DUALITIES, NATURAL AND OTHERWISE

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GAIA 2013

With acknowledgements, in particular, to Brian Davey and Leonardo Cabrer

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From Brian Davey's La Trobe staff profile

Professor Davey's research interests centre on general (= universal) algebra and lattice theory. He is interested in the topological representation of algebras, with an emphasis on natural duality theory, and is particularly interested in applications of duality theory in general and Priestley duality in particular to algebras with an underlying distributive lattice structure.

But how did it all begin?

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Mathematics Department Monash University, Clayton, Vic. 3168.

Dr. H. A. Priestly, Mathematical Institute, OXFORD, ENGLAND.

Dear Dr. Priestly,

I have just received a preprint of T.F. Speed's paper. "Profinite Posets", in which he refers to two of your papers, namely, "Representations of distributive lattices by means of ordered stone spaces" and "Ordered topological spaces and the representation of distributive lattices".

I am greatly interested in this topic and if possible I would like to purchase a copy of your thesis, but if that is not possible copies of these papers would be appreciated.

Yours sincerely,

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Outline

Setting the scene: natural dualities for quasivarieties and varieties of D-based algebras
 (D the variety of bounded distributive lattices)

An example: a hierarchy of prioritised default bilattices

- Natural dualities versus \mathcal{D} - \mathcal{P} -based dualities
- The piggyback strategy: more than just a method for finding economical dualities

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Natural dualities and \mathcal{D} - \mathcal{P} -based dualities related

Examples

• Coproducts in \mathfrak{D} -based varieties

Setting the scene

- Let D denote the variety of bounded distributive lattices.
 Take a finitely generated quasivariety A = ISP(M), where M is a finite set of finite D-based algebras. (Jónsson's Lemma ensures that every finitely generated variety of lattice-based algebras can be expressed this way.)
- Assume we have a forgetful functor $U \colon \mathcal{A} \to \mathcal{D}$.

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Priestley duality supplies a dual equivalence between

$$\mathfrak{D} = \mathbb{ISP}(2)$$
 and $\mathfrak{P} = \mathbb{IS}_{c}\mathbb{P}^{+}(2)$ (\equiv Priestley spaces)

set up by hom-functors $H = \mathfrak{D}(-, \mathbf{2})$ and $K = \mathfrak{P}(-, \mathbf{2})$. Here

$$\mathbf{2} = \langle \{0,1\}; \wedge, \vee, 0,1\rangle \quad \text{and} \quad \mathbf{2} = \langle \{0,1\}; \leqslant, \mathfrak{T}\rangle.$$

Here, and henceforth, \mathcal{T} denotes the discrete topology.

Objective

To find an alter ego ${\mathfrak M}$ for ${\mathfrak M}$ so there exists a dual equivalence between

$$\mathcal{A} = \mathbb{ISP}(\mathfrak{M})$$
 and $\mathfrak{X} = \mathbb{IS}_{c}\mathbb{P}^{+}(\mathfrak{M})$

set up by hom-functors $D = \mathcal{A}(-, \mathfrak{M})$ and $E = \mathfrak{X}(-, \mathfrak{M})$, with ED and DE given by natural evaluation maps.

Can we do this?: YES!

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Case in which \mathfrak{M} contains one algebra, \mathbf{M} , is probably familiar.

Quick recap: single-sorted case

- Let $\mathcal{A} = \mathbb{ISP}(\mathbf{M})$, where \mathbf{M} is finite, \mathcal{D} -based. Take $\underline{\mathbf{M}} = \langle M; R, H, \mathfrak{T} \rangle$ where
- $R = \mathbb{S}(\mathbf{M}^2)$ (all algebraic binary relations on M)
- H = some suitable set of (partial or total) algebraic

operations, with a bound on the arities.

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- The algebraicity requirements guarantee the hom- functors D and E are well defined, and are embeddings;
- M lattice-based, hence the specified R yields a duality: $A \cong ED(A)$ for all A [by the NU Duality Theorem].
- **Upgrading**: with a suitably chosen set H, the duality is **strong** and hence full: $\mathbf{X} \cong \mathsf{DE}(\mathbf{X})$ for all $\mathbf{X} \in \mathfrak{X}$.

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Multisorted case, $|\mathfrak{M}| > 1$ (Davey and Priestley (1984, at LTU)): All the above carries over: Relations and (partial and total) operations are algebraic, but now between members of \mathfrak{M} . The structure dual of an algebra \mathbf{A} is based on the disjoint union of hom-sets $\mathcal{A}(\mathbf{A}, \mathbf{M})$, for $\mathbf{M} \in \mathfrak{M}$. Example to follow.

A hierarchy of prioritised default bilattices

(Cabrer, Craig & Priestley, 2013) We set up strong dualities for $\mathbb{ISP}(\mathbf{K}_n)$ and (multisorted) for $\mathbb{HSP}(\mathbf{K}_n)$, where

$$\mathbf{K}_n = (K_n; \wedge_k, \vee_k, \wedge_t, \vee_t, \neg, \top, \bot).$$



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 \mathbf{K}_n in knowledge order (left) and truth order (right) with 0 < i < j < n.

Interpretation for default logic

The elements of K_n represent levels of truth and falsity. Knowledge represented by the truth values at level m + 1 is regarded as having lower priority priority than from those at level m. Also one thinks of \mathbf{t}_{m+1} as being 'less true' than \mathbf{t}_m and \mathbf{f}_{m+1} as 'less false' than \mathbf{f}_m .

Base cases: \mathbf{K}_0 and \mathbf{K}_1 are Ginsberg's bilattices \mathcal{FOUR} and \mathcal{SEVEN} , with \neg added. In \mathcal{SEVEN} , \mathbf{t}_1 and \mathbf{f}_1 may be given the connotation of 'true by default' and 'false by default'.



Algebraic facts

- Every element of K_n is term-definable; K_n has no proper subalgebras.
- For $m \leq n$ there exists a surjective homomorphism $h_{n,m} \colon K_n \to K_m$.
- For $0 \leqslant m \leqslant n$, there exists $\S_{n,m} \in \mathbb{S}(\mathbf{K}_n^2)$ with elements

$$\S_{n,m} = \Delta_n \cup \{ (a,b) \mid a, b \leqslant_k \top_{m+1} \text{ or } a \leqslant_k b \leqslant_k \top_m \}.$$

where $\Delta_n = \{ (a, a) \mid a \in K_n \}$. We have $S_{n,j} \subseteq S_{n,i}$ for $0 \leq i < j \leq n$.

- The subalgebras §_{n,m} entail every subalgebra of K²_n via converses and intersections.
- Each **K**_n is subdirectly irreducible.
- $\mathbb{ISP}(\mathbf{K}_n) = \mathbb{HSP}(\mathbf{K}_n)$ iff n = 0. For $n \ge 1$,

 $\mathbb{HSP}(\mathbf{K}_n) = \mathbb{ISP}(\mathbf{K}_n, \dots, \mathbf{K}_0).$

Algebraic binary relations n = 0: $\mathbb{ISP}(\mathbf{K}_0)$ is the variety \mathcal{DB} of distributive bilattices. $S_{0,0}$ is the knowledge order, \leq_k .

De Morgan algebras, \mathcal{M} , and \mathcal{DB} compared. $\mathcal{M} = \mathbb{ISP}(\mathbf{M})$ where \mathbf{M} has elements $\mathbf{t}.\mathbf{f}, \top, \bot$, with \wedge_t, \vee_t, \neg as in \mathbf{K}_0 . and constants \mathbf{t} and \mathbf{f} . Here $]\mathbb{S}(\mathbf{M}^2)] = 55$ whereas $]\mathbb{S}(\mathbf{K}_0^2)] = 4$. In both cases, alter ego for an optimal duality includes \leq_k ; for De Morgan we need also \mathbf{M} 's non-identity automorphism.

n = 1: Algebraic relations $S_{1,0}$ and $S_{1,1}$ on K_1 , depicted as quasi-orders.





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Duality theorems

Theorem

The structure $\mathbf{K_n} = \langle K_n; S_{n,n}, \dots, S_{n,0}, \mathfrak{T} \rangle$ yields a strong, and optimal, duality on $\mathbb{ISP}(\mathbf{K}_n)$.

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For n = 0 the dual category $\mathbb{IS}_{c}\mathbb{P}^{+}(\mathbf{K}_{0})$ is \mathfrak{P} .

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Theorem

Write $\mathbb{HSP}(\mathbf{K}_n)$ as $\mathbb{ISP}(\mathfrak{M})$, where $\mathfrak{M} = {\mathbf{K}_0 \dots, \mathbf{K}_n}$. Then the alter ego

$$\mathfrak{M} = \langle K_0 \, \dot{\cup} \, \ldots \, \dot{\cup} \, K_n; \{ S_{m,m} \}_{0 \leqslant m \leqslant n}, \{ h_{i,i-1} \}_{1 \leqslant i \leqslant n}, \mathfrak{T} \rangle,$$

yields a strong, and optimal, duality on $\mathbb{HSP}(\mathbf{K}_n)$.

The dual categories for $\mathbb{ISP}(\mathbf{K}_n)$ and $\mathbb{HSP}(\mathbf{K}_n)$ can be described for general *n*—comment on this later.

\mathcal{D} -based varieties: some old favourites

- De Morgan algebras;
- Kleene algebras;
- pseudocomplemented distributive lattices ℬ_ω and finitely generated subvarieties ℬ_n (ℬ₁ = Stone algebras);
- Heyting algebras, and Gödel algebras, 𝔅_n—generated by an *n*-element Heyting chain;

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n-valued Łukasiewicz–Moisil algebras;

etc, etc.

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Assume \mathcal{A} is a variety of \mathfrak{D} -based algebras.

- **1** Take the class $U(\mathcal{A})$ (the \mathcal{D} -reducts).
- 2 Seek to equip $\mathfrak{Z} := HU(\mathcal{A})$ with additional (relational or functional) structure so, for each $\mathbf{A} \in \mathcal{A}$, $KHU(\mathbf{A})$ becomes an algebra in \mathcal{A} isomorphic to \mathbf{A} ;
- 3 Identify a suitable class of morphisms, making \mathfrak{Z} into a category.

If this gives a dual equivalence between \mathcal{A} and \mathfrak{Z} , we say we have a \mathcal{D} - \mathcal{P} -based duality. Literature is full of examples!

Bringing things together

We have:

- Natural duality theory: for any finitely generated D-based variety we can call on the NU (Strong) Duality Theorem (single-sorted or multisorted).
- **2** \mathcal{D} - \mathcal{P} -based dualities
- **3** From, and for, logic: algebraic and relational (Kripke-style) semantics for non-classical propositional logics.

Both (1) and (2) provide valuable tools for studying algebraic properties of \mathcal{D} -based varieties.

Normally, For a given variety, a discrete duality (as in (3)) differs from a \mathcal{D} - \mathcal{P} -based just through the absence or presence of topology. Canonical extensions provide a systematic approach to (3).

Rivals?

	Pro	Con
natural duality	a strong duality can always be found	duality may be complicated (may need entailment to simplify)
	${f M}$ governs how the duality works	restriction to finitely gener- ated classes (usually)
	good categorical properties, notably w.r.t. free algebras: $D(F\mathcal{A})) = \underbrace{\mathbf{M}}^{S}$	concrete representation is via functions, not sets, if $\left M \right > 2$
D-P- based duality	close relationship to Kripke- style semantics	
	concrete representation via sets	products seldom cartesian; free algebras hard to find

or collaborators?

We'd like to be able to toggle between a natural duality and a $\mathfrak{D}\mathchar`-\mathfrak{P}\mathchar`-based one.$

To get the best of both worlds,

LET'S HITCH A

PIGGYBACK RIDE!

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The piggyback strategy: outline

For the case of $\mathcal{A}=\mathbb{ISP}(\mathbf{M}),$ where \mathbf{M} is finite and $\mathfrak{D}\text{-based}.$ Original objective

■ To choose a subset of S⁽M²) yielding an economical duality, so bypassing entailment arguments.

Davey & Werner' (1982, at LTU) Simple piggybacking achieved this for **certain** quasivarieties: **De Morgan algebras**: 55 relations reduced to 2. But no good, e.g., for **Kleene algebras**.

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Further objectives: to use piggybacking

- 1 given $\mathbf{A}\in \boldsymbol{\mathcal{A}},$ to relate $HU(\mathbf{A})$ to $\mathsf{D}(\mathbf{A})$ and
- 2 given (1), to use it to equip HU(A) with additional structure so as to set up a \mathcal{D} - \mathcal{P} -based duality for \mathcal{A} .

[Concentrate here on (1), and on duality rather than strong duality.]

Davey & Werner's key idea

Find a 1:1 map Δ such that diagram commutes. Then $\varepsilon_{U(\mathbf{A})}$ surjective $\implies e_{\mathbf{A}}$ surjective.

Given a X-morphism $\phi \colon D(\mathbf{A}) \to \mathbf{M}$, we seek a \mathcal{P} -morphism $\Delta(\phi) \colon HU(\mathbf{A}) \to \mathbf{2}$. IF $\phi \mapsto \Delta(\phi)$ is 1:1 then $e_{\mathbf{A}}$ is surjective [elementary!].



Construction of Δ for generalised piggybacking

Let $\Omega = HU(\mathbf{M}) = \mathcal{D}(U(\mathbf{M}), \mathbf{2})$ and $\omega \in \Omega$.

Idea: hope that, we can choose \underline{M} so that for each $A\colon$

- we get a well-defined map Δ_{ω} for each $\omega \in \Omega$;
- the union, Δ of the maps Δ_ω is well defined, has domain U(A) and has the properties required.



Choosing R

For $\omega_i, \omega_2 \in \Omega$, consider the sublattice of U(\mathbf{M}^2):

$$(\omega_1, \omega_2)^{-1} \leqslant) := \{ (a, b) \in \mathbf{M}^2 \mid \omega_1(a) \leqslant \omega_2(b) \}.$$

Let R_{ω_1,ω_2} be the set (possibly empty) of algebraic relations maximal w.r.t. being contained in $(\omega_1,\omega_2)^{-1}(\leqslant)$. Let

$$R = \bigcup \{ R_{\omega_1, \omega_2} \mid \omega_1, \omega_2 \in \Omega \}.$$

Theorem

Let $\mathcal{A} = \mathbb{ISP}(\mathbf{M})$, with \mathbf{M} D-based. Then R yields a duality on \mathcal{A} .

Key ingredient in proof: the maps Φ_{ω} , for $\omega \in \Omega$, jointly map onto $HU(\mathbf{A})$ for any $\mathbf{A} \in \mathcal{A}$. This relies on special properties of Priestley duality.

Fix $\mathbf{A} \in \mathcal{A}$, The dual space $D(\mathbf{A}) = \mathcal{A}(\mathbf{A}, \mathbf{M})$ is viewed as a closed substructure of \mathbf{M}^A and carries relations $r^{D(\mathbf{A})}$, for $r \in R$, obtained by pointwise lifting.

Relating $\mathsf{D}(\mathbf{A})$ and $\mathsf{HU}(\mathbf{A})$

Define \preccurlyeq on $D(\mathbf{A}) \times \Omega$ by $(x, \omega_1) \preccurlyeq (y, \omega_2) \iff (x, y) \in r^{\mathsf{D}(\mathbf{A})}$ for some $r \in R_{\omega_1, \omega_2}$.

Theorem

Let \approx be the equivalence relation $\preccurlyeq \cap \succcurlyeq$, Then \approx is ker Φ , where $\Phi: (x, \omega) \mapsto \omega \circ x$. Moreover, the map $\Psi: (\mathsf{D}(\mathbf{A}) \times \Omega) / \approx \rightarrow \mathsf{HU}(\mathbf{A})$ given by $\Psi([(x, \omega)]) = \omega \circ x$ is well defined and a Priestley space isomorphism. The map $\Psi: (\mathsf{D}(\mathbf{A}) \times \Omega) / \approx \rightarrow \mathsf{HU}(\mathbf{A})$ given by $\Psi([(x, \omega)]) = \omega \circ x$ is well defined and a Priestley space isomorphism.



Example 1: prioritised default bilattices revisited

Product representation —and beyond

 $\mathbb{HSP}(\mathbf{K}_0)$ is the much-studied variety $\underline{\mathbf{D}}$ of distributive bilattices. Such bilattices are **interlaced**: each pair of lattice operations is monotonic with to the other lattice order.

Bilattice theorists' favourite theorem: Every interlaced bilattice with negation decomposes as a product $\mathbf{L} \odot \mathbf{L}$, where \mathbf{L} is a lattice, and the truth operations are those of $\mathbf{L} \times \mathbf{L}$, the knowledge operations are those of $\mathbf{L} \times \mathbf{L}^{\partial}$ and $\neg(a.b) = (b, a)$. In the (bounded) distributive case, the product representation and our dualities are tied together.



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Beyond \mathbf{K}_0 : duality to the rescue!

For n > 0, the bilattice \mathbf{K}_n is NOT interlaced. Product representation FAILS and bilattice theory provides no structure theorem.

smallskip But in the natural duality for $\mathbb{HSP}(\mathbf{K}_n)$, we have we have an explicit description of the dual category $\mathfrak{X} = \mathbb{IS}_{c}\mathbb{P}^{+}(\mathfrak{M})$, where

$$\mathfrak{M} = \langle K_0 \cup \ldots \cup K_n; \{S_{m,m}\}_{0 \leq m \leq n}, \{h_{i,i-1}\}_{1 \leq i \leq n}, \mathfrak{T} \rangle,$$

Loosely, the objects in \mathfrak{X} are stratified: each of its (n+1) layers is a Priestley space, and successive layers linked by the maps $HU(h_{i,i-1})$.

This gives rise to a generalisation of the product

representation, whereby each $\mathbf{A} \in \mathbb{HSP}(\mathbf{K}_n)$ can be realised as sitting inside a product of distributive bilattices $(\mathbf{L}_0 \odot \mathbf{L}_0) \times \cdots \times (\mathbf{L}_n \odot \mathbf{L}_n).$

Example 2: Gödel algebras, $G_n = \mathbb{ISP}(n)$

Just a glimpse, to illustrate a point.

The quasivariety generated by the *n*-element Heyting chain *n* Is endodualisable (Davey (1976)); optimised version by Davey & Talukder (2005). Historically, it has been a mantra: 'Use endomorphisms as far as possible; avoid partial operations except where unavoidable.'

BUT G_n also has a simple piggyback duality using graphs of partial endomorphisms (Davey & Werner).

We shall deviate from our general strategy above: we dualise via simple piggybacking, with one endomorphism and two selected piggyback relations, to give a clear picture. This makes the translation from natural to \mathcal{D} - \mathcal{P} -based duality transparent. For an algebra $\mathbf{A} \in \mathcal{G}_n$, we can visualise D) \mathbf{A}) as having at most (n-1) layers: collapsing to HU(\mathbf{A}) occurs within layers, and a single endomorphism gives the ordering between layers.

Application (Cabrer & Priestley, work in progress) : description of coproducts of finite algebras in \mathcal{G}_n .

Illustration: from D(5) to HU(5)Elements of D(5) are uniquely determined by, and so labelled by,

Elements of $D(\mathbf{5})$ are uniquely determined by, and so labelled by, their ranges.

These relations act on D(5) as shown: endomorphism (black); piggyback relations (blue and orange, with converses shown dashed). The transitive closure of the union of these relations gives the quasi-order \approx .



Coproducts via duality

Cabrer and Priestley, 2012)

Assume $\mathcal{A} = \mathbb{ISP}(\mathfrak{M})$ is a finitely generated \mathcal{D} -based quasivariety. Take $\Omega \subseteq \bigcup_{\mathbf{M} \in \mathfrak{M}} \mathsf{HU}(\mathbf{M})$ to satisfy

 $(\text{Sep})_{\mathfrak{M},\Omega}: \text{ for all } \mathbf{M} \in \mathfrak{M}, \text{ given } a, b \in \mathbf{M} \text{ with } a \neq b, \text{ there exists} \\ \mathbf{M}_{a,a} \in \mathfrak{M}, \ u \in \mathcal{A}(\mathbf{M}, \mathbf{M}_{a,b}) \& \omega \in \Omega \cap \mathsf{HU}(\mathbf{M}_{a,b}) \text{ such} \\ \text{ that } \omega(u(a)) \neq \omega(u(b)).$

Ω = Ü HU(M) always works

 —as we used earlier for piggybacking when M = {M}.

 Simple piggybacking needs Sep_{M.ω}: 'one M, one ω' case.

Given $\mathfrak{K} \subseteq \mathcal{A}$, there is a \mathfrak{D} -homomorphism

$$\chi_{\mathbf{\mathfrak{K}}}:\ \coprod \mathsf{U}(\mathbf{\mathfrak{K}}) \to \mathsf{U}(\coprod \mathbf{\mathfrak{K}}),$$

We can work with the dual map $H(\chi_{\mathfrak{K}})$ to get iff conditions (E) for $\chi_{\mathfrak{K}}$ to be **injective** (for any \mathfrak{K});

(S) $\chi_{\mathfrak{K}}$ to be **surjective** (for any \mathfrak{K}).

Coproduct embedding and surjectivity theorems Theorem

The following are equivalent to (E):

- (1) there exists $\mathbf{M} \in \mathcal{A}_{\mathrm{fin}}$ and $\omega \in \mathsf{HU}(\mathbf{M})$ such that $\mathcal{A} = \mathbb{ISP}(\mathbf{M})$ and $(\mathrm{Sep})_{\mathbf{M},\omega}$ holds;
- (2) as (1) but with, additionally, M subdirectly irreducible.

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- (1) there exists $\mathbf{M} \in \mathcal{A}_{\mathrm{fin}}$ and $\omega \in \mathsf{HU}(\mathbf{M})$ such that $\mathcal{A} = \mathbb{ISP}(\mathbf{M})$ and $(\mathrm{Sep})_{\mathbf{M},\omega}$ holds;
- (2) as (1) but with, additionally, M subdirectly irreducible.

Theorem

The following are equivalent to (S):

(1) for every *n*-ary \mathcal{A} -term t $(n \ge 1) \exists$ unary \mathcal{A} -terms t_1, \ldots, t_n and an *n*-ary \mathcal{D} -term s such that for every $\mathbf{A} \in \mathcal{A}$ and every $a_1, \ldots, a_n \in \mathbf{A}$

$$t^{\mathbf{A}}(a_1,\ldots,a_n) = s^{\mathbf{A}}(t_1^{\mathbf{A}}(a_1),\ldots,t_n^{\mathbf{A}}(a_n));$$

- (2) $|R_{\omega_1,\omega_2}| \leqslant 1$ for every $\omega_1, \omega_2 \in \Omega$;
- (3) for $\mathbf{A} \in \mathcal{A}$ and every \mathfrak{D} -sublattice \mathbf{L} of $U(\mathbf{A})$, then either L contains no \mathcal{A} -subalgebra of $\{ \mathbf{B} \in \mathbb{S}(\mathbf{A}) \mid B \subseteq L \}$ is empty has a top element.

A medley of examples

E	Boolean Algebras De Morgan algebras <i>n</i> -valued pre-Łukasiewicz–Moisil algebras Stone algebras
E✓S×	Gödel algebra varieties \mathcal{G}_n $(n \ge 3)$ Varieties \mathcal{B}_n $(2 \le n < \omega)$ Q -lattice varieties \mathcal{D}_{p0} and \mathcal{D}_{q1} $(p \ge 1, q \ge 0)$
E×S✓	Kleene algebras MV-algebra varieties $\not\supseteq \mathbf{L}_{p \cdot q}$ (whenever $p \neq q$, primes) <i>n</i> -valued Łukasiewicz–Moisil algebras ($n \ge 2$)
$E^{\times}S^{\times}$	Non-singly generated varieties of Heyting algebras MV-algebra varieties $\supseteq \mathbf{L}_{pq}$ (for some $p \neq q$, primes) Q -lattice varieties \mathbf{D}_{pq} $(q \ge 2)$

Coproducts and piggybacking

Theorem

Let \mathcal{A} be a finitely generated \mathfrak{D} -based quasivariety.

- (i) \mathcal{A} admits free products iff it admits a single-sorted duality.
- (ii) U satisfies (E) iff \mathcal{A} admits a simple piggyback duality.
- (iii) U satisfies (S) iff \mathcal{A} admits a piggyback duality (singlesorted or multisorted), such that $|R_{\omega_1,\omega_2}| \leq 1$ for all $\omega_1, \omega_2 \in HU(\mathbf{M})$.
- (iv) U preserves coproducts (hat is, (E) and (S) hold) iff \mathcal{A} has a simple piggyback duality that is a \mathcal{D} - \mathcal{P} -based duality.

Familiar classics: reconciliations

- (iv) holds for De Morgan algebras and for Stone algebras.
- Can handle coproducts of Kleene algebras via a reflector from De Morgan algebras. This strategy generalises.

How to sum up?

Brian Davey has gone a very long way in convincing algebraists that topology and algebra make good bed-fellows.

In this he has followed in the footsteps of Marshall Stone (1938): A cardinal principle of modern mathematical research may be stated as a maxim: "One must always topologize."



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Pathways, Patios, Walls and More.

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