

Trees and near-linear stable sets

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Abstract

When H is a forest, the Gyárfás-Sumner conjecture implies that every graph G with no induced subgraph isomorphic to H and with bounded clique number has a stable set of linear size. We cannot prove that, but we prove that every such graph G has a stable set of size $|G|^{1-o(1)}$. If H is not a forest, there need not be such a stable set.

Second, we prove that when H is a “multibroom”, there *is* a stable set of linear size. As a consequence, we deduce that all multibrooms satisfy a “fractional colouring” version of the Gyárfás-Sumner conjecture.

Finally, we discuss extensions of our results to the multicolour setting.

1 Introduction

Graphs in this paper are finite, and have no loops or parallel edges. For a graph G , $|G|$ denotes the number of vertices of G , $\chi(G)$ is its chromatic number, and $\omega(G)$ and $\alpha(G)$ denote the sizes of its largest clique and stable set respectively. If G, H are graphs, G is H -free if no induced subgraph of G is isomorphic to H , and if \mathcal{H} is a set of graphs, G is \mathcal{H} -free if G is H -free for each $H \in \mathcal{H}$.

We are interested in how large $\alpha(G)$ must be, when $\omega(G)$ is bounded. For each integer $k \geq 1$, every K_{k+1} -free graph G satisfies $\alpha(G) \geq \Omega(|G|^{1/k})$, and for all $\varepsilon > 0$, there are K_{k+1} -free graphs G such that $\alpha(G) \leq |G|^{2/k+\varepsilon}$, so at least we have a rough idea of the right answer. But what happens if we fix a graph H and ask the same question for $\{H, K_{k+1}\}$ -free graphs G ? If H is a forest this question is particularly interesting, and is the focus of this paper.

The Gyárfás-Sumner conjecture ([5, 15]; see also [6]) says:

1.1 *For every forest H and every integer $k \geq 1$, there exists $c \geq 1$ such that $\chi(G) \leq c$ for every $\{H, K_{k+1}\}$ -free graph G .*

In particular, if the Gyárfás-Sumner conjecture is true, then every $\{H, K_{k+1}\}$ -free graph G has a stable set of size at least $|G|/c$. This is dramatically different from what happens when only K_{k+1} is excluded, and for such graphs G we ought at least to be able to prove a lower bound on $\alpha(G)$ better than $|G|^{1/k}$. Until now that had not been done, but we will prove:

1.2 *For every forest H and every integer $k \geq 1$, every $\{H, K_{k+1}\}$ -free graph G satisfies $\alpha(G) \geq |G|^{1-o(1)}$, and hence has chromatic number at most $O(|G|^{o(1)})$.*

(The second claim follows from the first by recursive deleting maximum stable sets.)

An appealing feature of the Gyárfás-Sumner conjecture is that it is sharp, in the sense that if H is not a forest, and $k \geq 2$, then there does not exist c such that $\chi(G) \leq c$ for every $\{H, K_{k+1}\}$ -free graph G . This is a consequence of a result of Erdős [4], who showed that for all $g \geq 1$ there exists $\varepsilon > 0$ and arbitrarily large graphs G of girth at least g in which every stable set has size at most $|G|^{1-\varepsilon}$. Thus such graphs G will be H -free if H has a cycle of length less than g , and will have chromatic number more than c when $|G|^\varepsilon > c$. The same example shows that 1.2 is sharp in the same sense: if H is not a forest, and $k \geq 2$, then not every $\{H, K_{k+1}\}$ -free graph G satisfies $\alpha(G) \geq |G|^{1-o(1)}$.

To prove 1.2, we could assume that H is a tree, and hence has a radius. Here is a more exact version of 1.2 (logarithms in this paper are to base two):

1.3 *Let $k, r \geq 2$ be integers, let $q := (r-1)(k-1)$, and let T be a tree of radius at most r . Then there exists $b > 0$ such that every $\{T, K_{k+1}\}$ -free graph G satisfies*

$$\alpha(G) \geq |G|^{1-b(\log |G|)^{-\frac{1}{q}}}.$$

(This bound is better than $\Omega(|G|^{1-\varepsilon})$, for any fixed $\varepsilon > 0$, but not as good as $\Omega(|G|/\text{polylog}(|G|))$.)

The same proof yields a stronger result, that implies 1.3 by setting $d = |G|$:

1.4 *Let $k, r \geq 2$ be integers, let $q := (r-1)(k-1)$, and let T be a tree of radius at most r . Then there exists $b > 0$ such that for every $d \geq 2$, every $\{T, K_{k+1}\}$ -free graph G with maximum degree at most d satisfies*

$$\alpha(G) \geq 2^{-b(\log d)^{1-\frac{1}{q}}} |G|.$$

For some trees T , we can do better than this: we can prove that all $\{T, K_{k+1}\}$ -free graphs have stable sets of linear size (as they should if the Gyárfás-Sumner conjecture is true), and that leads to the second main result of the paper. Let us explain. A *broom* is a tree obtained from a path $v_0-v_1-\dots-v_\ell$ (where $\ell \geq 1$) by adding some number of new vertices, each adjacent to v_ℓ . We call v_0 the *root* of the broom, and ℓ is its *length*. If $\ell_1, \dots, \ell_n \geq 1$, an (ℓ_1, \dots, ℓ_n) -*multibroom* is a graph obtained from brooms of lengths ℓ_1, \dots, ℓ_n by identifying their roots, and a *multibroom* is a (ℓ_1, \dots, ℓ_n) -multibroom for some choice of ℓ_1, \dots, ℓ_n .

It is not known whether all multibrooms satisfy the Gyárfás-Sumner conjecture; indeed, the only multibrooms currently known to do so are:

- $(1, \dots, 1, 2, \dots, 2)$ -multibrooms (due to Scott and Seymour [12], unifying earlier theorems of Gyárfás [5], Kierstead and Penrice [7], and Kierstead and Zhu [8]);
- (ℓ_1, ℓ_2) -multibrooms for all ℓ_1, ℓ_2 , and more generally, multibrooms made by identifying the roots of several brooms, where all but one of the brooms is a path (by Chudnovsky, Scott and Seymour [1]; Scott [11] had previously shown the result when all brooms are paths).

Incidentally, this is almost the complete list of trees known to satisfy the Gyárfás-Sumner conjecture¹. It is known [11] that for every tree T it suffices to exclude a finite list of subdivisions of T .

The second main result of this paper says that all multibrooms satisfy a “fractional” version of 1.1. If $a, b \geq 1$ are integers, an (a, b) -*fractional colouring* of G is a family $(A_i : 1 \leq i \leq a)$ of stable sets of G such that $|\{i \in \{1, \dots, a\} : v \in A_i\}| \geq b$ for each vertex v . The *fractional chromatic number* $\chi^*(G)$ of G is the minimum of a/b over all pairs a, b such that G admits an (a, b) -fractional colouring. We will prove:

1.5 *For every multibroom T , and every integer $k \geq 1$, there exists $c > 1$ such that $\chi^*(G) \leq c$ for every $\{T, K_{k+1}\}$ -free graph G .*

We will deduce this via linear programming duality from a weighted version of the following, which strengthens 1.3 for multibrooms to give a linear bound:

1.6 *For every multibroom T , and every integer $k \geq 1$, there exists $c > 0$ such that $\alpha(G) \geq c|G|$ for every $\{T, K_{k+1}\}$ -free graph G .*

In the final section, we discuss extensions of our results to the multicolour setting.

2 Excluding a general forest

This section contains the proof of 1.4. We follow a strategy of iterative sparsification: starting with a T -free graph G , we look for an induced subgraph G' of G such that G' is significantly sparser than G , but not too much smaller; provided we can control the trade-off between size and density, we can then iterate the process until we end up with a large stable set.

To make this strategy work, we need to be able to carry out the sparsification step without sacrificing too much in terms of size. We construct the sparse subgraph G' in small steps: we

¹The only other trees known to do so are those obtained by subdividing a star and adding one leaf, attached arbitrarily [1], and those obtained from two disjoint paths by adding an edge between them [14]. See [13] for discussion, and for other recent work see [9] and papers in the series [10].

repeatedly find a set A of vertices that is sparse to most of the graph, and then take a union of many such sets. The key part of the argument is to find a suitable set A . We do this in 2.1 by embedding T into G one vertex at a time: as G is T -free, this process must get stuck at some point, and we will use this to find the set A .

Before starting the argument, we make a few comments:

- The set A that we obtain must satisfy a suitable sparsity condition: in fact, we get two sets A and B so that the vertices of A have relatively few neighbours outside $A \cup B$. We then repeat the argument on $G \setminus (A \cup B)$, and continuing in this way, we obtain a sequence of disjoint sets A_i and B_i such that the A_i are sparse to later sets in the sequence. Provided the A_i are not too small and the B_i are not too large, we obtain the vertices of G' by taking a union of (some of) the sets A_i .
- We may get stuck at different stages when attempting to embed T . The size (and sparsity) of the sets A_i and B_i that we obtain depends on where we get stuck. So, in fact, we end up creating sets of several different types: one of these types will occur often enough so that we can obtain G' by taking all sets A_i of this type.
- As we grow the tree, the sets A_i get smaller. By embedding the vertices in a careful order, we can ensure that our bounds depend primarily on the *radius* of the tree T , and only weakly on the number of vertices of T . For this reason, and for convenience, we embed our tree in a depth-first search order.

Let us recall the definition of a depth-first enumeration of a tree. Let T be a tree with vertex set $\{\sigma_1, \dots, \sigma_t\}$. The numbering $(\sigma_1, \dots, \sigma_t)$ is a *dfs-enumeration for T rooted at σ_1* if, for each $i \in \{1, \dots, t-1\}$, σ_{i+1} has a neighbour in the path of T between σ_1, σ_i . By growing a depth-first tree in T , it follows that every tree T admits a dfs-enumeration rooted at any vertex $\rho \in V(T)$.

Dfs-enumerations have useful properties that will be helpful in the proof of 2.1. If $(\sigma_1, \dots, \sigma_t)$ is a dfs-enumeration for T , then for $1 \leq i \leq t$ the subgraph T_i of T induced on $\{\sigma_1, \dots, \sigma_i\}$ is a tree. We refer to the path P_i from σ_1 to σ_i as the *active path* of T_i . Later vertices in the dfs-enumeration of T can only attach to vertices of T_i that lie on the active path: if $j > i$ and σ_j has a neighbour σ' in T_i then σ' must lie on the active path P_i . Furthermore, if $i \geq 2$ and σ_i is a leaf then no vertex later in the dfs-enumeration is adjacent to σ_i .

The following two facts will be helpful:

2.1 *Let G be a graph with clique number at most k , where $k \geq 1$ is an integer, and let $n_0, \dots, n_k \geq 1$ be integers with $n_0 \leq |G|$. Then there exist $p \in \{1, \dots, k\}$ and an induced subgraph H of G with at least n_{p-1} vertices and maximum degree less than n_p .*

Proof. We prove this by induction on k . For $k = 1$ the statement is immediate as then G has no edges and we can take $H = G$. So suppose $k > 1$ and we have proved the statement for smaller k . If G has maximum degree less than n_1 then we can take $H = G$. Otherwise, pick a vertex of maximum degree and let G' be the subgraph of G induced by its neighbours: G' has clique number at most $k-1$, and applying the inductive statement to G' with parameters n_1, \dots, n_k gives the required induced subgraph H . ■

2.2 For all $d, n > 0$ where n is an integer, and for every graph G with maximum degree less than d , there is an induced subgraph of G with at least $|G|/n$ vertices and with maximum degree less than d/n .

Proof. Let (V_1, \dots, V_n) be a partition of $V(G)$ such that $\sum_{i=1}^n |E(G[V_i])|$ is minimum. We can choose i with $|V_i| \geq |G|/n$. Suppose that $v \in V_i$. For $j \neq i$, v has at least as many neighbours in V_j as in V_i , or else moving v to V_j would contradict minimality. Thus v has at most $d_g(v)/n \leq d/n$ neighbours in V_i . This proves 2.2. ▀

A *copy* of H in G is an isomorphism ϕ from H to an induced subgraph $\phi(H)$ of G . For $A, B \subset V(G)$ and $y \in (0, 1)$, we say that A is *y-sparse to B* if every vertex of A has at most $y|B|$ neighbours in B . A key lemma in the proof of 1.4 shows that if G does not contain a copy of T then we can find sets A and B such that B is not too large compared to A , and vertices of A are sparse to $G \setminus B$. We will apply this later to prove 2.4, which will allow us to perform the sparsification step.

Here is the key lemma.

2.3 Let $k, r, t \geq 2$ be integers, let $q := (k-1)(r-1)$, and let $y_0, y_1, \dots, y_q \in (0, 1)$ be such that $y_p \leq y_{p-1}/3t$ for $1 \leq p \leq q$. Let T be a tree with t vertices and radius at most r , and let G be a T -free graph with clique number at most k and maximum degree d , such that $d \geq 6t/y_{q-1}$. Then there exist $p \in \{1, \dots, q\}$ and disjoint $A, B \subseteq V(G)$ such that

- $|A| \geq y_{p-1}d/2$ and $|B| \leq 2rtd$;
- $G[A]$ has maximum degree less than $y_p d$; and
- every vertex in A has fewer than $\max(y_p d, k^t)$ neighbours in $G \setminus (A \cup B)$.

Proof. The proof proceeds by attempting to grow a copy of T in dfs-enumeration order: since G does not contain a copy of T we will get stuck at some stage, and we will use this to obtain the sets A and B . At each stage of the argument, we identify a copy of some subtree T_s of T , and some sets A_i of vertices that will allow us to grow the embedding further. Since our vertices are chosen in dfs order, later vertices of the tree cannot have arbitrary neighbours in T_s : they can only be adjacent to vertices on the active path P_s . So for each vertex v_i of P_s we keep a set of A_i of vertices that we can use to add a neighbour of v_i to our embedding at a later stage (vertices at distance r from the root are handled slightly differently, as they are not adjacent to later vertices in the dfs-enumeration). We need the sets A_i to satisfy several restrictions: the vertices of A_i should be adjacent to v_i and no other vertex of T_s ; each pair A_i, A_j of sets must satisfy some density condition (because adding a vertex from A_j to our embedding means that its neighbours in A_i can no longer be added to the embedding and have to be removed from A_i , and so we will only use vertices with few neighbours in A_i); and it will be convenient if the sets A_i are appropriately sparse, as we need to produce a set A with low density. The parameters depend on the distance from the root in T : as we get deeper, the sets A_i get smaller, and we need to adjust the numbers accordingly.

Let us make all this more precise. Let $\sigma_1 \in V(T)$ be a vertex with distance at most r to every other vertex in T , and let $(\sigma_1, \dots, \sigma_t)$ be a dfs-enumeration of T rooted at σ_1 . For $1 \leq s \leq t$ let T_s be the subtree of T induced on $\{\sigma_1, \dots, \sigma_s\}$.

Let $s \in \{1, \dots, t\}$ and let $v_1 \cdots v_\ell$ be the active path in T_s (between $v_1 = \sigma_1$ and $v_\ell = \sigma_s$). Thus $\ell \leq r+1$. Let ϕ be a copy of T_s in G , with vertex set $U = \phi(V(T_s))$, and for $i = 1, \dots, \ell$, let

$w_i = \phi(v_i)$. We say that ϕ is *good* if there are sets $(A_i)_{1 \leq i \leq \min(\ell, r)}$ and integers $(p_i)_{1 \leq i \leq \min(\ell, r-1)}$ such that:

- the A_i are pairwise disjoint subsets of $V(G) \setminus U$;
- for $i = 1, \dots, \ell$, every vertex in A_i is adjacent to w_i and has no other neighbours in U ;
- for $i = 1, \dots, \min(\ell, r-1)$, we have $p_i \in \{1, \dots, (k-1)i\}$, and

$$|A_i| \geq (1 - \frac{s}{2t})y_{p_i-1}d,$$

and $G[A_i]$ has maximum degree less than $y_{p_i}d$;

- if $\min\{\ell, r\} = r$ then A_r is stable and $|A_r| \geq t - s$; and
- for $1 \leq h < i \leq \min(\ell, r)$, A_i is $\frac{1}{2t-s}$ -sparse to A_h .

We refer to the collection of A_i and p_i as *references* for ϕ . Note that the case $i = r$ (if it occurs) is special, as we will only use A_r to embed leaves of T (and we will not need to define a set A_{r+1}).

We can now proceed with the main argument. So suppose, for a contradiction, that there do not exist $p \in \{1, \dots, q\}$ and $A, B \subseteq V(G)$ satisfying the lemma.

There is no good copy of $T_i = T$ in G , since G is T -free. On the other hand, there is a good copy of T_1 . To see this, let v be a vertex of degree d in G , and let D be the set of its neighbours. Thus $|D| = d \geq y_1d$. Applying 2.1 to $G[D]$ with k replaced by $k-1$, and $n_p = y_p d$ for $0 \leq p \leq k-1$, there must exist $p \in \{1, \dots, k-1\}$ and $A \subseteq D$ such that $|A| \geq y_{p-1}d \geq (1 - \frac{1}{2t})y_{p-1}d$ and $G[A]$ has maximum degree less than $y_p d$. The map sending ρ to v is therefore a good copy of T_1 .

Consequently there is a maximum value of $s \in \{1, \dots, t\}$ such that there is a good copy of T_s in G , and s is less than t . Let ϕ be a good copy of T_s , let $v_1 - \dots - v_\ell$ be the active path P_s of T_s (where $v_1 = \sigma_1$ and $v_\ell = \sigma_s$), and let $(A_i)_{1 \leq i \leq \min(\ell, r)}$ and $(p_i)_{1 \leq i \leq \min(\ell, r-1)}$ be references for ϕ .

Since we have a dfs-enumeration, σ_{s+1} is adjacent in T to some vertex v_j of P_s , where $j \leq r$. We will show that either we can find a set B so that the pair (A_j, B) satisfies the lemma, or we can extend ϕ by taking a suitable $x \in A_i$ as the image of σ_{s+1} ; this will give a contradiction, as we have assumed that neither outcome occurs.

We start by defining x and (if necessary) B . There are two cases, depending on the value of j .

If $j < r$, we define $x \in A_j$ and B as follows. Let B be the set of all vertices $v \in V(G)$ satisfying either of the following conditions:

- v is adjacent to or equal to some vertex of U ; or
- v has more than $\frac{1}{2t-s}|A_i|$ neighbours in A_i for some $i \in \{1, \dots, j\}$.

There are at most $s(d+1)$ vertices that are adjacent to or equal to some vertex of U . For $i \in \{1, \dots, j\}$, the sum of vertex degrees of vertices in A_i is at most $d|A_i|$ and so there are at most $(2t-s)d$ vertices with at least $\frac{1}{2t-s}|A_i|$ neighbours in A_i . Consequently

$$|B| \leq s(d+1) + j(2t-s)d \leq 2rtd.$$

Since A_j is one of the references for ϕ , we have that

$$|A_j| \geq (1 - \frac{s}{2t})y_{p_j-1}d \geq y_{p_j-1}d/2$$

and $G[A_j]$ has maximum degree at most $y_{p_j}d$. Since, by assumption, the pair (A_i, B) does not satisfy the lemma, it follows that there is a vertex $x \in A_j$ with at least $\max(y_{p_j}d, t^k)$ neighbours in $V(G) \setminus (A_j \cup B)$. This defines x .

The case $j = r$ is simpler, as then σ_{s+1} is a leaf of T and so has no neighbours later in the dfs-enumeration of T . In this case, we choose $x \in A_j$ arbitrarily and do not need to define B .

In either case, let ϕ' be the extension of ϕ to $V(T_{s+1})$ defined by $\phi'(\sigma_{s+1}) = x$ (and $\phi' = \phi$ otherwise). Clearly ϕ' is a copy of T_{s+1} in G : we claim, for a contradiction, that it is good.

To show the goodness of ϕ' , all that remains is to find references for it. Let $v'_1 \cdots v'_{j+1}$ be the active path P_{s+1} of T between σ_1 and σ_{s+1} ; thus, $v'_i = v_i$ for $1 \leq i \leq j$, and $v'_{j+1} = \sigma_{s+1}$. Let

$$A'_i = A_i \setminus (N(x) \cup \{x\})$$

for $1 \leq i \leq j$ (we will define A'_{j+1} later if we need it).

For $i < j$, since x is $\frac{1}{2t-s}$ -sparse to A_i , it follows that

$$|A'_i| \geq \left(1 - \frac{1}{2t-s}\right) |A_i| \geq \left(1 - \frac{1}{2t-s}\right) \left(1 - \frac{s}{2t}\right) y_{p_{i-1}}d = \left(1 - \frac{s+1}{2t}\right) y_{p_{i-1}}d.$$

For $1 \leq h < i \leq j$, since A_i is $\frac{1}{2t-s}$ -sparse to A_h , and $|A'_h| \geq (1 - \frac{1}{2t-s})|A_h|$, it follows that A'_i is $\frac{1}{2t-s-1}$ -sparse to A'_h .

If $j = r$, then (by the conditions for references of ϕ) A_j is stable and $|A_j| \geq t - s$. So A'_j is stable and $|A'_j| \geq t - s - 1$. It follows that A'_1, \dots, A'_r and p_1, \dots, p_{r-1} are references for ϕ' , and therefore ϕ' is good, a contradiction.

Thus $j < r$, and so (again, by the conditions for references of ϕ) $|A_j| \geq (1 - \frac{s}{2t})y_{p_{j-1}}d$, and $G[A_j]$ has maximum degree less than $y_{p_j}d$. Hence, noting that $x \in A_j$, we have

$$|A'_j| \geq |A_j| - y_{p_j}d - 1 \geq \left(1 - \frac{s}{2t}\right) y_{p_{j-1}}d - y_{p_j}d - 1 \geq \left(1 - \frac{s+1}{2t}\right) y_{p_{j-1}}d,$$

since $y_{p_{j-1}}d/(3t) \geq y_{p_j}d$ and $y_{p_{j-1}}d/(6t) \geq y_{q-1}d/(6t) \geq 1$. To complete the references for ϕ' , all that remains is to define A'_{j+1} .

As $j < r$, we have defined the set B above. Let $C = N(x) \setminus (A_j \cup B)$, so $|C| \geq \max(y_{p_j}d, t^k)$ by our choice of x . Each vertex in C is adjacent to x , has no neighbour in U , and is $\frac{1}{2t-s}$ -sparse to A_i and hence $\frac{1}{2t-s-1}$ -sparse to A'_i , for $1 \leq i \leq j$.

If $j = r - 1$, then since $|C| \geq k^t$ and G has clique number at most k , C contains a stable set A'_{j+1} of size t ; then A'_1, \dots, A'_r and p_1, \dots, p_{r-1} are references for ϕ' , again a contradiction. Finally, if $j \leq r - 2$, we apply 2.1 to $G[C]$, with $n_p = y_{p_j+p}d$ for $0 \leq p \leq k$: we obtain $C' \subseteq C$ and $p \in \{1, \dots, k\}$ such that $|C'| \geq y_{p_j+p-1}d$ and $G[C']$ has maximum degree at most $y_{p_j+p}d$. Let $A'_{j+1} := C'$ and $p'_{j+1} := p_j + p$. Then A'_1, \dots, A'_{j+1} and $p_1, \dots, p_j, p'_{j+1}$ are references for ϕ' , a contradiction. This proves 2.3. ■

We now use 2.3 to show that we can carry out the sparsification step in our strategy. The idea is to apply 2.3 repeatedly to give a sequence of pairs A_i, B_i such that the sets $A_i \cup B_i$ are pairwise disjoint, and each pair satisfies the conclusion of the lemma. We have to be a little careful, as 2.3 has several different outcomes; however, one of these will occur frequently enough that we can take the union of the corresponding A_i to obtain the sparse induced subgraph that we are looking for.

2.4 Let $k, r, t \geq 2$ be integers, let $q := (k-1)(r-1)$, and let $y_0, y_1, \dots, y_q \in (0, 1)$ be such that $y_i \leq y_{i-1}/3t$ for $1 \leq i \leq q$. Let T be a tree with t vertices and radius at most r , and let G be a T -free graph with clique number at most k and maximum degree at most d . Then there exist $p \in \{1, \dots, q\}$ and an induced subgraph of G with at least $(20qrtk^t)^{-1}y_{p-1}|G|$ vertices and maximum degree less than $y_p d$.

Proof. Let $n \geq 0$ be maximal such that for $1 \leq j \leq n$ there are subsets $A_j, B_j \subseteq V(G)$ and an integer $p_j \in \{1, \dots, q\}$, such that $A_1, \dots, A_n, B_1, \dots, B_n$ are pairwise disjoint, and for $1 \leq j \leq n$:

- $|A_j| > (4rt)^{-1}y_{p_j-1}|B_j|$;
- $G[A_j]$ has maximum degree less than $y_{p_j}d$; and
- let $C_j := A_1 \cup \dots \cup A_j$ and $D_j := B_1 \cup \dots \cup B_j$; then every vertex in A_j has at most $\max(y_{p_j}d, k^t)$ neighbours in $V(G) \setminus (C_j \cup D_j)$.

We claim:

$$(1) \quad C_n \cup D_n = V(G).$$

Suppose not, and let $F := G \setminus (C_n \cup D_n)$. Let d' be the maximum degree of F , and suppose first that $d' \geq 6ty_{q-1}^{-1}$. Then by 2.3 applied to F , there exist $p \in \{1, \dots, q\}$ and disjoint $A, B \subseteq V(F)$ with $|A| \geq y_{p-1}d'/2$ and $|B| < 2rd't$ such that $G[A]$ has maximum degree less than $y_p d'$, and every vertex in A has fewer than $\max(y_p d', k^t) \leq \max(y_p d, k^t)$ neighbours in

$$V(F) \setminus (A \cup B) = V(G) \setminus (C_n \cup D_n \cup A \cup B).$$

In particular $|A| \geq y_{p-1}d'/2 \geq (4rt)^{-1}y_{p-1}|B|$, and $d' \leq d$, so taking $A_{n+1} := A$ and $B_{n+1} := B$ and $p_{n+1} := p$ would contradict the maximality of n .

This implies that $d' < 6ty_{q-1}^{-1}$. But F as maximum degree d' , it has a stable set A of size at least $(d'+1)^{-1}|F|$. Let $B := V(F) \setminus A$. Then $|B| \leq \frac{d'}{d'+1}|F| \leq d'|A|$, and so

$$|A| \geq (6t)^{-1}y_{q-1}|B| \geq (4rt)^{-1}y_{q-1}|B|,$$

and setting $A_{n+1} := A$ and $B_{n+1} := B$ and $p_{n+1} = q$ would contradict the maximality of n . This proves (1).

Thus we have $C_n \cup D_n = V(G)$. For every $p \in \{1, \dots, q\}$, let $J_p := \{j \in \{1, \dots, n\} : p_j = p\}$; then

$$V(G) = \bigcup_{p=1}^q \bigcup_{j \in J_p} A_j \cup B_j,$$

and so there exists $p \in \{1, \dots, q\}$ with

$$\bigcup_{j \in J_p} |A_j \cup B_j| \geq q^{-1}|G|.$$

Let $C := \bigcup_{j \in J_p} A_j$ and $D := \bigcup_{j \in J_p} B_j$; then $|C| \geq (4rt)^{-1} y_{p-1} |D|$ since $|A_j| \geq (4rt)^{-1} y_{p-1} |B_j|$ for all $j \in J_p$. Hence $|D| \leq 4rt y_{p-1}^{-1} |C|$, and so $|C \cup D| \leq (4rt y_{p-1}^{-1} + 1) |C| \leq 5rt y_{p-1}^{-1} |C|$. We deduce that

$$|C| \geq (5rt)^{-1} y_{p-1} |C \cup D| \geq (5qrt)^{-1} y_{p-1} |G|.$$

Now, for each $j \in J_i$, at most

$$y_p d |A_j| / 2 + \max(y_p d, k^t) |A_j| \leq \frac{3}{2} \max(y_p d, k^t) |A_j|$$

edges have both ends in $\bigcup_{j' \in J_i, j' \geq j} A_{j'}$ and at least one end in A_j ; and so $G[C]$ has at most $\frac{3}{2} \max(y_p d, k^t) |C|$ edges.

If $G[C]$ has at most $\frac{3}{2} k^t |C|$ edges, then by Turán's theorem, $G[C]$ has a stable set S with

$$|S| \geq (3k^t + 1)^{-1} |C| \geq (4k^t)^{-1} |C| \geq (20qrtk^t)^{-1} y_{p-1} |G|$$

and the theorem holds. Otherwise, $G[C]$ has at most $\frac{3}{2} y_p d |C|$ edges. In this case let $S' \subseteq C$ be the set of vertices of degree at most $6y_p d$ in $G[C]$. Then $|C \setminus S'| \leq \frac{1}{2} |C|$ and so $|S'| \geq \frac{1}{2} |C|$. Since $G[S']$ has maximum degree at most $6y_p d$, there exists $S \subseteq S'$ with $|S| \geq |S'|/6$ such that $G[S]$ has maximum degree at most $y_p d$, by 2.2. So

$$|S| \geq |S'|/6 \geq |C|/12 \geq (60qrt)^{-1} y_{p-1} |G| \geq (20qrtk^t)^{-1} y_{p-1} |G|,$$

and again the theorem holds. This proves 2.4. ■

Finally, by iterating the sparsification given by 2.4, we obtain 1.4, which we restate:

2.5 *Let $k, r \geq 2$ be integers, let $q := (r-1)(k-1)$, and let T be a tree of radius at most r . Then there exists $b > 0$ such that for every $d \geq 2$, every $\{T, K_{k+1}\}$ -free graph G with maximum degree at most d satisfies*

$$\alpha(G) \geq 2^{-b(\log d)^{1-\frac{1}{q}}} |G|.$$

Proof. Let $t := |T|$ and $c = 20qrtk^t$. We will prove that the theorem holds with $b := \log(4c^2)$.

Let $d \geq 2$, and let G be a $\{T, K_{k+1}\}$ -free graph with maximum degree at most d . Let $x := (\log d)^{1/q}$. Suppose first that $x \leq b/2$, and so $b(\log d)^{-1/q} = b/x \geq 2$. Since G has maximum degree at most d , it has a stable set of size at least

$$|G|/(d+1) \geq d^{-2} |G| \geq d^{-b(\log d)^{-\frac{1}{q}}} |G|$$

(since $b, d \geq 2$), and so the theorem holds. Hence we may assume that $x \geq b/2$.

Let $G_0 := G$ and let d_0 be the maximum degree of G_0 ; then $d_0 \leq d$. For $i = 0, \dots, q$, let $y_i := 2^{-x^i}$; then $y_{i-1}/y_i = y_{i-1}^{1-x} \geq 2^{x-1} \geq 2^{b/2-1} \geq 3t$ for all $i \in \{1, \dots, q\}$. Inductively, 2.4 implies that for each $j \geq 1$, there exists $p_j \in \{1, \dots, q\}$ and an induced subgraph G_j of G_{j-1} with $|G_j| \geq c^{-1} y_{p_j-1} |G_{j-1}|$ and maximum degree $d_j \leq y_{p_j} d_{j-1}$.

Choose $n \geq 0$ minimal such that $d_n < 1$. The minimality of n implies

$$2^{-x^q} = 1/d \leq d_{n-1}/d \leq y_{p_1} y_{p_2} \cdots y_{p_{n-1}} \leq y_1^{n-1} \leq 2^{-x(n-1)}.$$

Consequently $x(n-1) \leq x^q$, and so $n \leq x^{q-1} + 1 \leq 2x^{q-1} = 2(\log d)/x$. Let

$$y := y_{p_1-1} y_{p_2-1} \cdots y_{p_n-1};$$

then $y^x = y_{p_1} y_{p_2} \cdots y_{p_n}$ and

$$1 \leq (y_{p_1} y_{p_2} \cdots y_{p_n-1}) d = (y^x / y_{p_n}) d \leq y^x d^2;$$

and so $y \geq d^{-2/x}$. Thus

$$|G_n| \geq c^{-n} y |G| \geq c^{-2(\log d)/x} d^{-2/x} |G| = d^{-2(\log c)/x - 2/x} |G| = d^{-b/x} |G| = d^{-b(\log d)^{-1/q}} |G|.$$

Since G_n has maximum degree less than 1, this proves 2.5. ■

3 Multibrooms and linear-sized stable sets

Now we move on to our second main theorem, 1.5, which we will prove in this section. For two integers $\ell, m \geq 0$, let us say an (ℓ, m) -broom is a tree obtained from a path of length ℓ , with ends a, b say, by adding m new vertices each adjacent to b . We call a the *root* of the broom. If G is a graph and $S \subseteq V(G)$, we denote by $N_G(S)$ the set of vertices in $V(G) \setminus S$ with a neighbour in S . A *weighting* on a graph G is a function $w: V(G) \rightarrow \mathbb{R}^+$; and for $S \subseteq V(G)$ we define $w(S) := \sum_{v \in S} w(v)$. We write $w(H)$ for $w(V(H))$ when H is an induced subgraph of G , and define the *maximum weighted degree* $\Delta(G, w) := \max_{v \in V(G)} w(N_G(v))$,

By the linear programming duality theorem, 1.5 is equivalent to the following:

3.1 *Let T be a multibroom and $k \geq 1$ an integer. Then there exists $c > 0$ such that if G is $\{T, K_{k+1}\}$ -free, and w is a weighting on G , then there is a stable set S of G with $w(S) \geq cw(G)$.*

One way to try to find a stable set containing a linear fraction of the total weight, is via degeneracy. A graph is *d-degenerate* if every non-null subgraph has a vertex of degree at most d . If we can find an induced subgraph G' of G that is d -degenerate (where d is some constant, depending on T and k only) and with total weight at least a constant fraction of the total weight of G , then since G' is $(d+1)$ -colourable, it has a stable set containing at least a $1/(d+1)$ fraction of the total weight of G' , and so we win. We will try to assemble such a large d -degenerate subgraph iteratively. Let us say a subset $X \subseteq V(G)$ is *d-degenerate in G* if it can be ordered $X = \{x_1, \dots, x_n\}$ such that for $1 \leq i \leq n$, x_i has at most d neighbours in $\{x_{i+1}, \dots, x_n\} \cup (V(G) \setminus X)$. The advantage of this more restrictive definition is that if X is d -degenerate in G , and X' is d -degenerate in $G \setminus X$, then $X \cup X'$ is d -degenerate in G , as is easily seen. This will allow us to grow a large d -degenerate subgraph by piecing together smaller subsets recursively.

We start with a lemma that we can use to grow a single branch of a multibroom. We show that, for an arbitrary vertex v in a weighted graph G , either we can grow a broom of any specified length rooted at v , or we can find two disjoint sets X, Y of vertices such that $w(X)$ is at least a constant fraction of $w(Y)$ and X has bounded degeneracy in $G \setminus (Y \cup \{v\})$. We will argue later by induction on the clique number k , so we assume that 3.1 holds for induced subgraphs with smaller clique number.

3.2 Let $k, \ell \geq 1$ and $m \geq 0$ be integers, and let $d \geq 1$. Let G be a graph with clique number at most k , let w be a weighting of G , and let v be a vertex of G . Suppose that, for every induced subgraph G' of G with clique number less than k , there is a stable set S of G' with $w(S) \geq w(G')/d$. Then one of the following holds:

- there is an induced (ℓ, m) -broom in G with root v ; or
- there exist disjoint $X, Y \subseteq V(G) \setminus \{v\}$ with $w(X) \geq w(X \cup Y)/(d2^{2\ell})$ and $w(X \cup Y) \geq w(N_G(v))$ such that X is k^m -degenerate in $G \setminus (Y \cup \{v\})$.

Proof. We proceed by induction on ℓ . Let G, v, w be as in the theorem, and let $\Delta = \Delta(G, w)$ be the maximum weighted degree. We may delete all vertices $u \neq v$ with $w(u) = 0$: at the end of an argument, either we get the desired broom, or we obtain sets X and Y as in the lemma, and we can add all the deleted vertices to Y . Thus we may assume that $w(u) > 0$ for all $u \neq v$.

One approach to finding a broom would be to start from v and explore into the graph by distance from v : for $i \geq 0$, let U_i be the set of vertices at distance i from v . If any vertex u in U_ℓ has k^m neighbours outside $U := \bigcup_{i=0}^{\ell} U_i$, then we can construct a broom by taking a shortest path from v to u and finding a stable set of size m among the neighbours of u outside U . Otherwise, vertices in U have few neighbours in the rest of the graph. However this approach is too simple: the sets U_i might be very large, or might quickly exhaust the graph. We therefore explore the graph more carefully. Starting with $R_0 = \{v\}$, we grow a sequence of sets R_i , where each set R_i is a subset of the neighbours L_i of R_{i-1} in the unexplored part of the graph. We choose R_{i-1} carefully, so that L_i is not too heavy.

Let $L_0 = \{v\}$ and $J_0 = \emptyset$. For $i = 1, \dots, \ell$, we define J_{i-1}, R_{i-1}, L_i inductively as follows:

- Let R_{i-1} be a maximal subset of L_{i-1} such that $N_G(R_{i-1}) \setminus (J_{i-1} \cup \{v\})$ has weight at most 2Δ .
- $L_i = N_G(R_{i-1}) \setminus (J_{i-1} \cup \{v\})$.
- $J_i = L_1 \cup \dots \cup L_i$.

Thus $L_1 = N_G(v)$ and $w(L_1) = w(N_G(v)) \leq \Delta$. Note that for any vertex w of L_i , we can find an induced path $P_w = u_1 - \dots - u_i$ such that $u_1 = v$, $u_i = w$ and $u_j \in R_j$ for each j ; furthermore, for $h < i$, u_h has no neighbours outside J_i .

(1) We may assume that $w(L_{j+1}) < \min(w(L_j), \Delta)$ for some $j \in \{1, \dots, \ell - 1\}$.

Suppose first that $w(L_i) \geq \min(w(L_{i-1}), \Delta)$ for all $i \in \{2, \dots, \ell\}$. Since $w(N_G(v)) = w(L_1) \leq \Delta$, it follows that $w(L_\ell) \geq w(N_G(v))$.

We claim that $w(J_{\ell-1}) \leq 2(\ell - 1)w(L_\ell)$. If $w(L_\ell) \geq \Delta$, this is clear, since $w(L_i) \leq 2\Delta$ by definition. If $w(L_\ell) < \Delta$, then by applying the condition $w(L_i) \geq \min(w(L_{i-1}), \Delta)$ repeatedly, we see that $\Delta > w(L_\ell) \geq \dots \geq w(L_1)$. So $w(J_{\ell-1}) \leq (\ell - 1)w(L_\ell)$. In either case, $w(J_{\ell-1}) \leq 2(\ell - 1)w(L_\ell)$.

Let $R_{\ell-1} = \{v_1, \dots, v_n\}$, numbered in arbitrary order, and for $1 \leq i \leq n$ let C_i be the set of vertices in L_ℓ that are adjacent to v_i and nonadjacent to each of v_1, \dots, v_{i-1} . Then C_1, \dots, C_n are pairwise disjoint and have union L_ℓ . For each i , $G[C_i]$ has clique number less than k , as C_i is contained in the neighbourhood of v_i ; and so by assumption there is a stable set S_i of $G[C_i]$ with

$w(S_i) \geq w(C_i)/d$. Let $S = S_1 \cup \dots \cup S_n$; thus $w(S) \geq w(L_\ell)/d$, and we have already shown that $w(L_\ell) \geq w(N_G(v))$.

Let $Y = J_{\ell-1} \cup (L_\ell \setminus S)$ and, for $1 \leq i \leq n$, let

$$Q_i = S_{i+1} \cup \dots \cup S_n \cup (V(G) \setminus (J_{\ell-1} \cup \{v\})).$$

No vertex $u \in C_i$ has m pairwise nonadjacent neighbours in Q_i , or we could find an induced (ℓ, m) -broom by taking the path P_w and adding these neighbours. Thus u has fewer than k^m neighbours in Q_i . Consequently S is k^m -degenerate in $G \setminus (Y \cup \{v\})$. Since $w(S \cup Y) \leq 2\ell w(L_\ell) \leq 2d\ell w(S) \leq d2^{2\ell}w(S)$, the theorem holds in this case. This proves (1).

We may therefore assume that $w(L_{j+1}) < \min(w(L_j, \Delta))$ for some $j \in \{1, \dots, \ell - 1\}$; choose j minimal with this property. From the minimality of j , it follows that $w(J_{j-1}) \leq 2(j-1)w(L_j)$, and $w(L_j) \geq w(N_G(v))$, as in the proof of (1).

Suppose first that $j = 1$. Since $G[L_1]$ has clique number less than k , and $L_1 \neq \emptyset$, there is a nonempty stable subset $S \subseteq L_1$ with $dw(S) \geq w(L_1)$. Let $Y = (L_1 \setminus C) \cup L_2$. Thus

$$w(S \cup Y) \leq w(L_1) + w(L_2) \leq 2w(L_1) \leq 2dw(S) \leq 2^{2\ell}dw(S),$$

and S is 0-degenerate and hence k^m -degenerate in $G \setminus (Y \cup \{v\})$. Hence the theorem holds in this case.

So we may assume that $j \geq 2$. Choose disjoint $X, Y' \subseteq L_j$ with $X \cup Y'$ maximal such that X is k^m -degenerate in $G[L_j \setminus Y']$ and $w(X \cup Y') \leq 2^{2(\ell-j+1)}dw(X)$.

Suppose that $X \cup Y' \neq L_j$. Choose $q \in L_j \setminus (X \cup Y')$, and let $p \in R_{j-1}$ be adjacent to q . Let $G' = G[(L_j \setminus (X \cup Y')) \cup \{p\}]$. There is a path P of length $j-1$ between v, p with interior in $R_1 \cup \dots \cup R_{j-2}$, and so there is no induced $(\ell - j + 1, m)$ -broom in G' with root p , since its union with P would be an induced (ℓ, m) -broom in G with root v . From the inductive hypothesis, applied to G' and p , there are disjoint $X'', Y'' \subseteq V(G') \setminus \{p\}$ with $X'' \neq \emptyset$ (since $w(q) > 0$) such that $w(X'' \cup Y'') \leq 2^{2(\ell-j+1)}dw(X'')$ and X'' is k^m -degenerate in $G' \setminus (Y'' \cup \{p\})$. But then $X \cup X''$ is k^m -degenerate in $G[L_j \setminus (Y' \cup Y'')]$, and

$$w(X \cup X'' \cup Y' \cup Y'') \leq 2^{2(\ell-j+1)}dw(X \cup X''),$$

contrary to the maximality of $X \cup Y'$.

Thus $X \cup Y' = L_j$. Then X is k^m -degenerate in $G \setminus (Y \cup \{v\})$, where $Y = Y' \cup J_{j-1} \cup L_{j+1}$. But

$$\begin{aligned} w(X \cup Y') &= w(L_j) \\ w(J_{j-1}) &\leq 2(j-1)w(L_j) \\ w(L_{j+1}) &\leq w(L_j) \end{aligned}$$

and it follows that

$$w(X \cup Y) \leq 2jw(L_j) \leq (2j)2^{2(\ell-j+1)}dw(X) \leq 2^{2\ell}dw(X).$$

Since

$$w(X) \geq (2^{2(\ell-j+1)}d)^{-1}w(X \cup Y') \geq (2^{2(\ell-j+1)}d)^{-1}w(N_G(v)) \geq 2^{-2\ell}d^{-1}w(N_G(v)),$$

the theorem holds in this case. This proves 3.2. ■

Now we will apply 3.2 to prove 3.1, which we restate:

3.3 *Let T be a multibroom and $k \geq 1$ an integer. Then there exists $c > 0$ such that if G is $\{T, K_{k+1}\}$ -free, and w is a weighting on G , then there is a stable set S of G with $w(S) \geq cw(G)$.*

Proof. We may assume that $k \geq 2$, and assume inductively that the theorem holds for smaller values of k . Thus there exists $d \geq 1$ such that for every $\{T, K_k\}$ -free graph G' , if w is a weighting on G' , then there is a stable set S of G' with $dw(S) \geq w(G')$.

Let B_1, B_2, \dots, B_b be vertex-disjoint brooms such that T is obtained by identifying their roots. For each $i \in \{1, \dots, b\}$, let B_i be an (ℓ_i, m_i) -broom, and let $\ell = \max(\ell_1, \dots, \ell_b)$, and $m = \max(m_1, \dots, m_b)$. We claim that $c = (d^2 2^{2\ell+3} |T| (k^m + 1))^{-1}$ satisfies the lemma. To see this, let G be $\{T, K_{k+1}\}$ -free, and let w be a weighting on G . Choose disjoint $X, Y \subseteq V(G)$ with $X \cup Y$ maximal such that:

- $d^2 2^{2\ell+3} |T| w(X) \geq w(X \cup Y)$; and
- X is k^m -degenerate in $G \setminus Y$.

Suppose first that $X \cup Y = V(G)$. Then $d^2 2^{2\ell+3} |T| w(X) \geq w(G)$, and since $G[X]$ is k^m -degenerate and hence $(k^m + 1)$ -colourable, it has a stable set S with $w(S) \geq w(X)/(k^m + 1) \geq cw(G)$, and the theorem holds. Thus, we may assume (for a contradiction) that $X \cup Y \neq V(G)$. Let $G' = G \setminus (X \cup Y)$, let $\Delta = \Delta(G', w')$, where w' is the restriction of w to $V(G')$, and choose $v \in V(G')$ with $w(N_{G'}(v)) = \Delta$. Since $G[N_{G'}(v)]$ has clique number less than k , there is a stable subset $S \subseteq N_{G'}(v)$ with $dw(S) \geq w(N_{G'}(v)) = \Delta$. Let Z be the set of vertices z of G' not in $S \cup \{v\}$, such that $w(N_G(z) \cap S) \geq w(S)/(2|T|)$. Let P be the sum of $w(u)w(z)$ over all adjacent u, z with $u \in S$ and $z \in Z$; then

$$w(Z)w(S)/(2|T|) = \sum_{z \in Z} w(z)w(S)/(2|T|) \leq P \leq \sum_{u \in S} w(u)\Delta = w(S)\Delta,$$

and so $w(Z) \leq 2|T|\Delta$. Let $Z' = N_{G'}(v) \setminus S$.

(1) *Let H be a connected induced subgraph of $G' \setminus (Z \cup Z')$, such that $v \in V(H)$ and $|H| \leq |T|$. Then for all ℓ', m' with $1 \leq \ell' \leq \ell$ and $0 \leq m' \leq m$, there is an (ℓ', m') -broom B in $G' \setminus (Z \cup Z')$, with root v , with $V(B \cap H) = \{v\}$, and such that $B \cup H$ is an induced subgraph of G' .*

Let W be the set of vertices of $G' \setminus (Z \cup Z' \cup \{v\})$ that are equal or adjacent to some vertex of $H \setminus \{v\}$. Thus $w(W \cup \{v\}) \leq |T|\Delta$. Since $w(N_{G'}(u)) \leq w(S)/(2|T|)$ for each vertex u of $H \setminus v$, it follows that $w(W \cap S) \leq w(S)/2$. Suppose that the desired broom B does not exist. Thus there is no (ℓ', m') -broom with root v in G'' , where $G'' = G' \setminus (Z \cup Z' \cup W)$. By 3.2 applied to G'' , and since $\ell' \leq \ell$ and $m' \leq m$, there exist disjoint $X', Y'' \subseteq V(G'') \setminus \{v\}$ with $d^{2\ell'} w(X') \geq w(X' \cup Y'') \geq w(N_{G''}(v))$ such that X' is $k^{m'}$ -degenerate in $G'' \setminus (Y'' \cup \{v\}) = G' \setminus Y'$, where

$$Y' = Y'' \cup Z \cup Z' \cup W \cup \{v\}.$$

Consequently $X \cup X'$ is k^m -degenerate in $G \setminus (Y \cup Y')$. But

$$w(N_{G''}(v)) \geq w(S) - w(W \cap S) \geq w(S)/2 \geq \Delta/(2d),$$

and so $d2^{2\ell}w(X') \geq \Delta/(2d)$. Furthermore,

$$\begin{aligned} w(Z) &\leq 2|T|\Delta \\ w(Z') &\leq \Delta \\ w(W \cup \{v\}) &\leq |T|\Delta \end{aligned}$$

and so

$$w(Z \cup Z' \cup W \cup \{v\}) \leq (3|T| + 1)\Delta \leq (3|T| + 1)d^22^{2\ell+1}w(X').$$

Consequently

$$w(X' \cup Y') \leq d2^{2\ell}w(X') + (3|T| + 1)d^22^{2\ell+1}w(X') \leq d2^{2\ell}w(X')(1 + 2(3|T| + 1)d) \leq d^22^{2\ell+3}|T|w(X').$$

But $X \cup X'$ is k^m -degenerate in $G \setminus (Y \cup Y')$, contradicting the maximality of $X \cup Y$. This proves (1).

By b applications of (1), it follows that G' contains a copy of T , a contradiction. This proves 3.3. ■

4 A multicolour version

One can look for “multicolour” versions of these theorems. One simple multicolour extension of the Gyárfás-Sumner conjecture is:

4.1 Conjecture: *For every forest H , and all integers $k, t \geq 1$, there exists $c > 0$ such that if G is a graph with clique number at most k , and the edges of G can be coloured with t colours in such a way that no monochromatic induced subgraph of G is isomorphic to H , then $\chi(G) \leq c$.*

(*Monochromatic* means with all edges the same colour.) Calling this a “conjecture” is a stretch, because we have very little faith in it. So far we cannot even prove it when H is the four-vertex path (although we can if in addition $t \leq 3$).

On the other hand, surprisingly to us, 4.1 is “nearly” true. Our first main theorem 1.2 can be extended to the following:

4.2 *For every forest H , and all integers $k, t \geq 1$, if G is a graph with clique number at most k , and the edges of G can be coloured with t colours in such a way that no monochromatic induced subgraph of G is isomorphic to H , then $\alpha(G) \geq |G|^{1-o(1)}$.*

It is straightforward to adapt the proof of 1.2 to give this more general result, and we omit the details. Pleasingly, this is tight, in the sense that if \mathcal{H} is a finite set of t -edge-coloured graphs (that is, graphs H with $E(H)$ partitioned into t subsets labelled $1, \dots, t$), and we want to know whether every t -edge-coloured graph with bounded clique number that contains no member of \mathcal{H} as an induced subgraph (in the natural sense) has bounded chromatic number, then 4.2 tells us the answer: if and only if for $1 \leq i \leq t$, there is a forest in \mathcal{H} with all edges coloured i .

Another multicolour extension of 1.1 generalizes a theorem of Chudnovsky and Seymour [2], who proved the following when $t = 2$:

4.3 Let $t \geq 1$ be an integer, and let \mathcal{H} be a finite set of t -edge-coloured complete graphs. Assuming the truth of 1.1, the following are equivalent:

- there exists $c > 0$ such that for every t -edge-coloured complete graph G with no induced subgraph in \mathcal{H} , $V(G)$ is the union of c monochromatic cliques
- for every two colours $i, j \in \{1, \dots, t\}$, \mathcal{H} contains a forest with all edges in colour i and all nonedges in colour j , and \mathcal{H} contains a complete multipartite graph with all edges in colour i and all nonedges in colour j .

This can be deduced from a result of [3] (we omit the details), although deriving the second bullet from the first assumes the Gyárfás-Sumner conjecture. But without assuming that, we can prove something similar:

4.4 Let $t \geq 1$ be an integer, and let \mathcal{H} be a finite set of t -edge-coloured complete graphs. The following are equivalent:

- every t -edge-coloured complete graph with no induced subgraph in \mathcal{H} has a monochromatic clique of size $|G|^{1-o(1)}$;
- for every two colours $i, j \in \{1, \dots, t\}$, \mathcal{H} contains a forest with all edges in colour i and all nonedges in colour j , and \mathcal{H} contains a complete multipartite graph with all edges in colour i and all nonedges in colour j .

This can easily be deduced from 1.2 and the theorem of [3], and we omit the details.

References

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