

Thm If G is f.g. and non-amenable,
then $\delta_G \geq e^n$.

Proof

$$G = \langle S \rangle, |S| \in \mathbb{O}.$$

G is non-amenable, $\Rightarrow \exists \epsilon > 0$:

$$|\partial B_S(1, n)| \geq \epsilon |B_S(1, n)|$$

Clearly $\epsilon < 1$, and

$$(1 - \epsilon) |B_S(1, n)| \geq |B_S(1, n-1)|$$

$$\text{So } \delta_{G,S}(n) \geq \left(\frac{1}{1 - \epsilon} \right)^n \text{ exponential. } \square$$

\Leftarrow
Groups of subexponential growth
are amenable \Leftarrow

Van Neumann - Day Problem

Is the converse true? If you have no \mathbb{Z}_2 ,
are you a wreath?

No. (Olshanskii 1980) constructed

non-amenable Tarski monster,
i.e. all subgroups are \mathbb{Z} or \mathbb{Z}_2 .

Def A group is a Tarski monster
if it is non-cyclic, proper non-trivial
subgroups are cyclic and isomorphic.

[A bit like \mathbb{Q} , but f.g.]

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The (Jordanho-Mass)

\exists infinite simple abelian groups.

In particular, not every abelian group admits a vector space V .

1.3 Wolf-Milnor Theorem

The (Milnor-Wolf 1968)

Let G be a f.g. group.

1) If G is vit. nilpotent then

δ_G is polynomial

2) If G is vit. solvable but not v.

nilpotent then γ_6 is exponentiated.

We need some preliminaries.

Lemma

G f.g., $N \trianglelefteq G$ of finite index.

Then $\gamma_n \times \gamma_6$.

Proof.

Let S be a finite gen set of N ,

$$\bar{T} = S \cup S' \quad \text{---} \quad \text{---} \quad G$$

with $G \rightarrow G/N$ taking S' to G/N
bijectively.

Let $k \in \mathbb{N}$ be such that $S^{S'} =$

$$\downarrow \circlearrowleft^t : t \in S', \circlearrowleft^t \subseteq B_S(1, 4).$$

$$\circlearrowleft^{-t} \text{ and } S \cdot S' \subseteq B_S(1, 4) \mid S'.$$

$$\delta_{N, S}(m) = |B_S(1, m)| \leq |B_T(1, m)|$$

\downarrow
 $\delta_{6, T}(m).$

$$\Rightarrow \delta_N \leq \delta_6.$$

Consider a word in \bar{T} . Going from

the left, we replace

$$t \succ \dots \succ t', \quad t, t' \in S', \quad t \in S$$

by $t \succ \dots \succ t^{-1} \succ t''$ where

$$\succ t'' = t^{-1} t', \quad t'' \in S', \quad t'' \in S.$$

Hence every element in $B_T(1, m)$
 can be written as a word of length at
 most mk in S times a single element
 from S' .

There are at most $\gamma_S(mk) \cdot |G'|$
 such elements in G' and so

$$\delta_G \leq \delta_{G'}.$$

Cor $H \leq G$, G f.g., then

$$\delta_H \leq \delta_G.$$

Proof $\exists N \trianglelefteq G$, $N \trianglelefteq H$. \square

and \times is transitive.

Lemma

If G f.g. contains the free monoid
on two generators $M(a, b)$ then

$$\gamma_G \cong e^n.$$

Proof.

$$\gamma_G \leq e^n \Rightarrow \text{ETS } e^n \leq \gamma_G.$$

$$\exists h: a, b \in B_G(1, h).$$

$$\text{So } \gamma_G(mh) \geq \beta_M(1, m) \cong e^m.$$

$$\gamma_G(m) \geq e^{\lfloor \frac{m}{h} \rfloor} \geq e^{\frac{m}{h} - 1} \cong e^m.$$

□

Preparations for nilpotent groups.

Conventions: $g, h \in G$.

$$g^h = \bar{h}^{-1} g h = \bar{h} g h$$

$$[g, h] = g h \bar{g} \bar{h}.$$

Very useful commutator relation:

$$[gh, x] = [g, x][h, x]^{x \bar{g} \bar{x}}.$$

In particular, if $[h, x]$ is central, then

$$[gh, x] = [g, x][h, x], \quad x \in G.$$

Key Lemma

G nilpotent of step k .

(G_i) Lower central series, $G_0 = G$

G/G_i l.g. $\varphi_i: G_i \rightarrow G_i/G_{i+1}$ $G_{k+1} = 1$.

\exists finite set $T_i \subseteq G_i$ s.t.:

1) $\varphi_i(T_i)$ is a minimal gen set of the abelian group G_i/G_{i+1}

2) $\forall i$, T_i consists of commutators from

$$[T_0, T_{i-1}] = \{ (t_0, t_{i-1}) : t_0 \in T_0 \wedge t_{i-1} \in T_{i-1} \}$$

Proof Induction on i .

Base case : $i=0$. just pick T_0

Induction step : we have picked T_0, \dots, T_{i-1} .

Consider $\rho: G \rightarrow H = G/G_{i+1}$

Let $H_i = \rho(G_i)$.

Let $S_j = \rho(T_j)$; note S_j has the same properties.

Key: H_i is central in H .

Every element of H_i is a product of commutators (α, β) with $\alpha \in H, \beta \in H_{i-1}$.

Since H_i is central, we may modify
every α, β by multiplying it by elements of
 H_i . Since S_0 generates H/H_i

(as T_0 generates G/A_i), we may assume
that α is a word in S_0 .

Similarly, β is a word in S_{i-1} .

$$\alpha = x_1 \dots x_n, \quad x_i \in S_0$$

$$\beta = y_1 \dots y_b, \quad y_i \in S_{i-1}$$

$$[\alpha, \beta] = \prod_{i,j} [x_i, y_j] \text{ by the}$$

commutator identity.

S_0 , $[S_0, S_{i-1}]$ generate H_i .

Pick a minimal subset, call S_i , put

back by ρ to form T_i . \square