hold for this structure. But e.g.  $(0' \cdot 0) = 0$ , so  $\forall x : (x \cdot S0) = x$  is false for this structure; similarly (0 + 0') = 0 so  $\forall x : (0 + x) = x$  is false in the structure. (Remark: it's common to work with a stronger system than TNT', known as Robinson arithmetic and denoted by Q, which adds to TNT' the axiom  $\forall x : \langle x = 0 \lor \exists y : Sy = x \rangle$  (which is a TNTtheorem). Our structure does not satisfy this last axiom. However,  $\forall x : (x \cdot S0) = x$  and  $\forall x : (0 + x) = x$  are also not theorems of Q. There's a reasonably pleasant example of a structure demonstrating this on page 7 of Peter Smith's *Gödel Without Tears* (linked to from the the course website), obtained from N by adding two extra elements.)

(b)  $(SS0 \cdot SS0) = ((SS0 \cdot S0) + SS0)$  $(SS0 \cdot S0) = ((SS0 \cdot 0) + SS0)$  $(SS0 \cdot 0) = 0$  $((SS0 \cdot 0) + SS0) = (0 + SS0)$  $(SS0 \cdot S0) = (0 + SS0)$ (0 + SS0) = S(0 + S0)(0 + S0) = S(0 + 0)(0 + 0) = 0S(0 + 0) = S0(0 + S0) = SS0(0 + S0) = SS0 $(SS0 \cdot S0) = SS0$  $(SS0 \cdot S0) = SS0$  $(SS0 \cdot S0) = SS0$  $(SS0 \cdot SS0) = (SS0 + SS0)$  Remark: We had to prove here (0 + SS0) = SS0: this is a theorem of TNT', even though  $\forall x : (0 + x) = x$  is not.

- (c) This is false in  $\mathbb{N}$ .
- (d) This is false in  $Mat_2(\mathbb{N})$ .