

# Max $k$ -cut and judicious $k$ -partitions

Béla Bollobás <sup>\*†‡</sup>      Alex Scott <sup>§</sup>

## Abstract

Alon, Bollobás, Krivelevich and Sudakov [1] proved that every graph with a large cut has a bipartition in which each vertex class contains correspondingly few edges. We prove an analogous result for partitions into  $k \geq 3$  classes; along the way we prove a result for biased bipartitions.

## 1 Introduction

Let  $G$  be a graph with  $m$  edges. It is easy to show that  $G$  has a cut (or, equivalently, a bipartite subgraph) of size least  $m/2$ . It is much less obvious (but nevertheless true) that there is a cut of this size such that the remaining edges are roughly evenly distributed between the two sides of the cut: in other words, each vertex class contains no more than (roughly)  $m/4$  edges. Now suppose that  $G$  has a cut that is much larger than  $m/2$ . In this case we might hope for more: if  $G$  has a cut of size  $m/2 + \alpha$ , then a near-optimal cut that divides the remaining edges roughly equally between the two vertex classes would have roughly  $m/4 - \alpha/2$  edges in each class. Alon, Bollobás, Krivelevich and Sudakov [1] showed that, for  $\alpha$  not too large, this is indeed possible (for  $\alpha$  large, they proved a complementary result: if  $\alpha \geq m/30$ , there is a bipartition in which each class contains at most  $m/4 - m/100$  edges).

---

<sup>\*</sup>Department of Pure Mathematics and Mathematical Statistics, University of Cambridge, Cambridge CB3 0WB, UK and

<sup>†</sup>Department of Mathematical Sciences, University of Memphis, Memphis TN 38152, USA

<sup>‡</sup>Research supported in part by NSF grants DMS-0505550, CNS-0721983 and CCF-0728928, and ARO grant W911NF-06-1-0076

<sup>§</sup>Mathematical Institute, 24-29 St Giles', Oxford, OX1 3LB, UK; email: scott@maths.ox.ac.uk

The aim of this paper is to generalize these results in two directions: we first give results on “biased” cuts, in which edges in the two vertex classes are counted with different weights; we then continue by giving results in partitions into more than two parts. In each case, as with Alon, Bollobás, Krivelevich and Sudakov [1], we obtain matching results for the cases  $\alpha$  small and  $\alpha$  large.

The remainder of this introduction is divided into two parts. In the first part, we discuss some background to the problem; the second part describes our results and gives a little notation.

## 1.1 Previous work

For a graph  $G$ , let us define

$$f(G) = \max_{V(G)=V_1 \dot{\cup} V_2} e(V_1, V_2) = \max_{V(G)=V_1 \dot{\cup} V_2} (m - e(V_1) - e(V_2))$$

to be the maximum size of a cut in  $G$ . Then, for  $m \geq 1$ , we set

$$f(m) = \min_{e(G)=m} f(G).$$

The extremal Max Cut problem asks for the value of  $f(m)$ , and has been extensively studied. It is easy to see that  $f(m) \geq m/2$ , for instance by considering random partitions or a suitable greedy algorithm. Edwards [10, 11] showed that

$$f(m) \geq \frac{m}{2} + \sqrt{\frac{m}{8} + \frac{1}{64}} - \frac{1}{8}, \tag{1}$$

which is sharp for complete graphs of odd order. More precise bounds for other values of  $m$  were given by Alon [2], Alon and Halperin [3], and in [5]. From the other side, it is easily seen by considering random graphs  $G \in \mathcal{G}(n, 1/2)$  that  $f(m) = m/2 + o(m)$ .

The Max Cut problem asks for a bipartition in which  $e(V_1, V_2)$  is large, and hence  $e(V_1) + e(V_2)$  is small. However, it does not place strong constraints on the number of edges in each vertex class separately. Problems in which constraints are placed on all vertex classes simultaneously are known as *judicious partitioning problems* (see [16] and [4] for an overview). In this case, we define a judicious partitioning problem as follows. For a graph  $G$ , let

$$g(G) = \min_{V(G)=V_1 \dot{\cup} V_2} \max\{e(V_1), e(V_2)\},$$

and, for  $m \geq 1$ , set

$$g(m) = \max_{e(G)=m} g(G).$$

Determining the behaviour of  $g(m)$  seems significantly harder than analyzing  $f(m)$ . For instance, proving that  $f(m) \sim m/2$  is trivial, but there does not seem to be any simple way to prove that  $g(m) \sim m/4$  (which turns out to be true). Bounds on  $g(m)$  were proved by several authors, including Porter [12, 13, 14], Porter and Bin Yang [15], and Bollobás and Scott [9]. An analogue of the Edwards bound was finally proved in [7], where it was shown that every graph  $G$  with  $m$  edges has a bipartition  $V(G) = V_1 \cup V_2$  such that

$$\max\{e(V_1), e(V_2)\} \leq \frac{m}{4} + \sqrt{\frac{m}{32} + \frac{1}{256}} - \frac{1}{16} \quad (2)$$

and in addition  $e(V_1, V_2)$  satisfies (1). More generally, there is a vertex partition into  $k$  classes, each of which contains at most

$$\frac{m}{k^2} + \frac{k-1}{2k^2} \left( \sqrt{2m + \frac{1}{4}} - \frac{1}{2} \right) \quad (3)$$

edges.

The bounds (2) and (1) are closely related, and it is natural to ask whether graphs with a very large cut (i.e. much larger than that guaranteed by (1)) also have a correspondingly good judicious partition. If  $G$  is a graph with  $m$  edges, and  $f(G) = m/2 + \alpha$ , then it is clear that  $g(G) \geq m/4 - \alpha/2$ , since we cannot do better than a maximum cut with the remaining edges divided equally between the two vertex classes. Alon, Bollobás, Krivelevich and Sudakov [1] showed that it is possible to get pretty close to this bound: if  $\alpha \leq m/30$  then

$$g(G) \leq \frac{m}{4} - \frac{\alpha}{2} + 3\sqrt{m} + \frac{10\alpha^2}{m}. \quad (4)$$

For large  $\alpha$ , this bound is less useful. However, they also showed the complementary result that if  $\alpha \geq m/30$  (and  $m$  is sufficiently large) then

$$g(G) \leq \frac{m}{4} - \frac{m}{100}. \quad (5)$$

## 1.2 Our results

The aim of this paper is to extend the results of Alon, Bollobás, Krivelevich and Sudakov [1] in two directions: to biased partitions, and to partitions into  $k \geq 3$  parts.

In Section 2, we give results on biased partitions. For  $p \in [0, 1]$  and  $q = 1 - p$ , define

$$m_p(G) = \min_{V(G)=V_1 \dot{\cup} V_2} qe(V_1) + pe(V_2)$$

Note that this is a ‘biased’ generalization of Max Cut: if we take  $p = 1/2$  then we get  $m_{1/2}(G) = \frac{1}{2}(m - f(G))$ .

Considering a random bipartition where each vertex independently has probability  $p$  of being in  $V_1$ , we get  $\mathbb{E}e(V_1) = p^2m$  and  $\mathbb{E}e(V_2) = q^2m$ . It follows that every graph  $G$  with  $m$  edges has  $m_p(G) \leq pqm$ , while complete graphs or not too sparse random graphs show that we can have  $m_p(G) = (1 + o(1))pqm$ . A corresponding judicious result was proved in [7], where it was shown that there is in fact a bipartition such that there are no more than about  $p^2m$  edges in  $V_1$  and  $q^2m$  edges in  $V_2$ . More precisely, there is a bipartition in which

$$e(V_1) \leq p^2m + h(p, m) \tag{6}$$

and

$$e(V_2) \leq q^2m + h(p, m), \tag{7}$$

where

$$h(p, m) = pq(\sqrt{m/2 + 1/16} - 1/4).$$

Note that when  $p = 1/2$ , we recover (2).

Our aim in section 2 is to prove bounds similar to (4) and (5) in this context. Suppose that  $m_p(G) = pqm - \alpha$ . If  $\alpha \leq c(p)m$ , we will show in Theorem 1 that there is a bipartition  $V(G) = V_1 \cup V_2$  such that  $V_1$  and  $V_2$  satisfy inequalities of form

$$e(V_1) \leq p^2m - \alpha + O(\sqrt{m} + \frac{\alpha^2}{m}) \tag{8}$$

and

$$e(V_2) \leq q^2m - \alpha + O((\sqrt{m} + \frac{\alpha^2}{m})). \tag{9}$$

Note that we get  $\alpha$  rather than  $\alpha/2$  here: this reflects the definition of  $m_p(G)$ : for example, compare  $m_{1/2}(G)$  with  $f(G)$ .

If  $\alpha \geq c(p)m$  then (8) and (9) are no longer useful: we show in Theorem 3 that there is a bipartition  $V(G) = V_1 \cup V_2$  such that

$$e(V_1) \leq p^2m - c^*(p)m$$

and

$$e(V_2) \leq q^2m - c^*(p)m.$$

More precise statements of these results can be found at Theorems 1 and 3 below.

In Section 3, we turn to partitions into more than 2 pieces. For  $k \geq 2$ , let us define  $\text{mc}_k(G)$  to be the maximum size of a  $k$ -cut of  $G$ . It is easily seen by considering a random partition that every graph  $G$  with  $m$  edges has

$$\text{mc}_k(G) \geq \frac{k-1}{k}m.$$

We show (Theorem 5) that if there is a significantly larger cut then we get a very good judicious partition. If

$$\text{mc}_k(G) = \frac{k-1}{k}m + \alpha$$

then the following holds: if  $\alpha \leq c(k)m$  then there is a  $k$ -cut in which each class has at most

$$\frac{m}{k^2} - \frac{\alpha}{k} + O\left(\sqrt{m} + \frac{\alpha^2}{m}\right) \quad (10)$$

edges (once again, a more precise statement is given below). For  $\alpha > c(k)m$  there is (Theorem 8) a  $k$ -cut in which each class has at most  $m/k^2 - c^*(k)m$  edges. Note that if  $\alpha$  is not too large, then (10) is similar to (3), except for the constant in the error term.

In both sections, our proof strategy is to start with a good biased partition or  $k$ -cut and then move vertices one at a time out of a ‘bad’ vertex class while tracking their effect on the distribution of edges. This was used in [7] and refined in [1]. Our strategy is similar to the approach used in [1]. However, there are some additional obstacles that need to be overcome.

Throughout the paper, we use the following notation. Let  $G$  be a graph. For  $W \subset V(G)$ , we write  $e(W)$  for the number of edges spanned by  $W$ ; for disjoint  $X, Y \subset V(G)$  we write  $e(X, Y)$  for the number of edges  $xy \in E(G)$  with  $x \in X$  and  $y \in Y$ .

## 2 Biased partitions

Let  $G$  be a graph with  $m$  edges and  $p \in [0, 1]$ ,  $q = 1 - p$ . In this section, we consider partitions  $V(G) = V_1 \cup V_2$  that minimize  $qe(V_1) + pe(V_2)$ . Recall that

$$m_p(G) = \min_{V(G)=V_1 \cup V_2} qe(V_1) + pe(V_2).$$

For a random partition in which each vertex independently is placed in  $V_1$  with probability  $p$  or in  $V_2$  with probability  $q$ , we have  $\mathbb{E}(qe(V_1) + pe(V_2)) = pqm$ . We shall show that if  $m_p(G) = pqe(G) - \alpha$ , with  $\alpha \gg \sqrt{m}$ , then we get a very good judicious partition.

Note that in a partition with  $qe(V_1) + pe(V_2)$  minimal, every  $v \in V_1$  must satisfy

$$q|\Gamma(v) \cap V_1| \leq p|\Gamma(v) \cap V_2|, \quad (11)$$

or else we would have moved  $v$  to  $V_2$ , and a similar inequality holds for vertices in  $V_2$ . We shall refer to (11) as the *local inequality*.

For any partition  $V(G) = V_1 \cup V_2$  that satisfies the local inequality, summing over  $V_1$  implies that

$$e(V_1, V_2) \geq \frac{2q}{p}e(V_1)$$

and so

$$\begin{aligned} e(V_2) &= m - e(V_1) - e(V_1, V_2) \\ &\leq m - e(V_1) - \frac{2q}{p}e(V_1) \\ &= m - \frac{1+q}{p}e(V_1). \end{aligned}$$

Therefore

$$\begin{aligned} qe(V_1) + pe(V_2) &\leq qe(V_1) + pm - (1+q)e(V_1) \\ &= pm - e(V_1). \end{aligned}$$

Thus if  $m_p(G) = pqm - \alpha$ , and  $V_1$  and  $V_2$  satisfy the local inequality, we have

$$e(V_1) \leq p^2m + \alpha. \quad (12)$$

We begin with a result for  $\alpha$  of moderate size, and prove a result for large  $\alpha$  later (Theorem 3).

**Theorem 1.** Let  $0 < p < 1$ ,  $q = 1 - p$ , and let  $c(p) = \frac{1}{2} \min\{p^2, q^2\}$ . Suppose  $G$  is a graph with  $m$  edges such that

$$m_p(G) = pqm - \alpha, \quad (13)$$

where  $\alpha \leq c(p)m$ . Then there is a partition  $V(G) = V'_1 \cup V'_2$  such that

$$e(V'_1) \leq p^2m - \alpha + \sqrt{32mp^2} + \frac{16\alpha^2}{q^3m} \quad (14)$$

and

$$e(V'_2) \leq q^2m - \alpha + \sqrt{32mq^2} + \frac{16\alpha^2}{p^3m}. \quad (15)$$

Note that this improves on (6) and (7) only in the range  $\alpha = O(\min\{p^3, q^3\})m$ .

Our main tool in the proof of Theorem 1 is the following.

**Lemma 2.** Suppose  $G$  has  $m$  edges and satisfies (13), where  $\alpha \leq p^2m/2$ . Suppose  $W \subset V = V(G)$  and, for all  $v \in W$ ,

$$|\Gamma(v) \cap V \setminus W| \geq \frac{q}{p} |\Gamma(v) \cap W|. \quad (16)$$

If  $e(W) > p^2m - \alpha$  then there is  $v \in W$  with

$$|\Gamma(v) \cap W| \leq \sqrt{32mp^2} \quad (17)$$

and

$$|\Gamma(v) \cap V \setminus W| \leq \left( \frac{q}{p} + \frac{8\alpha}{p^3m} \right) |\Gamma(v) \cap W|. \quad (18)$$

As above, we will refer to inequality (16) as the *local inequality*.

*Proof.* Define

$$T_1 = \{v \in W : |\Gamma(v) \cap W| > \sqrt{32mp^2}\} \quad (19)$$

and

$$T_2 = \{v \in W : |\Gamma(v) \cap V \setminus W| > \left( \frac{q}{p} + \frac{8\alpha}{p^3m} \right) |\Gamma(v) \cap W|\}. \quad (20)$$

Summing the inequality satisfied by vertices in (20) over  $T_2$ , and summing (16) over the rest of  $W$ , we see that

$$\begin{aligned} e(W, V \setminus W) &\geq \frac{q}{p} \sum_{v \in W} |\Gamma(v) \cap W| + \frac{8\alpha}{p^3m} \sum_{v \in T_2} |\Gamma(v) \cap W| \\ &= \frac{2q}{p} e(W) + \frac{8\alpha}{p^3m} \sum_{v \in T_2} |\Gamma(v) \cap W|. \end{aligned}$$

Thus

$$\begin{aligned}
qe(W) + pe(V \setminus W) &= qe(W) + p(m - e(W) - e(W, V \setminus W)) \\
&\leq qe(W) + p \left( m - e(W) - \frac{2q}{p}e(W) - \frac{8\alpha}{p^3m} \sum_{v \in T_2} |\Gamma(v) \cap W| \right) \\
&= pm - e(W) - \frac{8\alpha}{p^2m} \sum_{v \in T_2} |\Gamma(v) \cap W| \\
&< pqm + \alpha - \frac{8\alpha}{p^2m} \sum_{v \in T_2} |\Gamma(v) \cap W|.
\end{aligned}$$

Thus, by (13),

$$\frac{8\alpha}{p^2m} \sum_{v \in T_2} |\Gamma(v) \cap W| < 2\alpha$$

and so

$$\sum_{v \in T_2} |\Gamma(v) \cap W| < \frac{p^2m}{4}. \quad (21)$$

On the other hand, since  $W$  and  $V \setminus W$  satisfy the local inequality, by (12) we have

$$e(W) \leq p^2m + \alpha \leq 2p^2m$$

and so

$$\sum_{v \in T_1} |\Gamma(v) \cap W| \leq 2e(W) \leq 4p^2m,$$

which, by the definition of  $T_1$ , implies

$$|T_1| \leq \frac{4p^2m}{\sqrt{32mp^2}} = \sqrt{\frac{mp^2}{2}}.$$

Thus

$$\sum_{v \in T_1} |\Gamma(v) \cap W| \leq e(T_1) + e(W) \leq \binom{|T_1|}{2} + e(W) \leq \frac{mp^2}{4} + e(W). \quad (22)$$

Since  $e(W) > p^2m - \alpha \geq p^2m/2$ , (22) and (21) give

$$\sum_{v \in T_1 \cup T_2} |\Gamma(v) \cap W| < \frac{p^2m}{2} + e(W) \leq 2e(W)$$

and so  $T_1 \cup T_2 \neq W$ . The lemma follows immediately.  $\square$



We can now turn to the proof of Theorem 1.

*Proof of Theorem 1.* Let  $V_1 \cup V_2$  be a partition with

$$qe(V_1) + pe(V_2) = pqm - \alpha.$$

If (14) and (15) are satisfied for  $V_1$  and  $V_2$ , we are done. Otherwise, exchanging  $p$  and  $q$  if necessary (and noting that this also exchanges (14) and (15)), we may assume that

$$e(V_1) > p^2m - \alpha.$$

If

$$e(V_1) = p^2m - \alpha + \lambda$$

then

$$\begin{aligned} pe(V_2) &= pqm - \alpha - qe(V_1) \\ &= pqm - \alpha - qp^2m + q\alpha - q\lambda \\ &= pq^2m - p\alpha - q\lambda \end{aligned}$$

and so

$$e(V_2) = q^2m - \alpha - \frac{q}{p}\lambda. \quad (23)$$

Note that  $(V_1, V_2)$  satisfies the local inequality (16) (with  $W = V_1$ ).

We now successively move vertices from  $V_1$  to  $V_2$ , at each step choosing a vertex satisfying (17) and (18). We can find such a vertex, as the local inequality (16) remains true if we remove vertices from  $V_1$  and so we can apply Lemma 2. We continue until we obtain  $V'_1$  such that  $p^2m - \alpha \leq e(V'_1) \leq p^2m - \alpha + \sqrt{32mp^2}$  (note that (17) guarantees that our steps are sufficiently small that we don't overshoot). Since we have decreased  $e(V_1)$  by at most  $\lambda$ , (18) implies that we have increased  $e(V_2)$  by at most

$$\left( \frac{q}{p} + \frac{8\alpha}{p^3m} \right) \lambda$$

and so, by (23), we end up with  $V'_2$  satisfying

$$\begin{aligned} e(V'_2) &\leq e(V_2) + \left( \frac{q}{p} + \frac{8\alpha}{p^3m} \right) \lambda \\ &\leq q^2m - \alpha + \frac{8\alpha}{p^3m}\lambda. \end{aligned} \quad (24)$$

By (12) we have  $\lambda \leq 2\alpha$ , and so the result follows from (24) by taking the partition  $(V'_1, V'_2)$ .  $\square$

We now deal with the case when  $\alpha$  is large.

**Theorem 3.** *Let  $0 < p < 1$  and  $q = 1 - p$ . Let  $0 < c < \min\{p^2, q^2\}$  and  $c^*(p) = \min\{cp/12, cq/12\}$ . Suppose that  $G$  is a graph with  $m$  edges and*

$$m_p(G) = pqm - \alpha, \quad (25)$$

where  $\alpha \geq cm$ . Then, provided that  $m$  is sufficiently large (in terms of  $c$  and  $p$ ), there is a partition  $V(G) = V_1 \cup V_2$  such that

$$e(V_1) \leq p^2m - c^*m \quad (26)$$

$$e(V_2) \leq q^2m - c^*m. \quad (27)$$

The best fit with Theorem 1 is obtained by specializing to a particular value of  $c$ . However, it will be useful in the next section to allow any  $c > 0$ .

The proof of Theorem 3 is based on the following lemma.

**Lemma 4.** *Let  $0 < p < 1$ ,  $q = 1 - p$ , and suppose that  $0 < c^* < p^3/9$ . Suppose that  $G$  is a graph with  $m$  edges. Suppose  $W \subset V = V(G)$  satisfies*

$$e(W) > p^2m - c^*m$$

and, for every  $w \in W$ ,

$$q|\Gamma(w) \cap W| \leq p|\Gamma(w) \cap V \setminus W|.$$

Then, provided  $m$  is sufficiently large (in terms of  $p$  and  $c^*$ ), either there is  $w \in W$  such that

$$|\Gamma(w) \cap W| < c^*m \quad (28)$$

and

$$|\Gamma(w) \cap V \setminus W| \leq \left(\frac{q}{p} + \frac{1}{2}\right)|\Gamma(w) \cap W|, \quad (29)$$

or there is  $W' \subset W$  such that  $(V_1, V_2) = (W', V \setminus W')$  satisfies

$$e(V_1) \leq p^2m - c^*m \quad (30)$$

$$e(V_2) \leq q^2m - c^*m. \quad (31)$$

*Proof.* Let

$$T_1 = \{v \in W : |\Gamma(v) \cap W| > c^*m\}$$

and

$$T_2 = \{v \in W : |\Gamma(v) \cap V \setminus W| \geq \left(\frac{q}{p} + \frac{1}{2}\right) |\Gamma(v) \cap W|\}.$$

Let  $T = T_1 \cup T_2$ . We consider two cases.

*Case 1.*  $e(T) \geq p^2m - c^*m$ .

Since every graph with  $m$  edges has a vertex of degree at most  $\sqrt{2m}$ , we can delete vertices from  $T \subseteq W$  one at a time until we obtain  $V_1 \subseteq T$  with

$$p^2m - c^*m - \sqrt{2m} < e(V_1) \leq p^2m - c^*m. \quad (32)$$

Then, writing  $V_2 = V \setminus V_1$ , and using the local inequality and the fact that  $V_1 \subseteq T$ , we have

$$\begin{aligned} e(V_1, V_2) &= \sum_{v \in V_1} |\Gamma(v) \cap V_2| \\ &\geq \sum_{v \in V_1} |\Gamma(v) \cap V \setminus T| \\ &\geq \frac{q}{p} \sum_{v \in V_1} |\Gamma(v) \cap T| + \frac{1}{2} \sum_{v \in V_1 \cap T_2} |\Gamma(v) \cap T| \\ &\geq \frac{2q}{p} e(V_1) + \frac{1}{2} \sum_{v \in V_1 \cap T_2} |\Gamma(v) \cap V_1|. \end{aligned} \quad (33)$$

Now, since  $T \subseteq W$ ,  $\sum_{v \in T_1} |\Gamma(v) \cap W| \leq 2e(W) \leq 2m$ , and so  $|T_1| < 2/c^*$ . Thus  $e(T_1) < 2/(c^*)^2$ , and so

$$\begin{aligned} \sum_{v \in V_1 \cap T_1} |\Gamma(v) \cap V_1| &\leq e(V_1 \cap T_1) + e(V_1) \\ &\leq e(T_1) + e(V_1) \\ &< \frac{2}{(c^*)^2} + e(V_1). \end{aligned}$$

Since  $V_1 \subseteq T$ , it follows that

$$\begin{aligned} \sum_{v \in V_1 \cap T_2} |\Gamma(v) \cap V_1| &\geq 2e(V_1) - \sum_{v \in V_1 \cap T_1} |\Gamma(v) \cap V_1| \\ &> e(V_1) - \frac{2}{(c^*)^2} \\ &\geq e(V_1)/2, \end{aligned}$$

provided  $m$  is sufficiently large. Thus, by (33),

$$e(V_1, V_2) > \left( \frac{2q}{p} + \frac{1}{4} \right) e(V_1)$$

and so, using (32),

$$\begin{aligned} e(V_2) &= m - e(V_1) - e(V_1, V_2) \\ &< m - \left( 1 + \frac{2q}{p} + \frac{1}{4} \right) e(V_1) \\ &\leq m - \left( 1 + \frac{2q}{p} + \frac{1}{4} \right) (p^2m - c^*m - \sqrt{2m}) \\ &= q^2m - \frac{1}{4}p^2m + \left( 1 + \frac{2q}{p} + \frac{1}{4} \right) (c^*m + \sqrt{2m}) \\ &< q^2m - c^*m, \end{aligned}$$

provided  $m$  is sufficiently large. Thus  $(V_1, V_2)$  satisfies (30) and (31), as required.

*Case 2.*  $e(T) < p^2m - c^*m$ .

In this case, there is some vertex  $w \in W \setminus T$ ; this vertex will satisfy the required inequalities. □

*Proof of Theorem 3.* Let  $V(G) = V_1 \cup V_2$  be a partition such that

$$qe(V_1) + pe(V_2) = pqm - \alpha. \quad (34)$$

If  $V_1$  and  $V_2$  satisfy (26) and (27) then we are done. Otherwise (exchanging  $V_1$  and  $V_2$  and  $p$  and  $q$  if necessary, and noting that  $c^*$  is unchanged) we may assume  $V_1$  fails (26). Suppose that

$$e(V_1) = p^2m - c^*m + \lambda, \quad (35)$$

so that, by (34),

$$\begin{aligned} e(V_2) &= \frac{1}{p}(pqm - \alpha - qe(V_1)) \\ &= \frac{1}{p}(pq^2m - \alpha - q\lambda + qc^*m) \\ &= q^2m - \frac{q\lambda + \alpha}{p} + \frac{qc^*m}{p} \end{aligned}$$

Provided  $m$  is sufficiently large, we can move vertices from  $V_1$  to  $V_2$  using Lemma 4. At each stage, we either obtain the partition required by the theorem, or by (29) move a vertex that decreases  $e(V_1)$  by some integer  $d$  and increases  $e(V_2)$  by at most  $(\frac{q}{p} + \frac{1}{2})d$ . We halt when we reach  $V'_1 \subset V_1$  with

$$p^2m - 2c^*m \leq e(V'_1) \leq p^2m - c^*m;$$

here, (28) guarantees that we do stop. We have decreased  $e(V_1)$  by

$$e(V_1) - e(V'_1) \leq \lambda + c^*m$$

and so, writing  $V'_2 = V \setminus V'_1$ ,

$$\begin{aligned} e(V'_2) &\leq e(V_2) + \left(\frac{q}{p} + \frac{1}{2}\right)(\lambda + c^*m) \\ &= q^2m - \frac{q\lambda + \alpha}{p} + \frac{qc^*m}{p} + \left(\frac{q}{p} + \frac{1}{2}\right)(\lambda + c^*m) \\ &= q^2m - \frac{1}{p}\alpha + \frac{1}{2}\lambda + \left(\frac{2q}{p} + \frac{1}{2}\right)c^*m. \end{aligned}$$

By (12), (35) and (34), we have  $\lambda \leq \alpha + c^*m$ , so

$$\begin{aligned} e(V'_2) &\leq q^2m - \left(\frac{1}{p} - \frac{1}{2}\right)\alpha + 2\left(\frac{q}{p} + \frac{1}{2}\right)c^*m \\ &\leq q^2m - \frac{1}{2}\alpha + \frac{4}{p}c^*m \\ &< q^2m - c^*m. \end{aligned}$$

Thus  $(V'_1, V'_2)$  will do for our partition. □

### 3 Partitions into $k$ vertex classes

In this section, we show that graphs with a large  $k$ -cut have a good judicious partition into  $k$  vertex classes. As in the previous section, we begin with a result for moderate values of  $\alpha$ , and then prove a result (Theorem 8) for large  $\alpha$ .

Our first result is the following.

**Theorem 5.** Let  $k \geq 2$ . Suppose that  $G$  is a graph with  $m$  edges such that

$$\text{mc}_k(G) = \left(1 - \frac{1}{k}\right) m + \alpha, \quad (36)$$

where  $\alpha \leq m/k^6$ . Then there is a  $k$ -cut in which each class has at most

$$\frac{m}{k^2} - \frac{\alpha}{k} + \frac{k^5 \alpha^2}{m} + 4\sqrt{m} \quad (37)$$

edges.

Before we prove this result, let us make a few simple observations. Note first that if  $\bigcup_{i=1}^k V_i$  is a maximum  $k$ -cut of  $G$  then, for  $i \neq j$  and  $v \in V_i$ , we have

$$|\Gamma(v) \cap V_j| \geq |\Gamma(v) \cap V_i|, \quad (38)$$

or else we could move  $v$  from  $V_i$  to  $V_j$  to obtain a larger cut. Thus every vertex class  $V_i$  satisfies, for all  $v \in V_i$ , the inequality

$$|\Gamma(v) \cap V \setminus V_i| \geq (k-1)|\Gamma(v) \cap V_i|. \quad (39)$$

Once again, we shall refer to this as the *local inequality*.

Summing (38) over vertices in  $V_i$ , we find that

$$e(V_i, V_j) \geq 2e(V_i). \quad (40)$$

It is easily seen (for instance, by considering a random  $k$ -cut, or partitioning greedily one vertex at a time) that

$$\text{mc}_k(G) \geq \frac{k-1}{k} e(G). \quad (41)$$

Given a partition of some subset  $W \subset V(G)$  into  $k$  sets, we can extend greedily to a  $k$ -cut of  $G$  by adding vertices one at a time to whichever class maximizes the partial cut at each step. We see that, if  $H = G[W]$ , then

$$\text{mc}_k(G) \geq \text{mc}_k(H) + \frac{k-1}{k} (e(G) - e(H)). \quad (42)$$

We can also obtain a  $k$ -cut by choosing one vertex class and then taking a  $(k-1)$ -cut of the remainder of the graph. In particular, for any  $W \subset V = V(G)$ ,

$$\text{mc}_k(G) \geq e(W, V \setminus W) + \text{mc}_{k-1}(G \setminus W). \quad (43)$$

In addition to these observations, our proof of Theorem 5 will be based on the following two lemmas.

**Lemma 6.** *Suppose that  $G$  is a graph with  $m$  edges such that*

$$\text{mc}_k(G) = \frac{k-1}{k}m + \alpha \quad (44)$$

*and  $W \subset V$  satisfies the local inequality*

$$|\Gamma(v) \cap (V \setminus W)| \geq (k-1)|\Gamma(v) \cap W| \quad (45)$$

*for all  $v \in W$ . Then*

$$e(W) \leq \frac{m}{k^2} + \frac{k-1}{k}\alpha. \quad (46)$$

*Proof.* Let  $V = V(G)$ . Using (43) and (41), we see that

$$\begin{aligned} \text{mc}_k(G) &\geq e(W, V \setminus W) + \text{mc}_{k-1}(G \setminus W) \\ &\geq e(W, V \setminus W) + \frac{k-2}{k-1}(m - e(W) - e(W, V \setminus W)) \\ &= \frac{k-2}{k-1}m + \frac{1}{k-1}e(W, V \setminus W) - \frac{k-2}{k-1}e(W). \end{aligned}$$

Summing (45) over vertices in  $W$ , we see  $e(W, V \setminus W) \geq 2(k-1)e(W)$ . So

$$\begin{aligned} \text{mc}_k(G) &\geq \frac{k-2}{k-1}m + 2e(W) - \frac{k-2}{k-1}e(W) \\ &= \frac{k-2}{k-1}m + \frac{k}{k-1}e(W). \end{aligned}$$

The result now follows by a simple calculation.  $\square$

The proof of Theorem 5 involves moving certain vertices between the vertex classes of a partition. The fact that we can find suitable vertices is guaranteed by the following lemma.

**Lemma 7.** *Suppose that  $\alpha < m/2k$  and*

$$\text{mc}_k(G) = \frac{k-1}{k}m + \alpha. \quad (47)$$

*Suppose that  $W \subset V$  and the local inequality (45) holds for every  $v \in W$ . If*

$$e(W) \geq \frac{m}{k^2} - \frac{\alpha}{k} \quad (48)$$

then there is a vertex  $v \in W$  with

$$|\Gamma(v) \cap W| \leq 4\sqrt{m} \quad (49)$$

and

$$|\Gamma(v) \cap (V \setminus W)| \leq \left(k - 1 + 4k^3 \frac{\alpha}{m}\right) |\Gamma(v) \cap W| \quad (50)$$

*Proof.* Let

$$\begin{aligned} T_1 &= \{v \in W : |\Gamma(v) \cap W| > 4\sqrt{m}\} \\ T_2 &= \left\{v \in W : |\Gamma(v) \cap (V \setminus W)| > \left(k - 1 + 4k^3 \frac{\alpha}{m}\right) |\Gamma(v) \cap W|\right\} \end{aligned}$$

It is enough to show that  $W \setminus (T_1 \cup T_2)$  is nonempty.

By (45), we have  $e(W, V \setminus W) \geq 2(k-1)e(W)$  and as  $e(W) + e(W, V \setminus W) \leq m$ , we get  $e(W) \leq m/(2k-1)$ . Since  $\sum_{v \in T_1} |\Gamma(v) \cap W| \leq 2e(W) \leq 2m/(2k-1)$ , we have

$$|T_1| \leq \frac{2e(W)}{4\sqrt{m}} \leq \frac{\sqrt{m}}{2(2k-1)}$$

and so

$$e(T_1) \leq \binom{|T_1|}{2} \leq \frac{m}{8(2k-1)^2}.$$

It follows that

$$\sum_{v \in T_1} |\Gamma(v) \cap W| \leq e(W) + e(T_1) \leq e(W) + \frac{m}{8(2k-1)^2}. \quad (51)$$

We now concentrate on bounding  $\sum_{v \in T_2} |\Gamma(v) \cap W|$ . Calculating as in the proof of Lemma 6, we have

$$\text{mc}_k(G) \geq \frac{k-2}{k-1}m - \frac{k-2}{k-1}e(W) + \frac{1}{k-1}e(W, V \setminus W). \quad (52)$$

Now (45) and the definition of  $T_2$  imply that

$$\begin{aligned} e(W, V \setminus W) &= \sum_{v \in W} |\Gamma(v) \cap (V \setminus W)| \\ &\geq (k-1) \sum_{v \in W} |\Gamma(v) \cap W| + \frac{4k^3\alpha}{m} \sum_{v \in T_2} |\Gamma(v) \cap W| \\ &= 2(k-1)e(W) + \frac{4k^3\alpha}{m} \sum_{v \in T_2} |\Gamma(v) \cap W|. \end{aligned} \quad (53)$$



It therefore follows from (52) that

$$\text{mc}_k(G) \geq \frac{k-2}{k-1}m + \frac{k}{k-1}e(W) + \frac{4k^3\alpha}{m(k-1)} \sum_{v \in T_2} |\Gamma(v) \cap W|. \quad (54)$$

By (48) the right hand side is at least

$$\frac{k-1}{k}m - \frac{\alpha}{k-1} + \frac{4k^3\alpha}{m(k-1)} \sum_{v \in T_2} |\Gamma(v) \cap W|.$$

But then (47) implies that

$$\frac{4k^3\alpha}{m(k-1)} \sum_{v \in T_2} |\Gamma(v) \cap W| \leq \alpha + \frac{\alpha}{k-1} = \frac{k}{k-1}\alpha,$$

and so

$$\sum_{v \in T_2} |\Gamma(v) \cap W| \leq \frac{m}{4k^2}.$$

It therefore follows from (51) that

$$\sum_{v \in T_1 \cup T_2} |\Gamma(v) \cap W| \leq e(W) + \left( \frac{1}{4k^2} + \frac{1}{8(2k-1)^2} \right) m.$$

Since  $\alpha < m/2k$ , we have  $e(W) > m/2k^2$ . Since

$$\frac{1}{4k^2} + \frac{1}{8(2k-1)^2} < \frac{1}{2k^2},$$

we have

$$\sum_{v \in T_1 \cup T_2} |\Gamma(v) \cap W| \leq e(W) + \frac{m}{2k^2} < 2e(W),$$

and so  $T_1 \cup T_2 \neq W$ , as claimed.  $\square$

After this, we are ready to prove Theorem 5.

*Proof of Theorem 5.* We argue by induction on  $k$ . Let  $(V_1, \dots, V_k)$  be a maximum cut, and suppose that  $e(V_1) \geq \dots \geq e(V_k)$ . If  $e(V_1)$  satisfies (37) we are done. Otherwise,

$$e(V_1) = \frac{m}{k^2} - \frac{\alpha}{k} + \lambda, \quad (55)$$

where Lemma 6 implies that  $\lambda \leq \alpha$ .

We proceed by moving vertices one at a time from  $V_1$  to other vertex classes. Suppose we have reached a stage with vertex classes  $V'_1, \dots, V'_k$  (where  $V'_1 \subseteq V_1$ ). Applying Lemma 7, we find a vertex  $v$  satisfying (49) and (50), and move  $v$  to whichever class  $V'_i$ ,  $i > 1$ , contains fewest neighbours of  $v$ . This decreases  $e(V'_1)$  by  $|\Gamma(v) \cap V'_1| \leq 4\sqrt{m}$  and, by (50), decreases the size of the  $k$ -cut by at most

$$\begin{aligned} \min_{i>1} \{|\Gamma(v) \cap V'_i| - |\Gamma(v) \cap V'_1|\} &\leq \frac{1}{k-1} |\Gamma(v) \cap (V \setminus V'_1)| - |\Gamma(v) \cap V'_1| \\ &\leq \frac{4k^3\alpha}{m(k-1)} |\Gamma(v) \cap V'_1|. \end{aligned} \quad (56)$$

Since moving  $v$  does not affect the local inequality (45), we can continue to move vertices until  $V_1$  is reduced to  $W_1$  with

$$\frac{m}{k^2} - \frac{\alpha}{k} \leq e(W_1) \leq \frac{m}{k^2} - \frac{\alpha}{k} + 4\sqrt{m}. \quad (57)$$

Note that inequality (49) implies that we do eventually obtain  $W_1$  with  $e(W_1)$  in this range.

We end up with a set  $W_1 \subseteq V_1$  that satisfies (37), and sets  $W_2, \dots, W_k$  with  $W_i \supseteq V_i$  for each  $i$ . Since (55) and (57) imply that  $e(V_1) - e(W_1) \leq \lambda \leq \alpha$ , it follows from (56) that the size of the  $k$ -cut we end up with is at least

$$\frac{k-1}{k}m + \alpha - \alpha \cdot \frac{4k^3\alpha}{m(k-1)}. \quad (58)$$

Since  $(W_1, \dots, W_k)$  satisfies (57), by (58) we have

$$\begin{aligned} \sum_{i \geq 2} e(W_i) &\leq m - \left( \frac{k-1}{k}m + \alpha - \frac{4k^3\alpha^2}{m(k-1)} \right) - e(W_1) \\ &= \frac{m}{k} - \alpha + \frac{4k^3\alpha^2}{m(k-1)} - e(W_1) \\ &\leq \frac{k-1}{k^2}m - \frac{k-1}{k}\alpha + \frac{4k^3\alpha^2}{m(k-1)}. \end{aligned} \quad (59)$$

If  $k = 2$ , this implies (37) immediately. Otherwise, we consider the subgraph  $H = G[V \setminus W_1]$ , and partition it into  $k-1$  classes.

Suppose first that  $e(H) \leq \binom{k-1}{k}^2 m - \frac{(k-1)^2}{k} \alpha$ . We can find a judicious partition of  $H$  into  $k-1$  classes, each of which satisfies (3). Extending to a  $k$ -partition of  $G$  by taking  $W_1$  as the  $k$ th vertex class gives a partition satisfying (37).

Otherwise,  $e(H) > \binom{k-1}{k}^2 m - \frac{(k-1)^2}{k} \alpha$ . Note that since  $V_1$  satisfies the local inequality (39), so does  $W_1$ , and so  $e(W_1, V \setminus W_1) \geq 2(k-1)e(W_1)$ . Now

$$\text{mc}_{k-1}(H) \leq \frac{k-2}{k-1}e(H) + \frac{k}{k-1}\alpha, \quad (60)$$

or else, using (57), (43) and the local inequality,

$$\begin{aligned} \text{mc}_k(G) &\geq \text{mc}_{k-1}(H) + e(W_1, V \setminus W_1) \\ &> \frac{k-2}{k-1}(m - e(W_1) - e(W_1, V \setminus W_1)) + \frac{k}{k-1}\alpha + e(W_1, V \setminus W_1) \\ &= \frac{k-2}{k-1}m + \frac{1}{k-1}e(W_1, V \setminus W_1) - \frac{k-2}{k-1}e(W_1) + \frac{k}{k-1}\alpha \\ &\geq \frac{k-2}{k-1}m + 2e(W_1) - \frac{k-2}{k-1}e(W_1) + \frac{k}{k-1}\alpha \\ &= \frac{k-2}{k-1}m + \frac{k}{k-1}e(W_1) + \frac{k}{k-1}\alpha \\ &\geq \frac{k-2}{k-1}m + \frac{k}{k-1} \cdot \frac{m}{k^2} - \frac{k}{k-1} \frac{\alpha}{k} + \frac{k}{k-1}\alpha \\ &= \frac{k-1}{k}m + \alpha, \end{aligned}$$

which contradicts (36). Thus, writing

$$\text{mc}_{k-1}(H) = \frac{k-2}{k-1}e(H) + \gamma, \quad (61)$$

by (60) and our assumptions on the size of  $e(H)$  and  $\alpha$ ,

$$\begin{aligned} \gamma/e(H) &\leq \frac{k\alpha/(k-1)}{(k-1)^2m/k^2 - (k-1)^2\alpha/k} \\ &\leq \frac{(m/k^4) \cdot k/(k-1)}{m(k-1)^2/k^2 - (k-1)^2m/k^5} \\ &\leq 1/(k-1)^4. \end{aligned}$$

Applying the inductive hypothesis to  $H$ , we obtain a partition  $W'_2, \dots, W'_k$  with

$$\max_{i>1} e(W'_i) \leq \frac{e(H)}{(k-1)^2} - \frac{\gamma}{k-1} + (k-1)^5 \frac{\gamma^2}{e(H)} + 4\sqrt{e(H)} \quad (62)$$

Now, by (61),

$$\begin{aligned} \frac{e(H)}{(k-1)^2} - \frac{\gamma}{k-1} &= \frac{1}{k-1} (e(H) - \text{mc}_{k-1}(H)) \\ &\leq \frac{1}{k-1} \sum_{i \geq 2} e(W_i) \end{aligned} \quad (63)$$

and, since  $\gamma \leq \frac{k}{k-1}\alpha$  (by (60)) and  $e(H) \geq \left(\frac{k-1}{k}\right)^2 m - (k-1)^2 \alpha/k \geq \left(\frac{k-1}{k}\right)^2 m - (k-1)^2 m/k^5$ ,

$$\begin{aligned} \frac{\gamma^2}{e(H)} &\leq \left(\frac{k}{k-1}\right)^2 \frac{\alpha^2}{m} \frac{1}{(k-1)^2/k^2 - (k-1)^2/k^5} \\ &= \frac{\alpha^2}{m} \cdot \frac{k^7}{(k-1)^4(k^3-1)}. \end{aligned}$$

It follows from (62), (63) and (59) that

$$\begin{aligned} \max_{i \geq 1} e(W'_i) &\leq \frac{1}{k-1} \sum_{i \geq 2} e(W_i) + (k-1)^5 \frac{\gamma^2}{e(H)} + 4\sqrt{e(H)} \\ &\leq \frac{m}{k^2} - \frac{\alpha}{k} + \frac{4k^3}{(k-1)^2} \frac{\alpha^2}{m} + \frac{(k-1)^5 k^7}{(k-1)^4(k^3-1)} \frac{\alpha^2}{m} + 4\sqrt{m} \\ &\leq \frac{m}{k^2} - \frac{\alpha}{k} + k^5 \frac{\alpha^2}{m} + 4\sqrt{m}, \end{aligned}$$

for  $k \geq 3$ . The result now follows immediately by taking the partition  $W_1, W'_2, \dots, W'_k$ .  $\square$

Finally, we turn to the case when the maximum  $k$ -cut is very large. As in Alon, Bollobás, Krivelevich and Sudakov [1], we use a rather cruder argument.

**Theorem 8.** *Let  $k \geq 2$ . Suppose that  $G$  is a graph with  $m$  edges such that*

$$\text{mc}_k(G) = \frac{k-1}{k}m + \alpha,$$

where  $\alpha > m/k^6$ . Then, provided that  $m$  is sufficiently large (in terms of  $k$ ), there is a partition of  $V(G)$  into  $k$  sets, each of which contains at most

$$\frac{m}{k^2} - \frac{m}{12k^{10}} \quad (64)$$

edges.

*Proof.* Let  $(V_1, \dots, V_k)$  be a cut of size  $(k-1)m/k + \alpha$ . Let  $i \in \{1, \dots, k\}$  be chosen uniformly at random, and consider the partition  $(V_i, V \setminus V_i)$ . Then, writing  $m' = \sum_{j=1}^k e(V_j) = m/k - \alpha$  and  $p = 1 - q = 1/k$ , we have

$$\begin{aligned} \mathbb{E}(qe(V_i) + pe(V \setminus V_i)) &= q\frac{1}{k}m' + p\left(\frac{k-1}{k}m' + \frac{\binom{k-1}{2}}{\binom{k}{2}}(m - m')\right) \\ &= \frac{2k-2}{k^2}m' + \frac{k-2}{k^2}(m - m') \\ &= \frac{k-2}{k^2}m + \frac{1}{k}m' \\ &= \frac{k-1}{k^2}m - \frac{\alpha}{k} \\ &= pqm - \frac{\alpha}{k}. \end{aligned}$$

Suppose that  $m_{1/k}(G) = pqm - \alpha'$ . Since  $\alpha' > m/k^7$ , we can apply Theorem 3 with  $p = 1/k$  and  $c = 1/k^7$  to get a bipartition  $V(G) = V'_1 \cup V'_2$  with  $e(V'_1) \leq m/k^2 - m/12k^8$  and  $e(V'_2) \leq (k-1)^2m/k^2 - m/12k^8$ . We refine the partition by splitting  $V'_2$  into  $k-1$  pieces satisfying (3) (for the  $(k-1)$ -partite case). Providing  $m$  is sufficiently large (in terms of  $k$ ), we obtain a partition of  $V(G)$  satisfying (64).  $\square$

## 4 Conclusion

It seems likely that our constants could be improved significantly. It would be interesting to have sharper constants both when  $\delta$  is small (for instance, in (37)), and when  $\delta$  is large (for instance, in (64)). Particularly when  $\delta = \Omega(m)$ , all the bounds are rather crude, and it would be very interesting to know the correct dependence of the error term on  $\delta$ , and to have some idea of the extremal graphs.

It would be very interesting to prove analogous results for hypergraphs (see, for instance, [6] and [8] for results on judicious partitions of hypergraphs).

Finally, it would also be of interest to consider bisections instead of cuts. More specifically, for a graph  $G$ , let

$$b(G) = \max\{e(V_1, V_2) : V(G) = V_1 \sqcup V_2, ||V_1| - |V_2|| \leq 1\}$$

be the maximum size of a bisection of  $G$ , and let  $g_b(G)$  be the minimum of  $\max\{e(V_1), e(V_2)\}$  over bisections of  $G$ . What can be said about the relationship between  $b(G)$  and  $g_b(G)$ ? Note that the star  $K_{1,n-1}$  has  $b(K_{1,n-1}) = \lceil n/2 \rceil \sim e(K_{1,n-1})/2$ , while  $g_b(K_{1,n-1}) = \lfloor n/2 \rfloor - 1 \sim e(K_{1,n-1})/2$ , which is about as bad as it could be. But what about graphs with bisections much larger than  $m/2$ ?

**Acknowledgement.** We would like to thank the referees for their careful reading of the paper.

## References

- [1] N. Alon, B. Bollobás, M. Krivelevich and B. Sudakov, Maximum cuts and judicious partitions in graphs without short cycles, *J. Combin. Theory Ser. B* **88** (2003), 329–346
- [2] N. Alon, Bipartite subgraphs, *Combinatorica* **16** (1996), 301–311
- [3] N. Alon and E. Halperin, Bipartite subgraphs of integer weighted graphs, *Discrete Math.* **181** (1998), 19–29
- [4] B. Bollobás and A.D. Scott, Problems and results on judicious partitions, *Random Structures and Algorithms* **21** (2002), 414–430
- [5] B. Bollobás and A.D. Scott, Better bounds for Max Cut, *in Contemporary combinatorics*, Bolyai Soc. Math. Stud. **10** (2002), 185–246
- [6] B. Bollobás and A.D. Scott, Judicious partitions of 3-uniform hypergraphs, *European Journal of Combinatorics* **21** (2000), 289–300
- [7] B. Bollobás and A.D. Scott, Exact bounds for judicious partitions of graphs, *Combinatorica* **19** (1999), 473–486

- [8] B. Bollobás and A.D. Scott, Judicious partitions of hypergraphs, *J. Comb. Theory Ser. A* **78** (1997), 15–31
- [9] B. Bollobás and A.D. Scott, Judicious partitions of graphs, *Period. Math. Hungar.* **26** (1993), 125–137
- [10] C.S. Edwards, An improved lower bound for the number of edges in a largest bipartite subgraph, *in* Recent advances in graph theory (Proc. Second Czechoslovak Sympos., Prague, 1974), pp. 167–181. Academia, Prague, 1975
- [11] C.S. Edwards, Some extremal properties of bipartite subgraphs, *Canad. J. Math.* **25** (1973), 475–485
- [12] T.D. Porter, Minimal partitions of a graph, *Ars Combin.* **53** (1999), 181–186
- [13] T. D. Porter, Graph partitions, *J. Combin. Math. Combin. Comput.* **15** (1994), 111–118
- [14] T. D. Porter, On a bottleneck bipartition conjecture of Erdős, *Combinatorica* **12** (1992), 317–321
- [15] T.D. Porter and Bing Yang, Graph partitions II, *J. Combin. Math. Combin. Comput.* **37** (2001), 149–158
- [16] A.D. Scott, Judicious partitions and related problems, *in* Surveys in Combinatorics 2005, B.S. Webb *ed.*, London Mathematical Society Lecture Note Series **327**, Cambridge University Press 2005