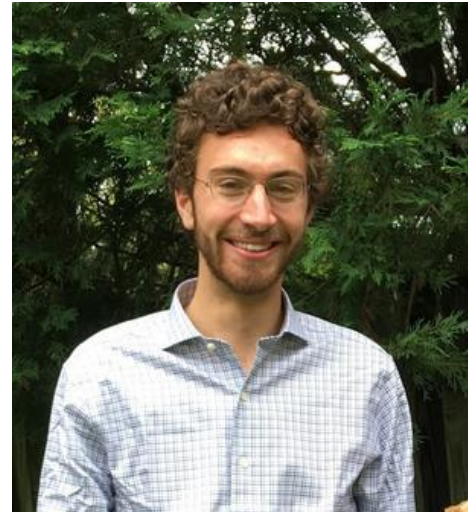


# Random Friends Walking on Random Graphs

**Noga Alon, Princeton and TAU**

**Joint work with Colin Defant and Noah Kravitz**

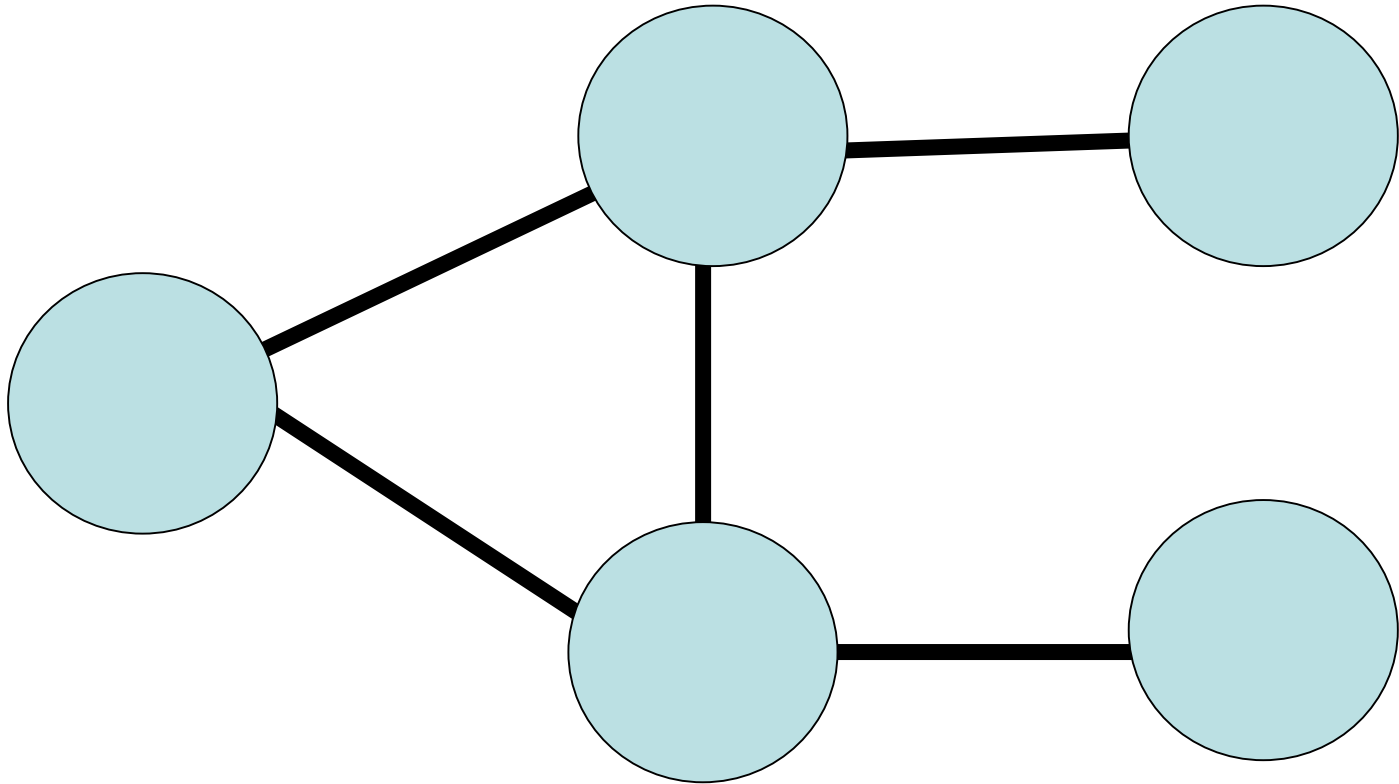


**Oxford, Jan. 2021**

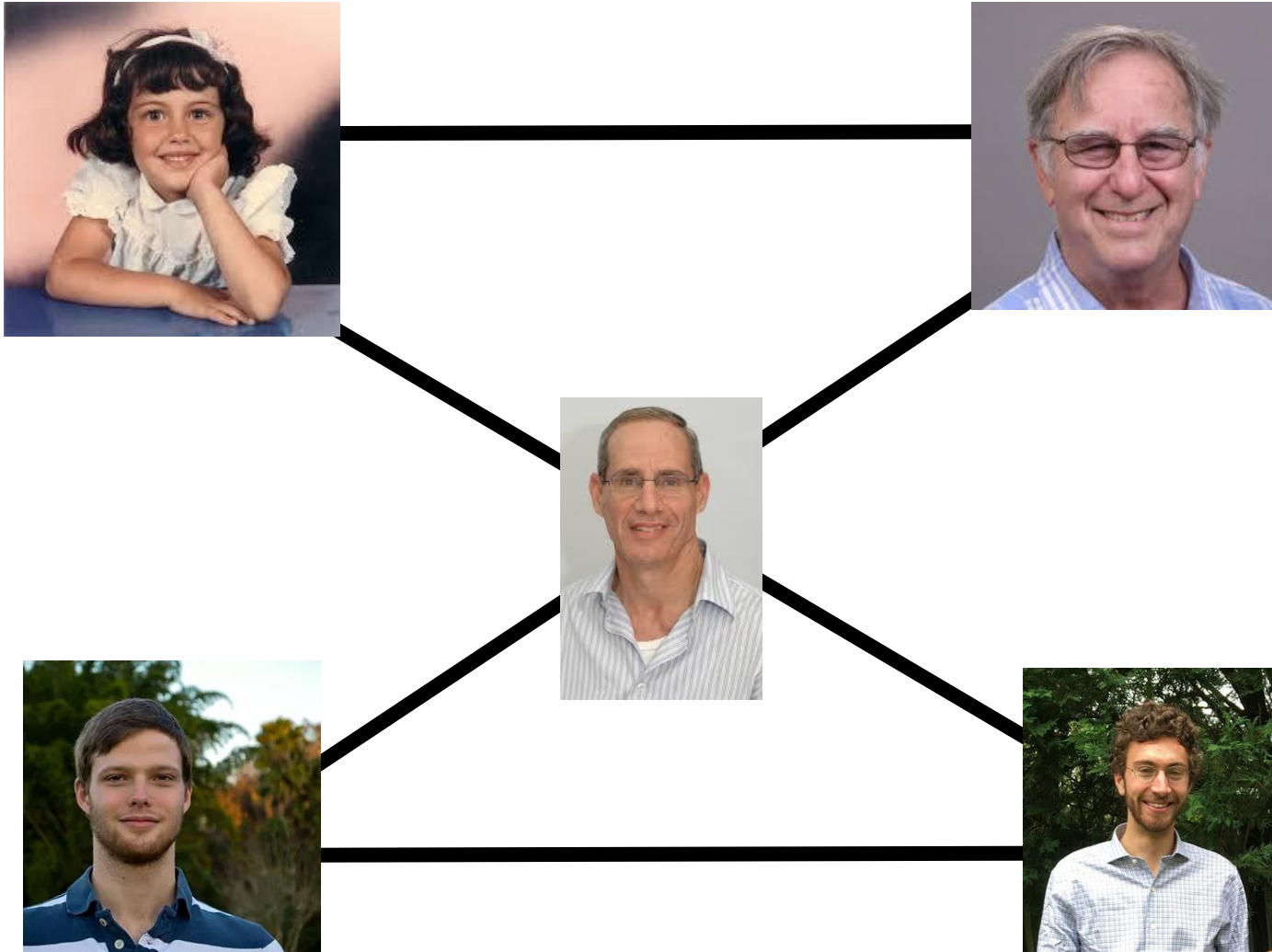
# Friends and Strangers Graphs

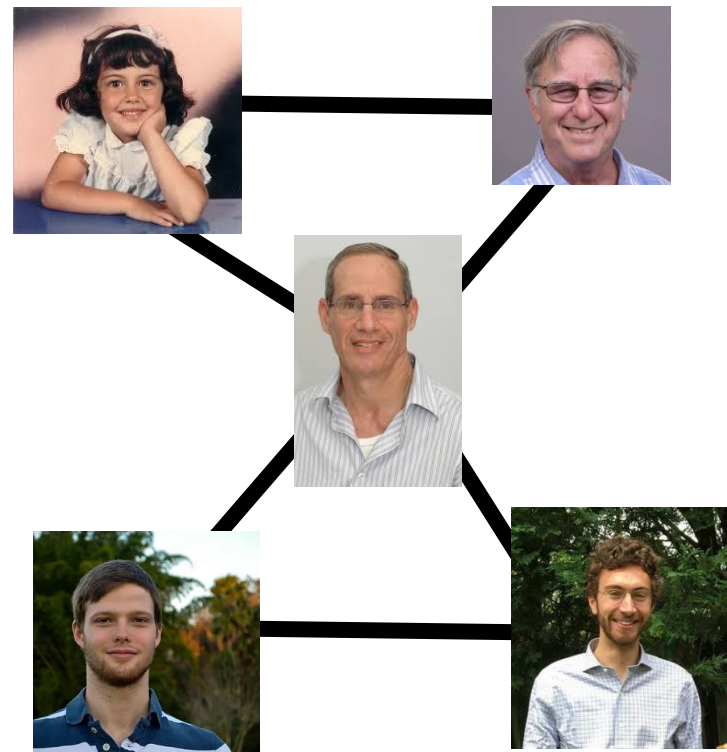
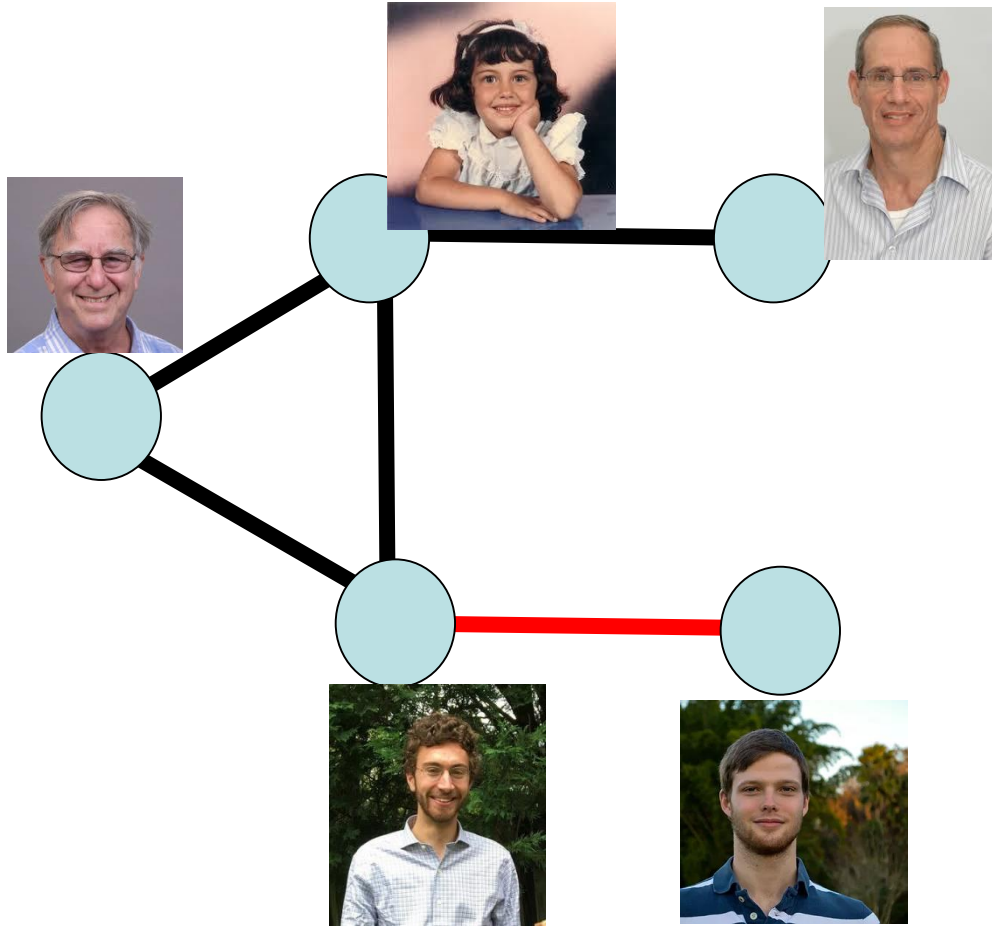


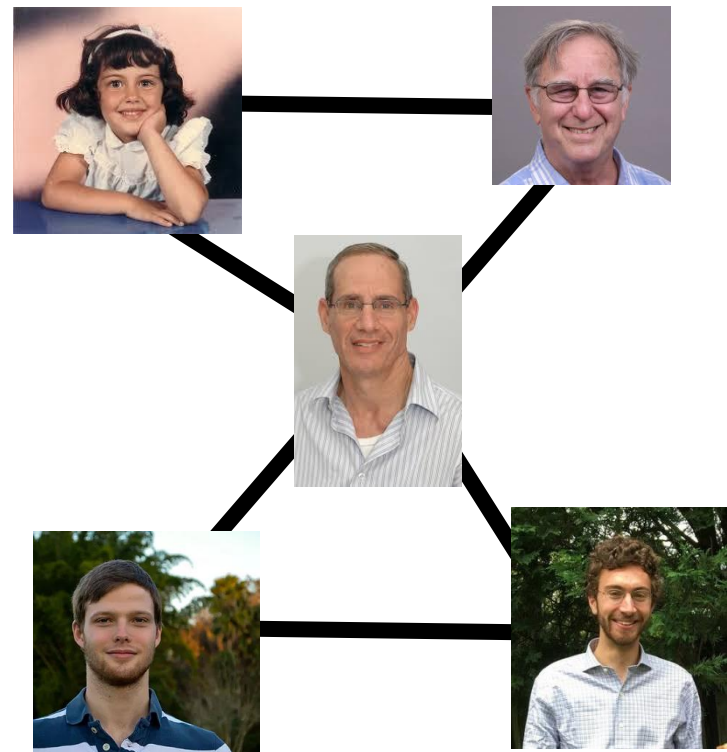
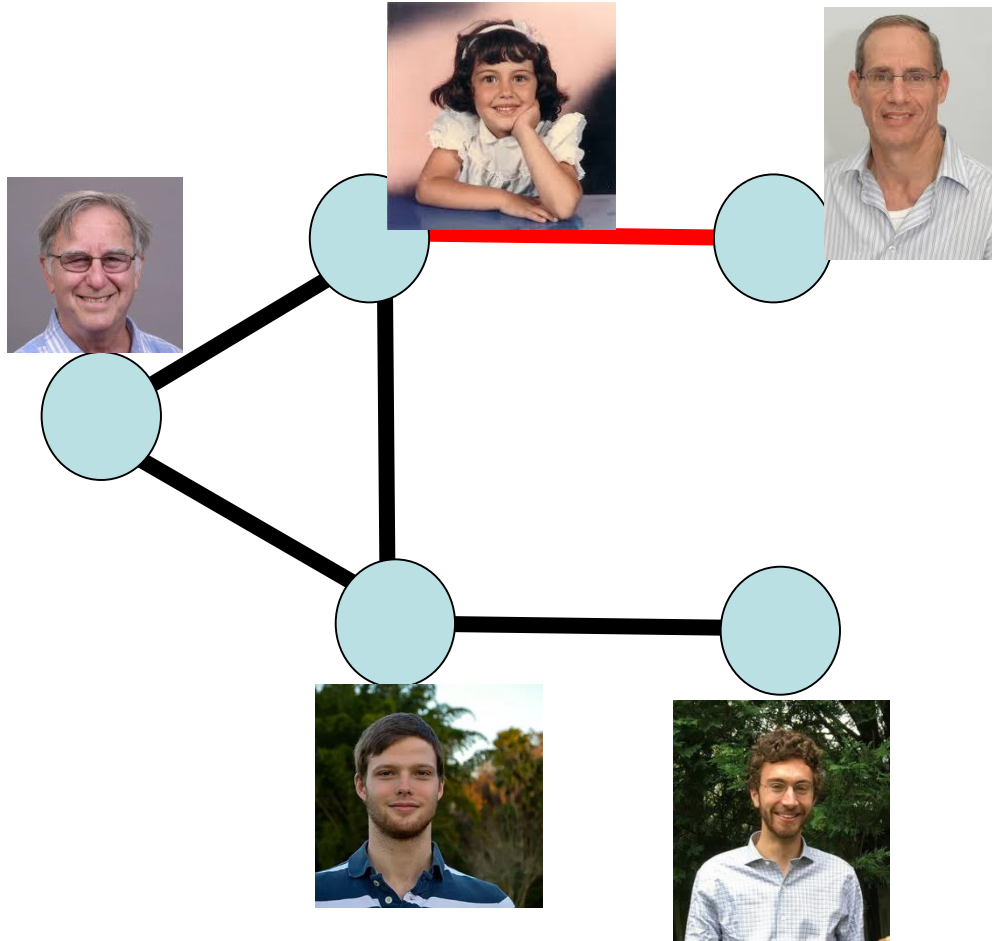
# Example

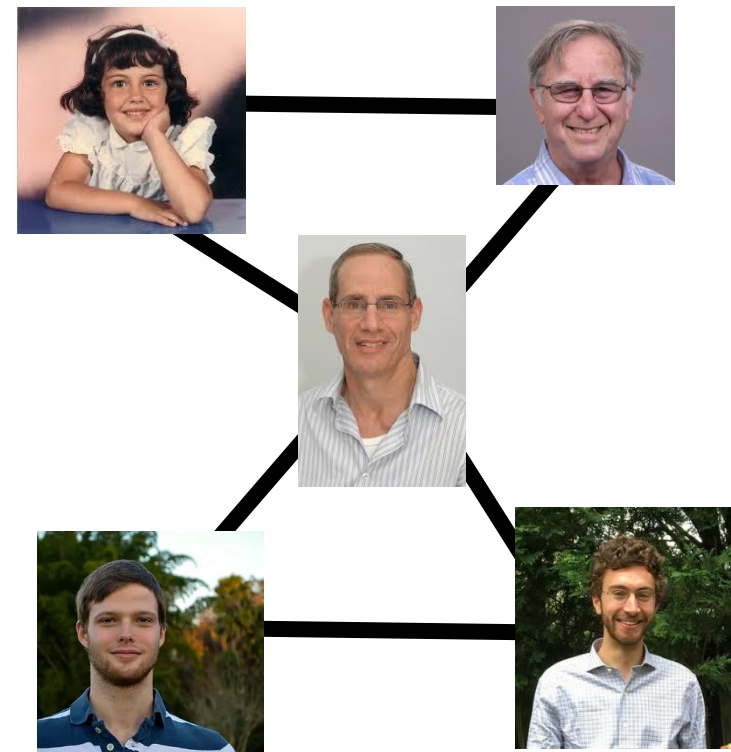
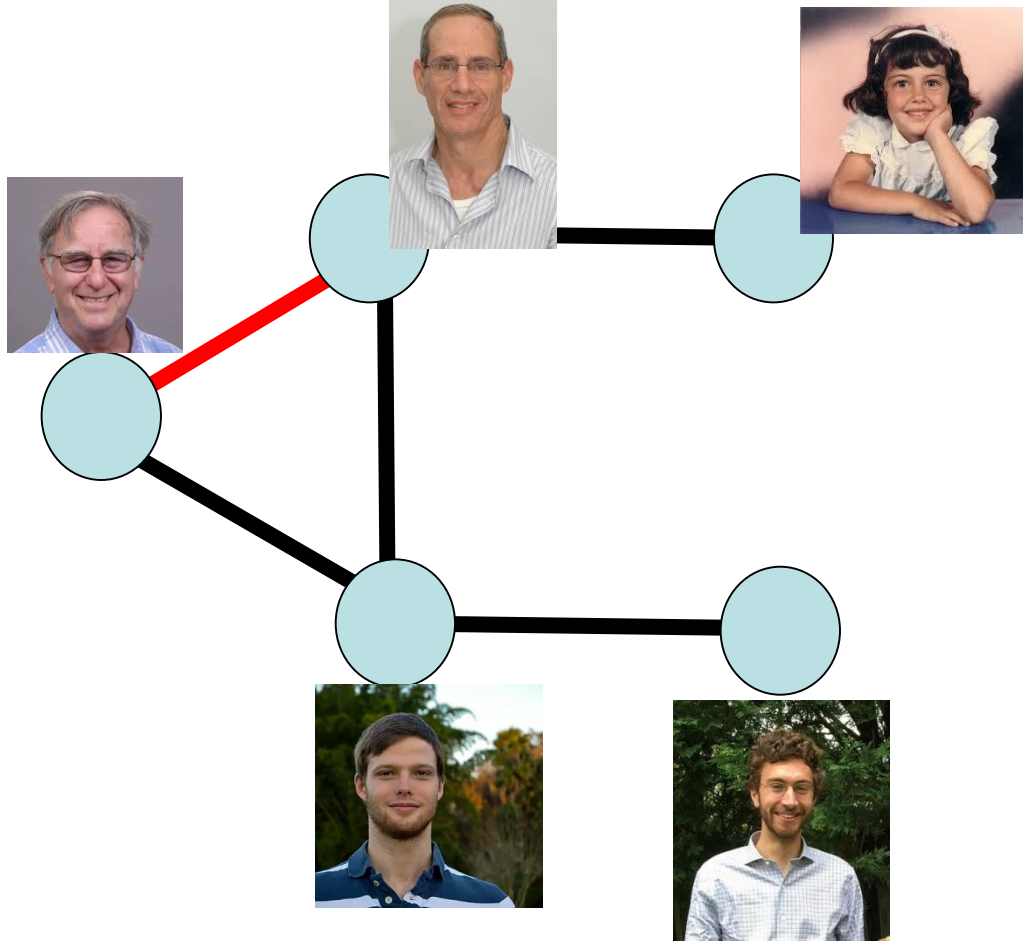


# Friendship Graph, example

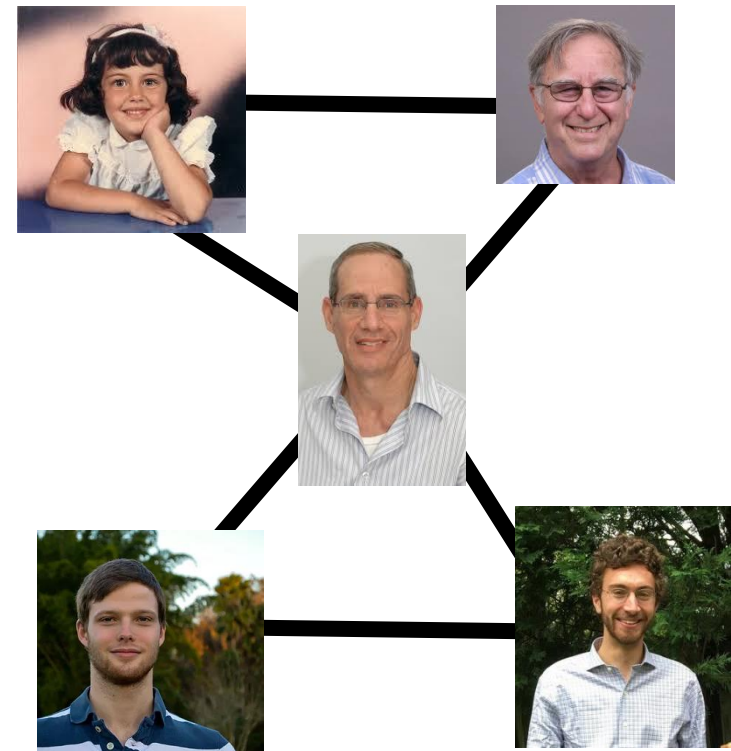
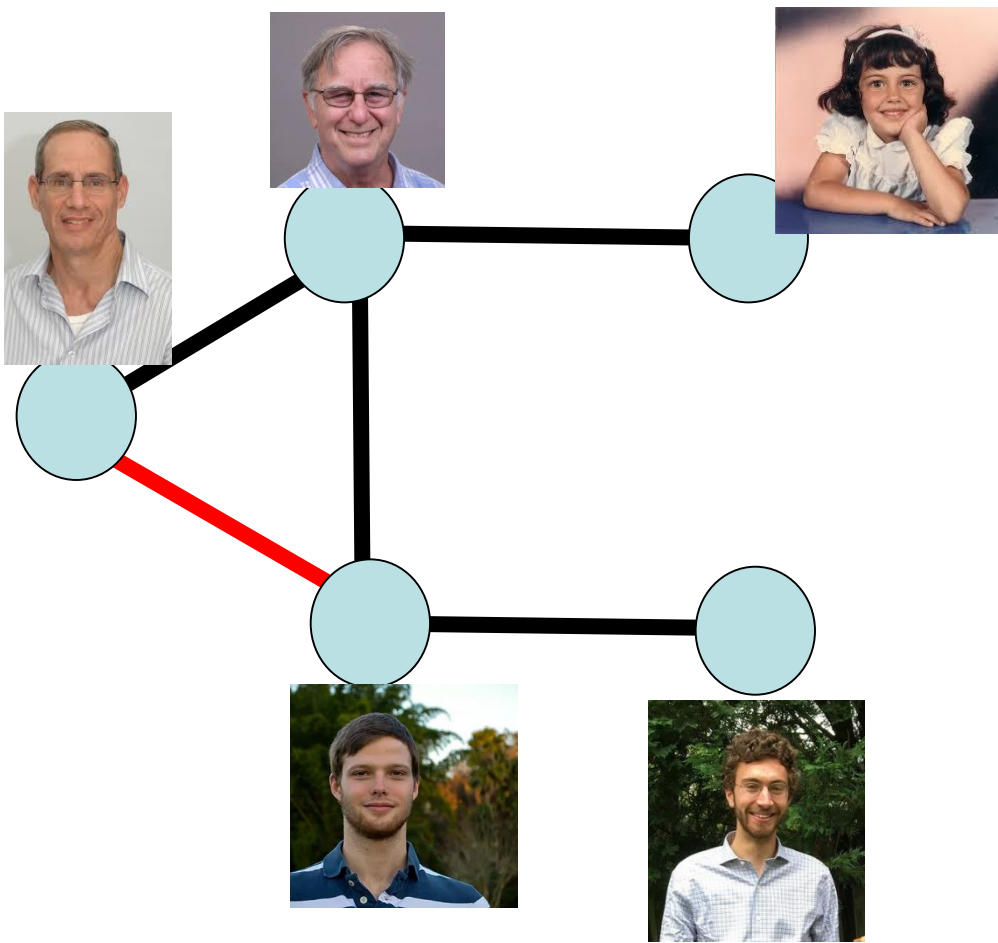




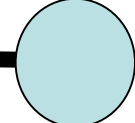
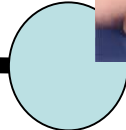
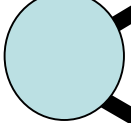
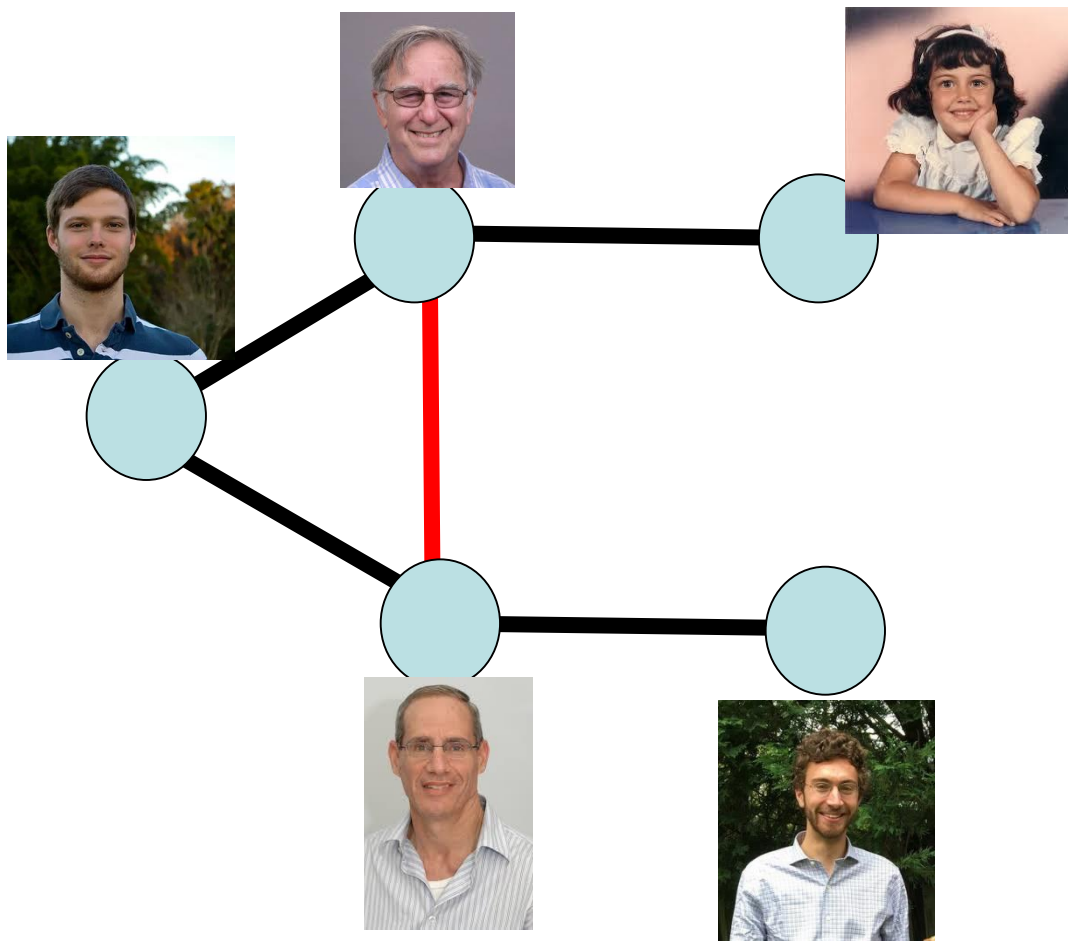


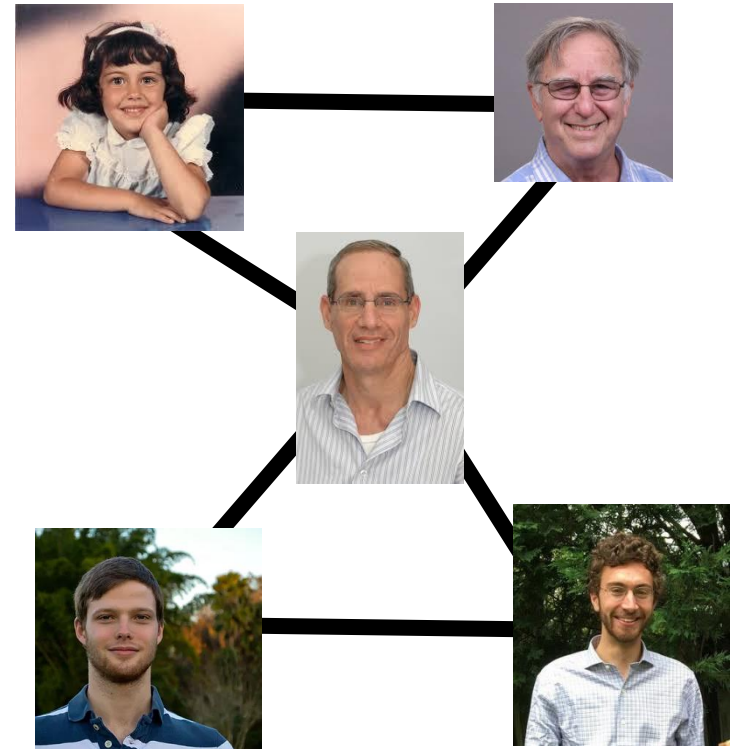
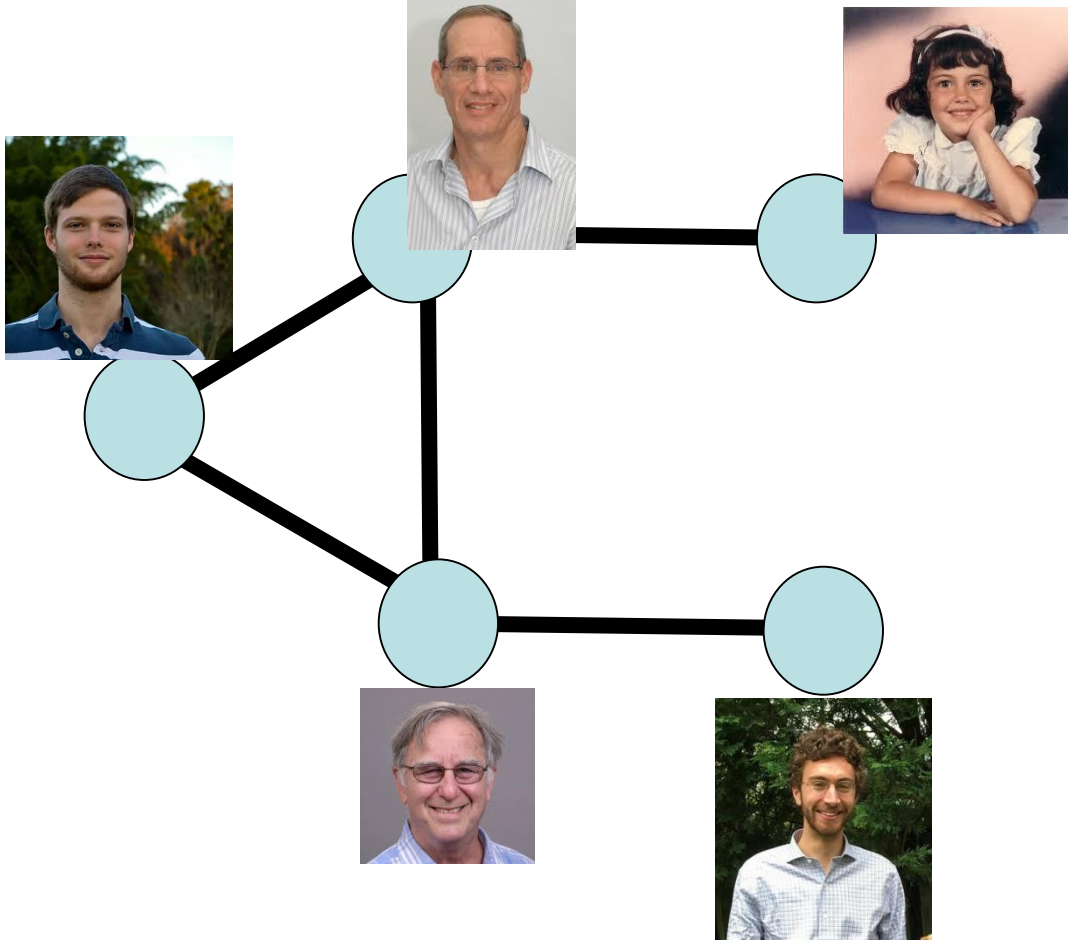






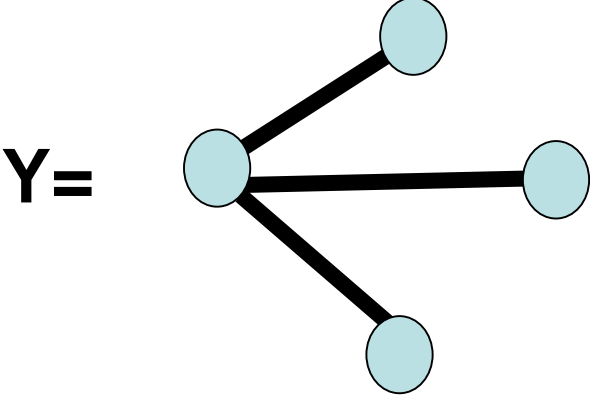
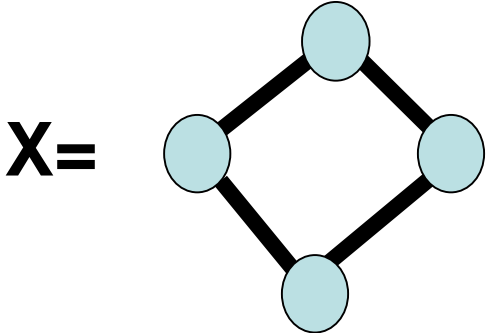






**Definition:** Let  $X, Y$  be two  $n$ -vertex graphs. The **Friends and Strangers Graph  $FS(X, Y)$**  is the graph whose vertices are the bijections from  $V(X)$  to  $V(Y)$  where two bijections are adjacent if one can be obtained from the other by a **friendly swap**.

**Example:**



$FS(X, Y) = 2C_{12}$

# Previous Work

For any graph  $Y$ ,  $FS(K_n, Y)$  is the **Cayley graph** of  $S_n$  generated by the transpositions corresponding to the edges of  $Y$

Analyzing the **15-puzzle game** is equivalent to analyzing  $FS(4 \text{ by } 4 \text{ grid}, K_{1,15})$



**Wilson (74)** studied the connected components of  $FS(X, K_{1,n-1})$  for arbitrary  $X$

**Stanley (12)** studied the connected components of  $FS(P_n, P_n)$ .

**Defant and Kravitz (20)** studied the connected components of  $FS(X, P_n)$ ,  $FS(X, C_n)$  for general  $X$

# Basic properties

Let  $X$  and  $Y$  be two  $n$ -vertex graphs

**Isolated vertices** of  $FS(X,Y)$  correspond to edge-disjoint packings of  $X,Y$  in  $K_n$

**$FS(X,Y)$**  is isomorphic to  **$FS(Y,X)$**

If  $X$  is **disconnected** so is  $FS(X,Y)$

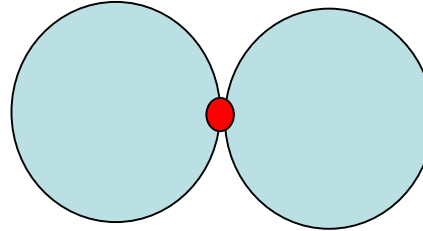
If  $X,Y$  are **bipartite** and  $n \geq 3$  then  $FS(X,Y)$  is disconnected

# The star graph $K_{1,n-1}$

When is  $FS(X, K_{1,n-1})$  **disconnected** ?

$X$  is **disconnected**

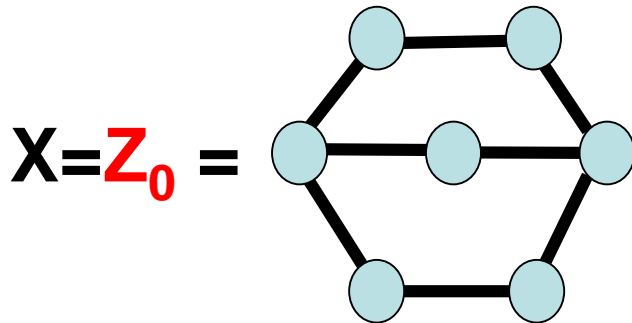
$X$  has a **cut-vertex**



$X$  is **bipartite** ( $n \geq 3$ )

$X$  is a **cycle** ( $n \geq 4$ )

$[FS(C_n, K_{1,n-1})$  has  $(n-2)!$  components]



$FS(Z_0, K_{1,6})$  has 6 components



**Theorem (Wilson (74)):** Let  $X$  be an  $n$  vertex graph,  $n \geq 3$ . Suppose  $X$  is biconnected, neither  $Z_0$  nor a cycle of length at least 4.

If  $X$  is **non-bipartite**, then  $FS(X, K_{1, n-1})$  is connected.

If  $X$  is **bipartite** then  $FS(X, K_{1, n-1})$  has exactly two connected components

Proof combines the **ear decomposition** of a 2-connected graph with some group theoretic arguments.

# Connectivity: Typical and Extremal Questions

## Minimum degree

**Question 1a:** what is the smallest  $d_n$  so that  $FS(X,Y)$  is connected for every two  $n$ -vertex graphs  $X,Y$  each having minimum degree at least  $d_n$  ?

**Question 1b:** what is the smallest  $d_{n,n}$  so that  $FS(X,Y)$  has exactly two connected components for every two subgraphs  $X,Y$  of  $K_{n,n}$  each having minimum degree at least  $d_{n,n}$  ?

# Random Graphs

**Question 2a:** Let  $X, Y$  be independent binomial random graphs in  $G(n, p)$ . For which  $p=p(n)$  is  $FS(X, Y)$  connected with high probability ?

**Question 2b:** Let  $X, Y$  be two independent bipartite random graphs in  $G(n, n, p)$ . For which  $p=p(n)$  does  $FS(X, Y)$  have exactly 2 connected components with high probability ?

# Results

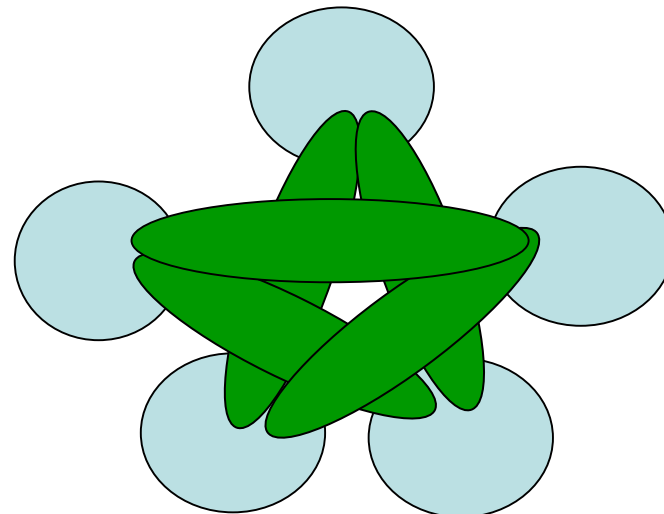
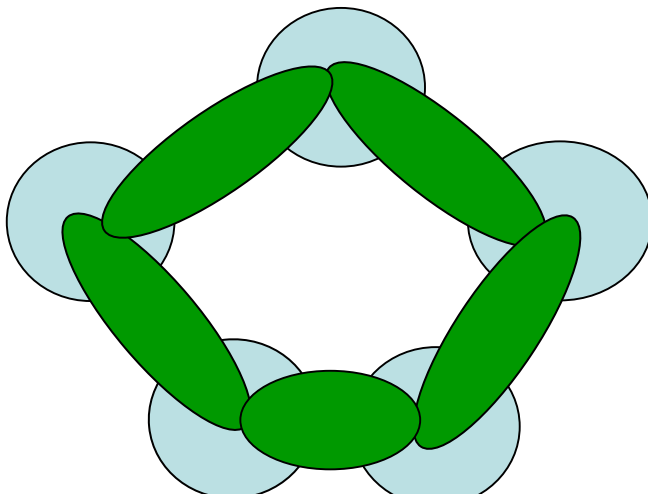
## Minimum degree

**Question 1a:** what is the smallest  $d_n$  so that  $FS(X,Y)$  is connected for every two  $n$ -vertex graphs  $X,Y$  each having minimum degree at least  $d_n$  ?

**Theorem (A,Defant,Kravitz):**

$$\frac{3n}{5} - 2 \leq d_n \leq \frac{9n}{14} + 1$$

**The lower bound**

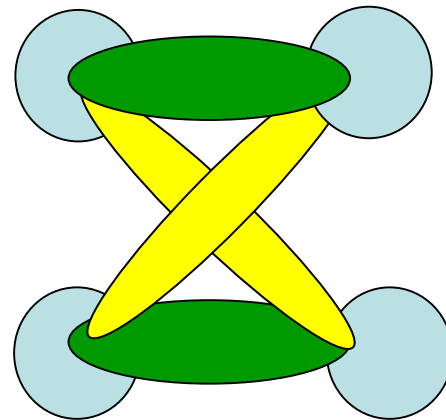
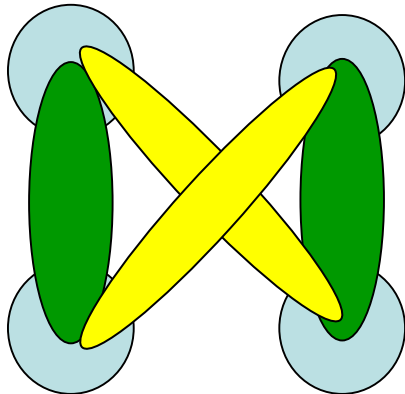


**Question 1b:** what is the smallest  $d_{n,n}$  so that  $FS(X,Y)$  has exactly two connected components for every two subgraphs  $X,Y$  of  $K_{n,n}$  each having minimum degree at least  $d_{n,n}$  ?

**Theorem (A, Defant, Kravitz):**

$$\left\lceil \frac{3n+1}{4} \right\rceil \leq d_{n,n} \leq \left\lceil \frac{3n+2}{4} \right\rceil$$

**The lower bound**



## Random Graphs

**Question 2a:** Let  $X, Y$  be independent binomial random graphs in  $G(n, p)$ . For which  $p=p(n)$  is  $FS(X, Y)$  connected with high probability ?

**Theorem (A, Defant, Kravitz):** The **threshold**  $p=p(n)$  for connectivity of  $FS(X, Y)$  is

$$p(n) = \frac{1}{n^{\frac{1}{2}+o(1)}}$$

**Question 2b:** Let  $X, Y$  be two independent bipartite random graphs in  $G(n, n, p)$ . For which  $p = p(n)$  does  $FS(X, Y)$  have exactly 2 connected components with high probability ?

**Theorem (A, Defant, Kravitz):** the **threshold**  $p(n)$  for having two components satisfies

$$\Omega\left(\frac{1}{n^{1/2}}\right) \leq p(n) \leq \tilde{O}\left(\frac{1}{n^{10}}\right)$$



# A bit about the proofs

**Theorem:** the threshold for connectivity of  $FS(X,Y)$  for  $X,Y$  in  $G(n,p(n))$  is  $n^{-1/2+o(1)}$

**Fact:** for  $p(n) \leq \frac{2^{-\frac{1}{2} - \epsilon}}{\sqrt{n}}$  then with high probability  $FS(X,Y)$  is disconnected (has isolated vertices).

**Proof:** Sauer and Spencer (78) showed that if

$$2 \Delta(X)\Delta(Y) < n$$

then  $X$  and  $Y$  have an edge disjoint **packing** in  $K_n$

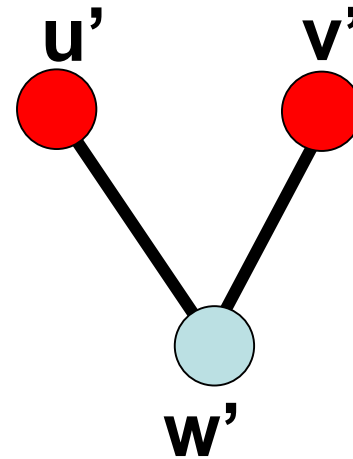
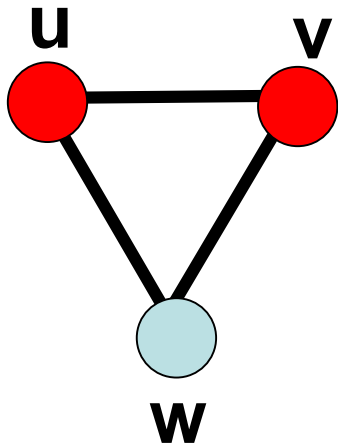
The main part of the proof is:

**Theorem:** if  $p(n) \geq \frac{\exp[(2 \log n)^{2/3}]}{n^{1/2}}$

then with high probability FS(X,Y) is connected.

This is proved by showing that with high probability, for every bijection  $f$  from  $V(X)$  to  $V(Y)$  every pair of **adjacent** elements  $u,v$  in  $Y$  are **exchangeable**.

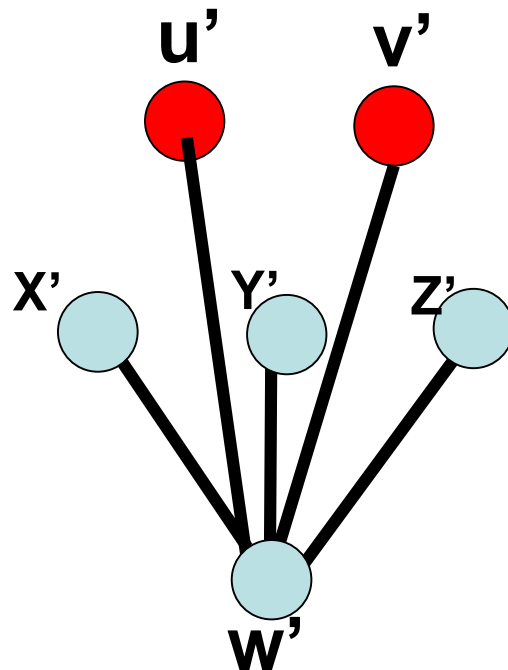
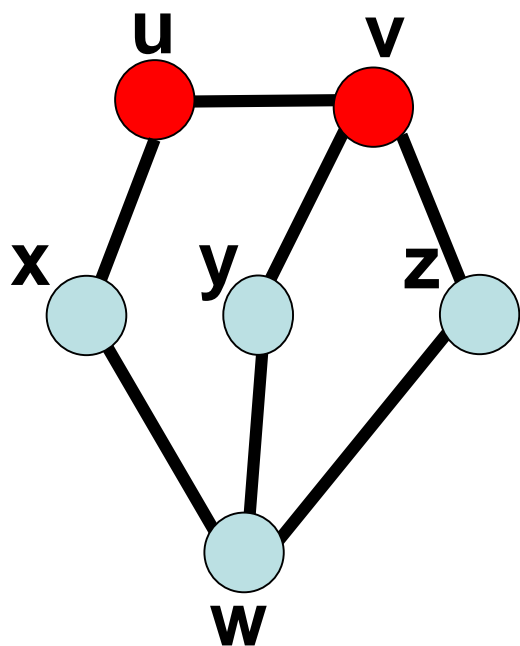
**First attempt** to establish that: hope that for every such  $u, v$  and  $f$  there is a common neighbor  $w$  of  $u$  and  $v$  in  $Y$  where  $w' = f^{-1}(w)$  is a common neighbor of  $u' = f^{-1}(u)$  and  $v' = f^{-1}(v)$ .



**But this fails with high probability for all**

$$p < 1/\sqrt{2}$$

**Second attempt:** show that for every such  $u, v$  and  $f$  the following two graphs appear in  $X, Y$

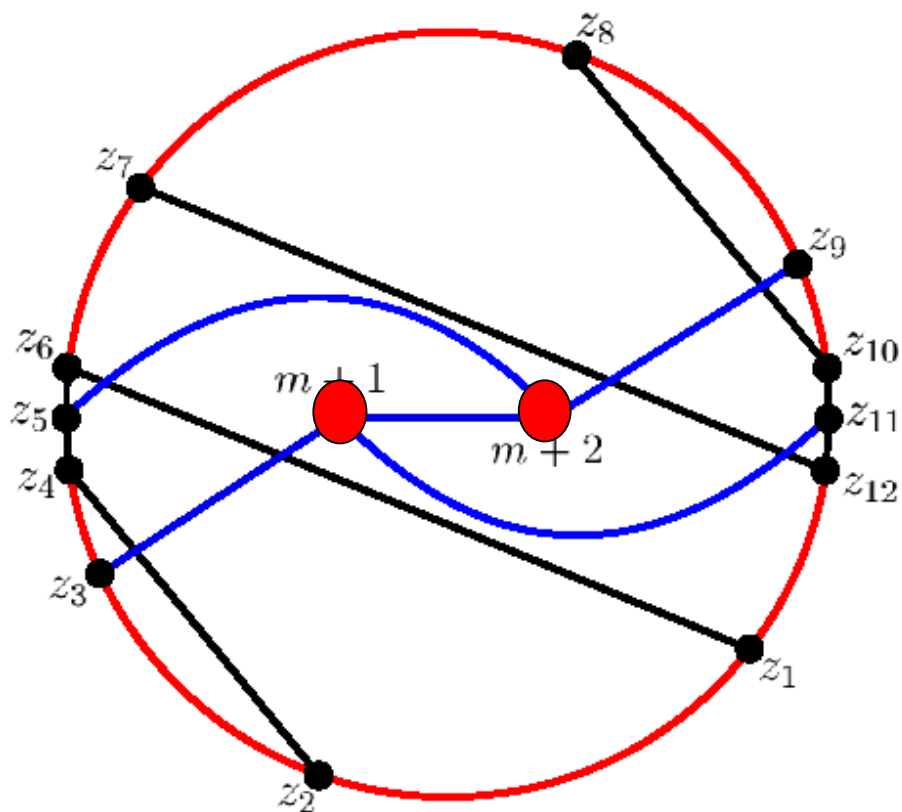


This will suffice by **Wilson's Theorem**

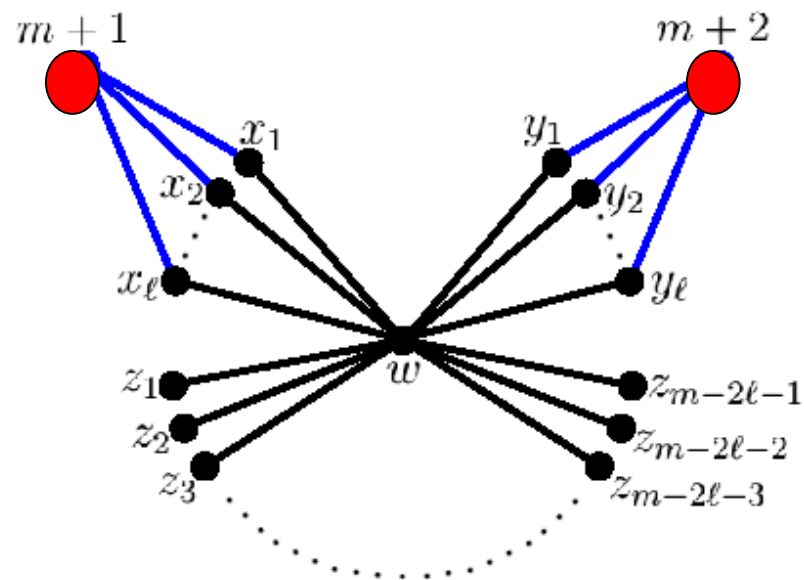
For  $p(n) \geq \tilde{\Omega}(1/n^{0.25})$  this holds with high probability, by **Janson Inequalities**

Yet it fails for  $p(n)=1/n^{1/2-o(1)}$

# Need more complicated graphs



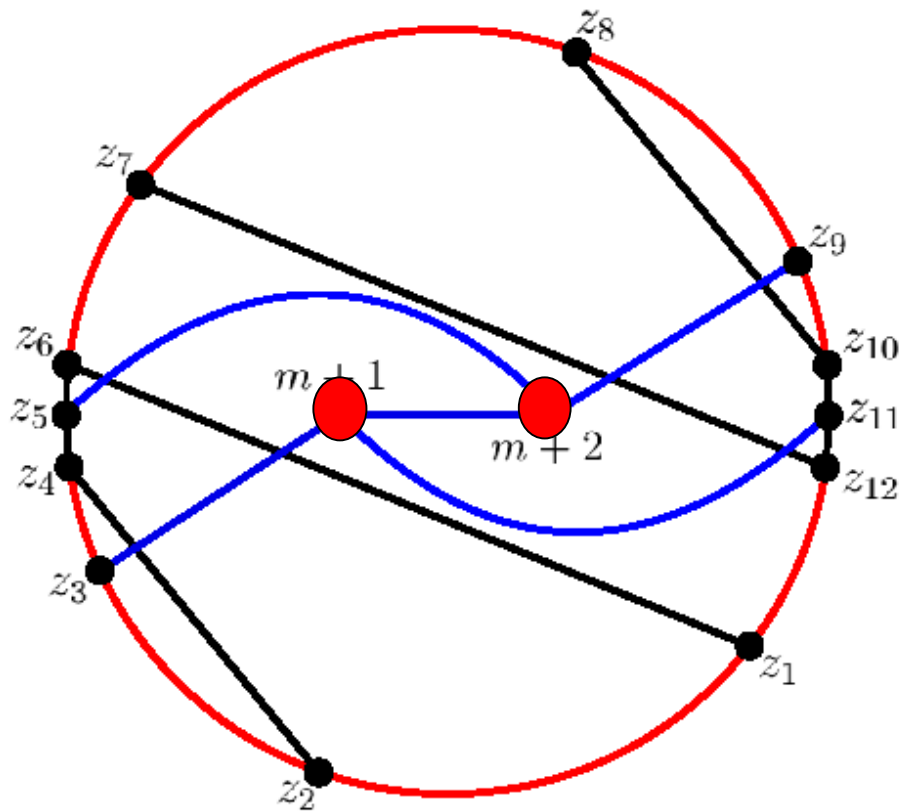
$G^*$



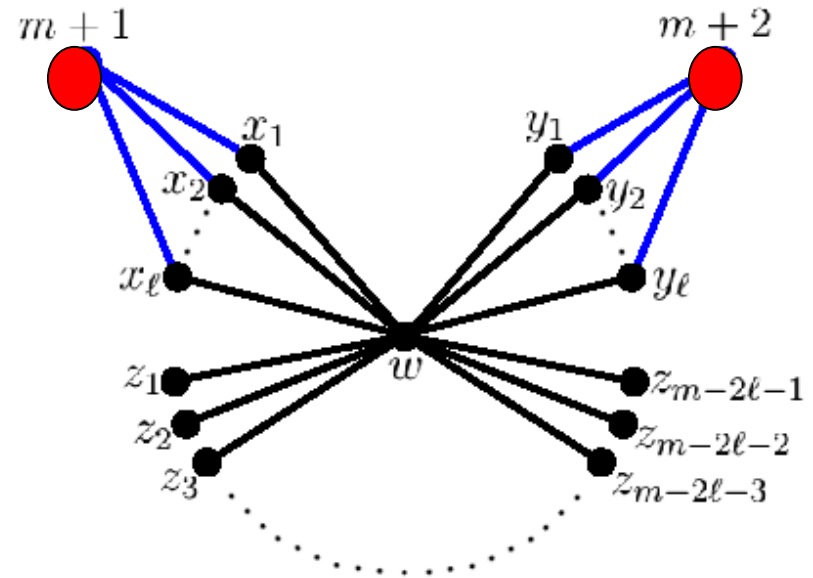
$H^*$

Here  $m = \lfloor \log n^{2/3} \rfloor$  grows with  $n$ .

# Need more complicated graphs



$G^*$

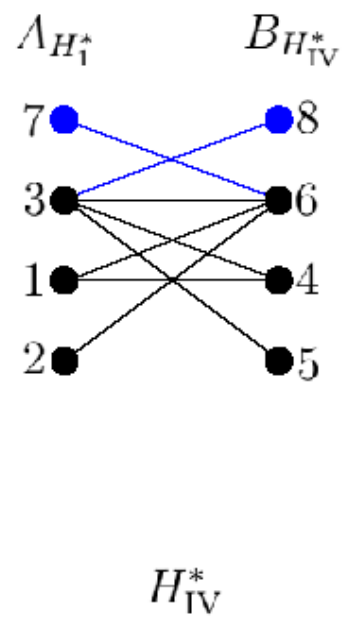
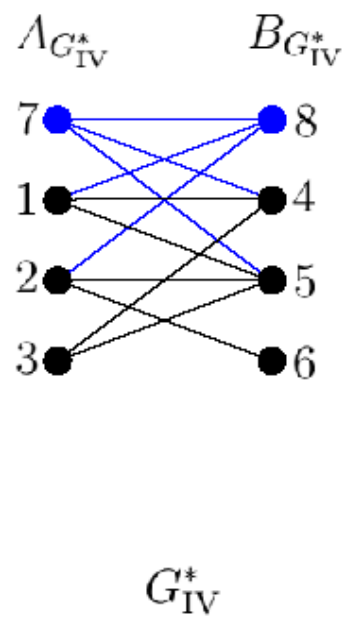
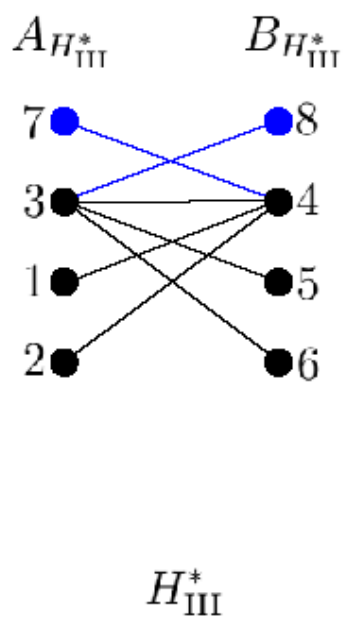
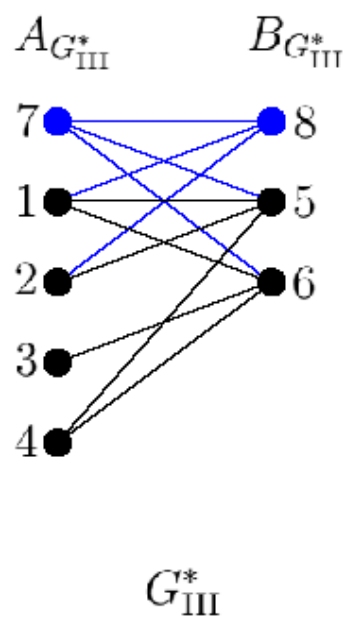
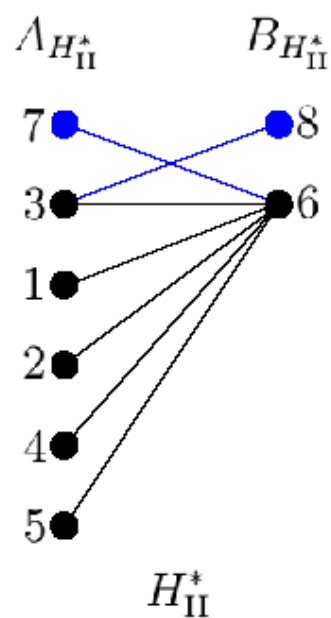
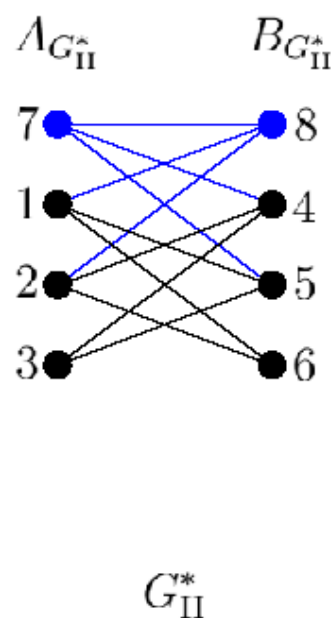
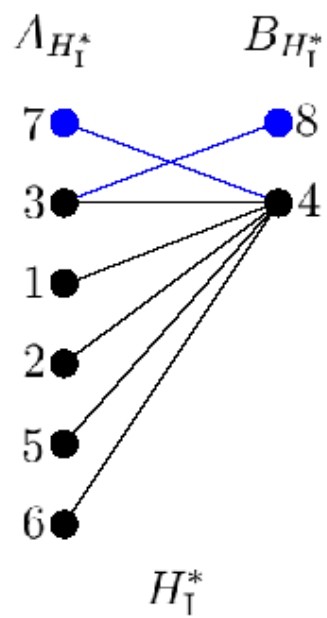
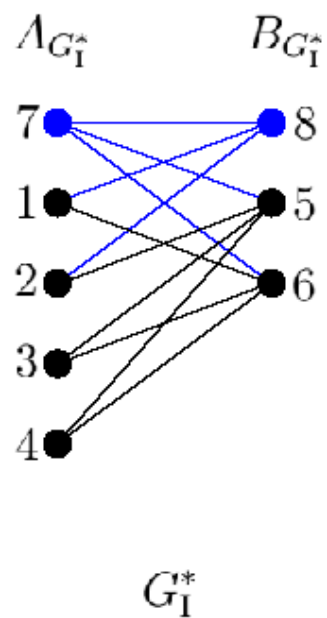


$H^*$

The proof these imply exchangeability  
applies **Wilson's Theorem** several times

**The (non-tight) result for the **bipartite** case is obtained by a similar reasoning using specific constant size pairs of graphs that supply exchangability by sequences of friendly swaps found by computer search.**





# Open Problems

Is  $p(n)=n^{-1/2+o(1)}$  the **threshold** for ensuring two connected components in the bipartite case too?

Is the smallest value  $d_n$  ensuring connectivity of  $FS(X, Y)$  for every pair of  $n$ -vertex graphs  $X$  and  $Y$  with **minimum degree** at least  $d_n$   $3n/5+O(1)$  ?

Is there a **hitting time result** ? Namely, starting with two edgeless graphs  $(X_0, Y_0)$  on  $n$  vertices each, let  $(X_i, Y_i)$  be a random sequence of pairs of graphs, where each  $X_{i+1}$  is obtained from  $X_i$  by adding to it a uniform random yet unchosen edge, and each  $Y_{i+1}$  is defined analogously.

Let  $t_{iso}$  denote the smallest  $i$  so that  $FS(X_i, Y_i)$  has **no isolated vertices**.

Let  $t_{con}$  denote the smallest  $i$  so that  $FS(X_i, Y_i)$  is **connected**.

Is  $t_{conn} = t_{iso}$  with high probability ? If not, is  $t_{con} = (1+o(1)) t_{iso}$  with high probability ?

Is the **diameter** of any connected component of  $FS(X, Y)$  at most  $n^{O(1)}$  ?

What about the mixing properties of the **random walk** on  $FS(X, Y)$  ?

[The case  $FS(X, K_n)$  is **Aldous spectral gap conjecture** settled by **Caputo, Liggett and Richthammer (10)** ]

