Inclusion-Exclusion

Let A_1 and A_2 be finite subsets of a set X. If A_1 and A_2 are disjoint, the addition principle tells us $|A_1 \cup A_2| = |A_1| + |A_2|$.

If they're not disjoint, this "double-counts" the elements of the intersection; we can fix this by subtracting the size of the intersection, yielding the general formula $|A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2|$.

For three finite subsets A_1, A_2, A_3 of some set X, similar reasoning yields

$$\begin{aligned} |A_1 \cup A_2 \cup A_3| &= \\ |A_1| + |A_2| + |A_3| - (|A_1 \cap A_2| + |A_2 \cap A_3| + |A_1 \cap A_3|) + |A_1 \cap A_2 \cap A_3| \end{aligned}$$

Example:

How many integers in [1,100] are divisible by 2, 5, or 7?

Let $D_n := \{k \in \mathbb{Z} \cap [1, 100] : n \mid k\}.$

Note $|D_n| = \lfloor \frac{100}{n} \rfloor$.

By inclusion-exclusion,

$$\begin{aligned} |D_2 \cup D_5 \cup D_7| &= |D_2| + |D_5| + |D_7| \\ &-(|D_2 \cap D_5| + |D_5 \cap D_7| + |D_2 \cap D_7|) \\ &+ |D_2 \cap D_5 \cap D_7| \\ &= |D_2| + |D_5| + |D_7| - (D_{10} + D_{35} + D_{14}) + D_{70} \\ &= 50 + 20 + 14 - (10 + 2 + 7) + 1 \\ &= 66. \end{aligned}$$

Theorem [Inclusion-Exclusion Principle]:

Let $A_1, ..., A_n$ be finite subsets of a set X. Then

$$|A_1 \cup \ldots \cup A_n| = \sum_i |A_i| - \sum_{i < j} |A_i \cap A_j| + \sum_{i < j < k} |A_i \cap A_j \cap A_k| - \ldots + (-1)^{n-1} |A_1 \cap \ldots \cap A_n|$$

Proof:

Let $x \in A_1 \cup ... \cup A_n$. We show that x "contributes 1" to the right hand side. Say x is in $m \ge 1$ of the n sets. Then x contributes 1 to $m = \binom{m}{1}$ of the $|A_i|$, to $\binom{m}{2}$ of the $|A_i \cap A_j|$, to $\binom{m}{3}$ of the $|A_i \cap A_j \cap A_k|$, and so on.

So x contributes $\sum_{k>0} (-1)^{k-1} {m \choose k}$ $= -\sum_{k>0} (-1)^k {m \choose k}$

$$= 1 - \sum_{k \ge 0} (-1)^k {m \choose k} \\ = 1 - 0 \\ = 1$$

Remark:

Neat alternative expression:

$$|\bigcup_{i} A_{i}| = \sum_{\emptyset \neq I \subseteq \{1, \dots, n\}} (-1)^{|I|-1} \left| \bigcap_{i \in I} A_{i} \right|.$$

Bonus:

Version of the proof using this notation:

$$\begin{split} \bigcup_{i} A_{i} \bigg| &= \sum_{x \subseteq \bigcup_{i} A_{i}} 1 \\ &= \sum_{x \subseteq \bigcup_{i} A_{i}} \left(-\sum_{k>0} (-1)^{k} \binom{\#\{i \mid x \in A_{i}\}}{k} \right) \right) \\ &= \sum_{x \subseteq \bigcup_{i} A_{i}, I \subseteq \{i \mid x \in A_{i}\}} (-1)^{|I|-1} \\ &= \sum_{\{(x,I) \mid x \in \bigcup_{i} A_{i}, I \subseteq \{1, \dots, n\}, x \subseteq \bigcap_{i \in I} A_{i}\}} (-1)^{|I|-1} \\ &= \sum_{\emptyset \neq I \subseteq \{1, \dots, n\}} (-1)^{|I|-1} \left| \bigcap_{i \in I} A_{i} \right| \end{split}$$

Example:

How many strings of 8 letters from the Roman alphabet contain 'j', 'q', 'x', 'y' and 'z'?

We could do this positively, but it would be fiddly.

Instead, let's count the number of strings which ${\bf don't}$ contain all of these letters, i.e. which omit 'j' or omit 'q' or... .

Let O_j be the strings which omit 'j', O_{jq} the strings which omit 'j' and 'q', and so on.

Then by inclusion-exclusion,

$$\begin{split} |O_{j} \cup O_{q} \cup O_{x} \cup O_{y} \cup O_{z}| &= |O_{j}| + |O_{q}| + \dots \\ &-(|O_{jq}| + |O_{jx}| + \dots) \\ &+(|O_{jqx}| + |O_{jqy}| + \dots) \\ &-(|O_{jqxy}| + |O_{jqxz}| + \dots) \\ &+|O_{jqxyz}| \end{split}$$

Now $|O_j| = |O_q| = \dots = 25^8$, and $|O_{jq}| = |O_{jx}| = \dots = 24^8$, and so on.

 So

$$\begin{array}{l} O_j \cup O_q \cup O_x \cup O_y \cup O_z | = \\ \binom{5}{1} 25^8 - \binom{5}{2} 24^8 + \binom{5}{3} 23^8 - \binom{5}{4} 22^8 + \binom{5}{5} 21^8 \end{array}$$

and the answer to the original question is $26^8 - |O_j \cup O_q \cup O_x \cup O_y \cup O_z|$, which comes to 87408720.

Combinations of multisets, revisited

Recall:

The number of r-combinations of a multiset with at least r of each of its t types is $\binom{r+t-1}{t-1}$.

If there are fewer than r of some of the types, we can use inclusion-exclusion.

This is clearest if we transform the problem.

An r-combination of a multiset

 $\{c_1 * a_1, ..., c_t * a_t\}$ corresponds to a solution in non-negative integers of the equation $x_1 + ... + x_t = r$ subject to the constraints $x_1 \leq c_1, ..., x_t \leq c_t.$

Concrete example:

I take 8 marbles from a bag containing 3 red marbles, 2 blue marbles, and 10 green marbles. How many possibilities are there for the numbers of each colour I get? Equivalently, what is

$$|\{(r, b, g) \mid r + b + g = 8, 0 \le r \le 3, 0 \le b \le 2\}|?$$

So we want to count the number of such solutions, and we know that the answer is $\binom{r+t-1}{t-1}$ if there are no constraints.

By the subtraction principle,

the number of solutions in the constrained case is the number in the unconstrained case minus the number which fail at least one constraint,

 $\binom{r+t-1}{t-1} - |F_1 \cup \dots \cup F_t|,$ where $F_i := \{(x_1, \dots, x_t) \mid x_1 + \dots + x_t = r, x_i > c_i\}.$

So we can use the inclusion-exclusion principle if we can determine the sizes of the intersections of the F_i .

If $c_i \ge r$, then $F_i = \emptyset$. Otherwise, subtracting $c_i + 1$ from x_i puts F_i in correspondence with $\{(y_1, ..., y_t) \mid y_1 + ... + y_t = r - (c_i + 1), y_i \ge 0\}$, so $|F_i| = \binom{r - (c_i + 1) + t - 1}{t - 1}$.

Similarly, $|F_i \cap F_j| = \binom{r-(c_i+1)-(c_j+1)+t-1}{t-1}$, and so on.

So inclusion-exclusion yields $\sum_{n=1}^{\infty} \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} \sum_{i=1}^{\infty$

$$\sum_{I \subseteq \{1,...,t\}} (-1)^{|I|} \binom{r - (\sum_{i \in I} (c_i + 1)) + t - 1}{t - 1}.$$

Marble example:

$$\binom{8+2}{2} - \binom{8-(3+1)+2}{2} - \binom{8-(2+1)+2}{2} + \binom{8-(3+1)-(2+1)+2}{2}$$

= $\binom{10}{2} - \binom{6}{2} - \binom{7}{2} + \binom{3}{2}$
= 12

Scrabble example:

How many 7-tile hands can be drawn from a standard 100-tile bag of scrabble tiles?

Using the above formula, my computer calculates it as 3199724.

(for the curious, here's the Haskell code I used to calculate this:

```
import Math.Combinatorics.Binomial (choose)
 combs :: Int -> [Int] -> Int
combs r cs =
   let
        t = length cs
        subs = subs' [] cs
        -- subs': returns relevant subsequences of cs, omitting those which
        -- will contribute 0 to the final sum (without this, the algorithm
        -- would have complexity exponential in t)
        subs' sub [] = [sub]
        subs' sub _ | (sum (map (+1) sub) > r) = []
        subs' sub (c:cs) = subs' (c:sub) cs ++ subs' sub cs
 in sum [ (-1)^{(length sub)} *
         choose (r - sum (map (+1) sub) + t-1) (t-1) | sub <- subs ]
scrabbleBag :: [Int]
scrabbleBag = concat [ replicate n c | (n,c) <-</pre>
         [(5,1), (10,2), (1,3), (4,4), (3,6), (1,8), (2,9), (1,12)]
main :: IO ()
main = print $ combs 7 scrabbleBag
)
```

Derangements

A derangement is a permutation which leaves nothing in its original position.

e.g.

(5,3,4,2,1) is a derangement of (1,2,3,4,5), and "endgreatmen" is a derangement of "derangement".

 $D_n :=$ the number of derangements of a sequence of length n, = number of derangements of (1, 2, ..., n).

We can use inclusion-exclusion to determine D_n .

A derangement of (1, ..., n) is a permutation $(a_1, ..., a_n)$ which satisfies the conditions $a_1 \neq 1, ..., a_n \neq n$.

Let P_i be the set of permutations which fail the *i*th of these conditions, i.e. such that $a_i = i$.

Easily, for $I \subseteq \{1, ..., n\}$, $\left|\bigcap_{i \in I} P_i\right| = (n - |I|)!.$

So by inclusion-exclusion,

$$D_n = n! - |\bigcup_i P_i|$$

= $n! - \sum_{\emptyset \neq I \subseteq \{1,...,n\}} (-1)^{|I|-1} |\bigcap_{i \in I} P_i|$
= $n! - \sum_{\emptyset \neq I \subseteq \{1,...,n\}} (-1)^{|I|-1} (n - |I|)!$
= $n! - \sum_{i=1}^n (-1)^{i-1} {n \choose i} (n - i)!$
= $\sum_{i=0}^n (-1)^i {n \choose i} (n - i)!$
= $\sum_{i=0}^n (-1)^i \frac{n!}{i!}$
= $n! \sum_{i=0}^n \frac{(-1)^i}{i!}$

Note then that the probability that a random n-permutation is a derangement is

$$Pr_n = \frac{|D_n|}{n!} = \sum_{i=0}^n \frac{(-1)^i}{i!},$$
$$\lim_{n \to \infty} Pr_n = \sum_{i=0}^\infty \frac{(-1)^i}{i!}$$
$$= e^{-1} \approx 0.368$$

The convergence is very fast;

e.g. $Pr_n \approx 0.368$ to 3 significant figures for $n \ge 6$.

Example:

 \mathbf{SO}

A deranged scientist removes the heads from a large number of different animals and re-attaches them at random. What is the probability that every resulting creature is a chimera, i.e. that no head is reattached to its own body?

Answer: About e^{-1} .